

Miscellaneous FOL Examples

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1 Natural numbers

```

theory Natural-Numbers
imports FOL
begin

```

Theory of the natural numbers: Peano's axioms, primitive recursion. (Modernized version of Larry Paulson's theory "Nat".)

```

typedecl nat
instance nat ::  $\langle term \rangle$  ..

```

axiomatization

```

Zero ::  $\langle nat \rangle$  ( $\langle 0 \rangle$ ) and
Suc ::  $\langle nat \Rightarrow nat \rangle$  and
rec ::  $\langle [nat, 'a, [nat, 'a] \Rightarrow 'a] \Rightarrow 'a \rangle$ 
where
  induct [case-names 0 Suc, induct type: nat]:
     $\langle P(0) \Rightarrow (!x. P(x) \Rightarrow P(Suc(x))) \Rightarrow P(n) \rangle$  and
  Suc-inject:  $\langle Suc(m) = Suc(n) \Rightarrow m = n \rangle$  and
  Suc-neq-0:  $\langle Suc(m) = 0 \Rightarrow R \rangle$  and
  rec-0:  $\langle rec(0, a, f) = a \rangle$  and

```

```

    rec-Suc:  $\langle \text{rec}(\text{Suc}(m), a, f) = f(m, \text{rec}(m, a, f)) \rangle$ 

lemma Suc-n-not-n:  $\langle \text{Suc}(k) \neq k \rangle$ 
proof (induct  $\langle k \rangle$ )
  show  $\langle \text{Suc}(0) \neq 0 \rangle$ 
  proof
    assume  $\langle \text{Suc}(0) = 0 \rangle$ 
    then show  $\langle \text{False} \rangle$  by (rule Suc-neq-0)
  qed
next
  fix n assume hyp:  $\langle \text{Suc}(n) \neq n \rangle$ 
  show  $\langle \text{Suc}(\text{Suc}(n)) \neq \text{Suc}(n) \rangle$ 
  proof
    assume  $\langle \text{Suc}(\text{Suc}(n)) = \text{Suc}(n) \rangle$ 
    then have  $\langle \text{Suc}(n) = n \rangle$  by (rule Suc-inject)
    with hyp show  $\langle \text{False} \rangle$  by contradiction
  qed
qed

definition add ::  $\langle \text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \rangle$  (infixl  $\langle + \rangle$  60)
  where  $\langle m + n = \text{rec}(m, n, \lambda x y. \text{Suc}(y)) \rangle$ 

lemma add-0 [simp]:  $\langle 0 + n = n \rangle$ 
  unfolding add-def by (rule rec-0)

lemma add-Suc [simp]:  $\langle \text{Suc}(m) + n = \text{Suc}(m + n) \rangle$ 
  unfolding add-def by (rule rec-Suc)

lemma add-assoc:  $\langle (k + m) + n = k + (m + n) \rangle$ 
  by (induct  $\langle k \rangle$ ) simp-all

lemma add-0-right:  $\langle m + 0 = m \rangle$ 
  by (induct  $\langle m \rangle$ ) simp-all

lemma add-Suc-right:  $\langle m + \text{Suc}(n) = \text{Suc}(m + n) \rangle$ 
  by (induct  $\langle m \rangle$ ) simp-all

lemma
  assumes  $\langle !n. f(\text{Suc}(n)) = \text{Suc}(f(n)) \rangle$ 
  shows  $\langle f(i + j) = i + f(j) \rangle$ 
  using assms by (induct  $\langle i \rangle$ ) simp-all

end

```

2 Examples for the manual “Introduction to Isabelle”

```
theory Intro
imports FOL
begin
```

2.0.1 Some simple backward proofs

```
lemma mythm:  $\langle P \vee P \longrightarrow P \rangle$ 
apply (rule impI)
apply (rule disjE)
prefer 3 apply (assumption)
prefer 2 apply (assumption)
apply assumption
done
```

```
lemma  $\langle (P \wedge Q) \vee R \longrightarrow (P \vee R) \rangle$ 
apply (rule impI)
apply (erule disjE)
apply (drule conjunct1)
apply (rule disjI1)
apply (rule-tac [2] disjI2)
apply assumption+
done
```

Correct version, delaying use of *spec* until last.

```
lemma  $\langle (\forall x y. P(x,y)) \longrightarrow (\forall z w. P(w,z)) \rangle$ 
apply (rule impI)
apply (rule allI)
apply (rule allI)
apply (drule spec)
apply (drule spec)
apply assumption
done
```

2.0.2 Demonstration of *fast*

```
lemma  $\langle (\exists y. \forall x. J(y,x) \longleftrightarrow \neg J(x,x)) \longrightarrow \neg (\forall x. \exists y. \forall z. J(z,y) \longleftrightarrow \neg J(z,x)) \rangle$ 
apply fast
done
```

```
lemma  $\langle \forall x. P(x,f(x)) \longleftrightarrow (\exists y. (\forall z. P(z,y) \longrightarrow P(z,f(x))) \wedge P(x,y)) \rangle$ 
apply fast
done
```

2.0.3 Derivation of conjunction elimination rule

```
lemma
```

```

    assumes major:  $\langle P \wedge Q \rangle$ 
    and minor:  $\langle \llbracket P; Q \rrbracket \Longrightarrow R \rangle$ 
    shows  $\langle R \rangle$ 
  apply (rule minor)
  apply (rule major [THEN conjunct1])
  apply (rule major [THEN conjunct2])
done

```

2.1 Derived rules involving definitions

Derivation of negation introduction

```

lemma
  assumes  $\langle P \Longrightarrow False \rangle$ 
  shows  $\langle \neg P \rangle$ 
  apply (unfold not-def)
  apply (rule impI)
  apply (rule assms)
  apply assumption
done

```

```

lemma
  assumes major:  $\langle \neg P \rangle$ 
  and minor:  $\langle P \rangle$ 
  shows  $\langle R \rangle$ 
  apply (rule FalseE)
  apply (rule mp)
  apply (rule major [unfolded not-def])
  apply (rule minor)
done

```

Alternative proof of the result above

```

lemma
  assumes major:  $\langle \neg P \rangle$ 
  and minor:  $\langle P \rangle$ 
  shows  $\langle R \rangle$ 
  apply (rule minor [THEN major [unfolded not-def, THEN mp, THEN FalseE]])
done

```

end

3 Theory of the natural numbers: Peano's axioms, primitive recursion

```

theory Nat
  imports FOL
begin

```

```

typedecl nat
instance nat :: <term> ..

```

axiomatization

```

  Zero :: <nat> (<0>) and
  Suc :: <nat  $\Rightarrow$  nat> and
  rec :: <[nat, 'a, [nat, 'a]  $\Rightarrow$  'a]  $\Rightarrow$  'a>
where
  induct: <[[P(0);  $\bigwedge x. P(x) \Rightarrow P(\text{Suc}(x))$ ]]  $\Rightarrow$  P(n)> and
  Suc-inject: <Suc(m)=Suc(n)  $\Rightarrow$  m=n> and
  Suc-neq-0: <Suc(m)=0  $\Rightarrow$  R> and
  rec-0: <rec(0,a,f) = a> and
  rec-Suc: <rec(Suc(m), a, f) = f(m, rec(m,a,f))>

```

```

definition add :: <[nat, nat]  $\Rightarrow$  nat> (infixl <+> 60)
  where <m + n  $\equiv$  rec(m, n,  $\lambda x y. \text{Suc}(y)$ )>

```

3.1 Proofs about the natural numbers

```

lemma Suc-n-not-n: <Suc(k)  $\neq$  k>
apply (rule-tac n = <k> in induct)
apply (rule notI)
apply (erule Suc-neq-0)
apply (rule notI)
apply (erule notE)
apply (erule Suc-inject)
done

```

```

lemma <(k+m)+n = k+(m+n)>
apply (rule induct)
back
back
back
back
back
back
oops

```

```

lemma add-0 [simp]: <0+n = n>
apply (unfold add-def)
apply (rule rec-0)
done

```

```

lemma add-Suc [simp]: <Suc(m)+n = Suc(m+n)>
apply (unfold add-def)
apply (rule rec-Suc)
done

```

```

lemma add-assoc: <(k+m)+n = k+(m+n)>

```

```

apply (rule-tac  $n = \langle k \rangle$  in induct)
apply simp
apply simp
done

lemma add-0-right:  $\langle m+0 = m \rangle$ 
apply (rule-tac  $n = \langle m \rangle$  in induct)
apply simp
apply simp
done

lemma add-Suc-right:  $\langle m+\text{Suc}(n) = \text{Suc}(m+n) \rangle$ 
apply (rule-tac  $n = \langle m \rangle$  in induct)
apply simp-all
done

lemma
  assumes prem:  $\langle \bigwedge n. f(\text{Suc}(n)) = \text{Suc}(f(n)) \rangle$ 
  shows  $\langle f(i+j) = i+f(j) \rangle$ 
apply (rule-tac  $n = \langle i \rangle$  in induct)
apply simp
apply (simp add: prem)
done

end

```

4 Theory of the natural numbers: Peano's axioms, primitive recursion

```

theory Nat-Class
  imports FOL
begin

```

This is an abstract version of `Nat.thy`. Instead of axiomatizing a single type *nat*, it defines the class of all these types (up to isomorphism).

Note: The *rec* operator has been made *monomorphic*, because class axioms cannot contain more than one type variable.

```

class nat =
  fixes Zero ::  $\langle 'a \rangle \ (\langle 0 \rangle)$ 
  and Suc ::  $\langle 'a \Rightarrow 'a \rangle$ 
  and rec ::  $\langle 'a \Rightarrow 'a \Rightarrow ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'a \rangle$ 
  assumes induct:  $\langle P(0) \Longrightarrow (\bigwedge x. P(x) \Longrightarrow P(\text{Suc}(x))) \Longrightarrow P(n) \rangle$ 
  and Suc-inject:  $\langle \text{Suc}(m) = \text{Suc}(n) \Longrightarrow m = n \rangle$ 
  and Suc-neq-Zero:  $\langle \text{Suc}(m) = 0 \Longrightarrow \text{False} \rangle$ 
  and rec-Zero:  $\langle \text{rec}(0, a, f) = a \rangle$ 
  and rec-Suc:  $\langle \text{rec}(\text{Suc}(m), a, f) = f(m, \text{rec}(m, a, f)) \rangle$ 
begin

```

definition $add :: \langle 'a \Rightarrow 'a \Rightarrow 'a \rangle$ (**infixl** $\langle + \rangle$ 60)
where $\langle m + n = rec(m, n, \lambda x y. Suc(y)) \rangle$

lemma $Suc\text{-}n\text{-}not\text{-}n$: $\langle Suc(k) \neq (k::'a) \rangle$
apply (*rule-tac* $n = \langle k \rangle$ **in** *induct*)
apply (*rule notI*)
apply (*erule Suc-neq-Zero*)
apply (*rule notI*)
apply (*erule notE*)
apply (*erule Suc-inject*)
done

lemma $\langle (k + m) + n = k + (m + n) \rangle$
apply (*rule induct*)
back
back
back
back
back
oops

lemma $add\text{-}Zero$ [*simp*]: $\langle 0 + n = n \rangle$
apply (*unfold add-def*)
apply (*rule rec-Zero*)
done

lemma $add\text{-}Suc$ [*simp*]: $\langle Suc(m) + n = Suc(m + n) \rangle$
apply (*unfold add-def*)
apply (*rule rec-Suc*)
done

lemma $add\text{-}assoc$: $\langle (k + m) + n = k + (m + n) \rangle$
apply (*rule-tac* $n = \langle k \rangle$ **in** *induct*)
apply *simp*
apply *simp*
done

lemma $add\text{-}Zero\text{-}right$: $\langle m + 0 = m \rangle$
apply (*rule-tac* $n = \langle m \rangle$ **in** *induct*)
apply *simp*
apply *simp*
done

lemma $add\text{-}Suc\text{-}right$: $\langle m + Suc(n) = Suc(m + n) \rangle$
apply (*rule-tac* $n = \langle m \rangle$ **in** *induct*)
apply *simp-all*
done

lemma


```

    assumes prem:  $\langle \bigwedge n. f(\text{Suc}(n)) = \text{Suc}(f(n)) \rangle$ 
    shows  $\langle f(i + j) = i + f(j) \rangle$ 
    apply (rule-tac  $n = \langle i \rangle$  in induct)
    apply simp
    apply (simp add: prem)
    done

end

end

```

5 Intuitionistic FOL: Examples from The Foundation of a Generic Theorem Prover

```

theory Foundation
imports IFOL
begin

lemma  $\langle A \wedge B \longrightarrow (C \longrightarrow A \wedge C) \rangle$ 
  apply (rule impI)
  apply (rule impI)
  apply (rule conjI)
  prefer 2 apply assumption
  apply (rule conjunct1)
  apply assumption
  done

```

A form of conj-elimination

```

lemma
  assumes  $\langle A \wedge B \rangle$ 
  and  $\langle A \Longrightarrow B \Longrightarrow C \rangle$ 
  shows  $\langle C \rangle$ 
  apply (rule assms)
  apply (rule conjunct1)
  apply (rule assms)
  apply (rule conjunct2)
  apply (rule assms)
  done

```

```

lemma
  assumes  $\langle \bigwedge A. \neg \neg A \Longrightarrow A \rangle$ 
  shows  $\langle B \vee \neg B \rangle$ 
  apply (rule assms)
  apply (rule notI)
  apply (rule-tac  $P = \langle \neg B \rangle$  in notE)
  apply (rule-tac [2] notI)
  apply (rule-tac [2]  $P = \langle B \vee \neg B \rangle$  in notE)
  prefer 2 apply assumption

```

```

apply (rule-tac [2] disjI1)
prefer 2 apply assumption
apply (rule notI)
apply (rule-tac  $P = \langle B \vee \neg B \rangle$  in notE)
apply assumption
apply (rule disjI2)
apply assumption
done

```

```

lemma
  assumes  $\langle \bigwedge A. \neg \neg A \implies A \rangle$ 
  shows  $\langle B \vee \neg B \rangle$ 
apply (rule assms)
apply (rule notI)
apply (rule notE)
apply (rule-tac [2] notI)
apply (erule-tac [2] notE)
apply (erule-tac [2] disjI1)
apply (rule notI)
apply (erule notE)
apply (erule disjI2)
done

```

```

lemma
  assumes  $\langle A \vee \neg A \rangle$ 
  and  $\langle \neg \neg A \rangle$ 
  shows  $\langle A \rangle$ 
apply (rule disjE)
apply (rule assms)
apply assumption
apply (rule FalseE)
apply (rule-tac  $P = \langle \neg A \rangle$  in notE)
apply (rule assms)
apply assumption
done

```

5.1 Examples with quantifiers

```

lemma
  assumes  $\langle \forall z. G(z) \rangle$ 
  shows  $\langle \forall z. G(z) \vee H(z) \rangle$ 
apply (rule allI)
apply (rule disjI1)
apply (rule assms [THEN spec])
done

```

```

lemma  $\langle \forall x. \exists y. x = y \rangle$ 
apply (rule allI)

```

```

apply (rule exI)
apply (rule refl)
done

```

```

lemma  $\langle \exists y. \forall x. x = y \rangle$ 
apply (rule exI)
apply (rule allI)
apply (rule refl)?
oops

```

Parallel lifting example.

```

lemma  $\langle \exists u. \forall x. \exists v. \forall y. \exists w. P(u, x, v, y, w) \rangle$ 
apply (rule exI allI)
apply (rule exI allI)
apply (rule exI allI)
apply (rule exI allI)
apply (rule exI allI)
oops

```

```

lemma
  assumes  $\langle (\exists z. F(z)) \wedge B \rangle$ 
  shows  $\langle \exists z. F(z) \wedge B \rangle$ 
apply (rule conjE)
apply (rule assms)
apply (rule exE)
apply assumption
apply (rule exI)
apply (rule conjI)
apply assumption
apply assumption
done

```

A bigger demonstration of quantifiers – not in the paper.

```

lemma  $\langle (\exists y. \forall x. Q(x, y)) \longrightarrow (\forall x. \exists y. Q(x, y)) \rangle$ 
apply (rule impI)
apply (rule allI)
apply (rule exE, assumption)
apply (rule exI)
apply (rule allE, assumption)
apply assumption
done

```

end

6 First-Order Logic: PROLOG examples

```

theory Prolog
imports FOL

```

```

begin

typedecl 'a list
instance list :: (⟨term⟩) ⟨term⟩ ..

axiomatization
  Nil    :: ⟨'a list⟩ and
  Cons   :: ⟨['a, 'a list] => 'a list⟩ (infixr ⟨:⟩ 60) and
  app    :: ⟨['a list, 'a list, 'a list] => o⟩ and
  rev    :: ⟨['a list, 'a list] => o⟩
where
  appNil: ⟨app(Nil,ys,ys)⟩ and
  appCons: ⟨app(xs,ys,zs) ==> app(x:xs, ys, x:zs)⟩ and
  revNil: ⟨rev(Nil,Nil)⟩ and
  revCons: ⟨[| rev(xs,ys); app(ys, x:Nil, zs) |] ==> rev(x:xs, zs)⟩

schematic-goal ⟨app(a:b:c:Nil, d:e:Nil, ?x)⟩
apply (rule appNil appCons)
apply (rule appNil appCons)
apply (rule appNil appCons)
apply (rule appNil appCons)
done

schematic-goal ⟨app(?x, c:d:Nil, a:b:c:d:Nil)⟩
apply (rule appNil appCons)+
done

schematic-goal ⟨app(?x, ?y, a:b:c:d:Nil)⟩
apply (rule appNil appCons)+
back
back
back
back
done

lemmas rules = appNil appCons revNil revCons

schematic-goal ⟨rev(a:b:c:d:Nil, ?x)⟩
apply (rule rules)+
done

schematic-goal ⟨rev(a:b:c:d:e:f:g:h:i:j:k:l:m:n:Nil, ?w)⟩
apply (rule rules)+
done

schematic-goal ⟨rev(?x, a:b:c:Nil)⟩

```

```

apply (rule rules)+ — does not solve it directly!
back
back
done

```

```

ML <
  fun prolog-tac ctxt =
    DEPTH-FIRST Thm.no-prems (resolve-tac ctxt @ {thms rules} 1)
  >

```

```

schematic-goal <rev(?x, a:b:c:Nil)>
apply (tactic <prolog-tac context>)
done

```

```

schematic-goal <rev(a:?x:c:?y:Nil, d:?z:b:?u)>
apply (tactic <prolog-tac context>)
done

```

```

schematic-goal <rev(a:b:c:d:e:f:g:h:i:j:k:l:m:n:o:p:Nil, ?w)>
apply (tactic <
  DEPTH-SOLVE (resolve-tac context ([@ {thm refl}, @ {thm conjI}] @ @ {thms
rules})) 1)>)
done

```

```

schematic-goal <a:b:c:d:e:f:g:h:i:j:k:l:m:n:o:p:Nil = ?x ∧ app(?x, ?x, ?y) ∧ rev(?y, ?w)>
apply (tactic <
  DEPTH-SOLVE (resolve-tac context ([@ {thm refl}, @ {thm conjI}] @ @ {thms
rules})) 1)>)
done

```

```

end

```

7 Intuitionistic First-Order Logic

```

theory Intuitionistic
imports IFOL
begin

```

Metatheorem (for *propositional* formulae): P is classically provable iff $\neg\neg P$ is intuitionistically provable. Therefore $\neg P$ is classically provable iff it is intuitionistically provable.

Proof: Let Q be the conjunction of the propositions $A \vee \neg A$, one for each atom A in P . Now $\neg\neg Q$ is intuitionistically provable because $\neg\neg(A \vee \neg A)$ is and because double-negation distributes over conjunction. If P is provable classically, then clearly $Q \rightarrow P$ is provable intuitionistically, so $\neg\neg(Q \rightarrow P)$

is also provable intuitionistically. The latter is intuitionistically equivalent to $\neg\neg Q \rightarrow \neg\neg P$, hence to $\neg\neg P$, since $\neg\neg Q$ is intuitionistically provable. Finally, if P is a negation then $\neg\neg P$ is intuitionistically equivalent to P . [Andy Pitts]

lemma $\langle \neg \neg (P \wedge Q) \longleftrightarrow \neg \neg P \wedge \neg \neg Q \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle \neg \neg ((\neg P \rightarrow Q) \rightarrow (\neg P \rightarrow \neg Q) \rightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

Double-negation does NOT distribute over disjunction.

lemma $\langle \neg \neg (P \rightarrow Q) \longleftrightarrow (\neg \neg P \rightarrow \neg \neg Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle \neg \neg \neg P \longleftrightarrow \neg P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle \neg \neg ((P \rightarrow Q \vee R) \rightarrow (P \rightarrow Q) \vee (P \rightarrow R)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle (P \longleftrightarrow Q) \longleftrightarrow (Q \longleftrightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle ((P \rightarrow (Q \vee (Q \rightarrow R))) \rightarrow R) \rightarrow R \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma
 $\langle (((G \rightarrow A) \rightarrow J) \rightarrow D \rightarrow E) \rightarrow (((H \rightarrow B) \rightarrow I) \rightarrow C \rightarrow J)$
 $\rightarrow (A \rightarrow H) \rightarrow F \rightarrow G \rightarrow (((C \rightarrow B) \rightarrow I) \rightarrow D) \rightarrow (A \rightarrow C)$
 $\rightarrow (((F \rightarrow A) \rightarrow B) \rightarrow I) \rightarrow E \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

Admissibility of the excluded middle for negated formulae

lemma $\langle (P \vee \neg P \rightarrow \neg Q) \rightarrow \neg Q \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

The same in a more general form, no ex falso quodlibet

lemma $\langle (P \vee (P \rightarrow R) \rightarrow Q \rightarrow R) \rightarrow Q \rightarrow R \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

7.1 Lemmas for the propositional double-negation translation

lemma $\langle P \rightarrow \neg \neg P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle \neg \neg (\neg \neg P \rightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle \neg \neg P \wedge \neg \neg (P \longrightarrow Q) \longrightarrow \neg \neg Q \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac}$ **context** 1 \rangle)

The following are classically but not constructively valid. The attempt to prove them terminates quickly!

lemma $\langle ((P \longrightarrow Q) \longrightarrow P) \longrightarrow P \rangle$
apply (*tactic* $\langle \text{IntPr.fast-tac}$ **context** 1 \rangle)?
apply (*rule asm-rl*) — Checks that subgoals remain: proof failed.
oops

lemma $\langle (P \wedge Q \longrightarrow R) \longrightarrow (P \longrightarrow R) \vee (Q \longrightarrow R) \rangle$
apply (*tactic* $\langle \text{IntPr.fast-tac}$ **context** 1 \rangle)?
apply (*rule asm-rl*) — Checks that subgoals remain: proof failed.
oops

7.2 de Bruijn formulae

de Bruijn formula with three predicates

lemma
 $\langle ((P \longleftrightarrow Q) \longrightarrow P \wedge Q \wedge R) \wedge$
 $((Q \longleftrightarrow R) \longrightarrow P \wedge Q \wedge R) \wedge$
 $((R \longleftrightarrow P) \longrightarrow P \wedge Q \wedge R) \longrightarrow P \wedge Q \wedge R \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac}$ **context** 1 \rangle)

de Bruijn formula with five predicates

lemma
 $\langle ((P \longleftrightarrow Q) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \wedge$
 $((Q \longleftrightarrow R) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \wedge$
 $((R \longleftrightarrow S) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \wedge$
 $((S \longleftrightarrow T) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \wedge$
 $((T \longleftrightarrow P) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T) \longrightarrow P \wedge Q \wedge R \wedge S \wedge T \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac}$ **context** 1 \rangle)

Problems from of Sahlin, Franzen and Haridi, An Intuitionistic Predicate Logic Theorem Prover. J. Logic and Comp. 2 (5), October 1992, 619-656.

Problem 1.1

lemma
 $\langle (\forall x. \exists y. \forall z. p(x) \wedge q(y) \wedge r(z)) \longleftrightarrow$
 $(\forall z. \exists y. \forall x. p(x) \wedge q(y) \wedge r(z)) \rangle$
by (*tactic* $\langle \text{IntPr.best-dup-tac}$ **context** 1 \rangle) — SLOW

Problem 3.1

lemma $\langle \neg (\exists x. \forall y. \text{mem}(y, x) \longleftrightarrow \neg \text{mem}(x, x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac}$ **context** 1 \rangle)

Problem 4.1: hopeless!

lemma
 $\langle (\forall x. p(x) \longrightarrow p(h(x)) \vee p(g(x))) \wedge (\exists x. p(x)) \wedge (\forall x. \neg p(h(x)))$
 $\longrightarrow (\exists x. p(g(g(g(g(x)))))) \rangle$
oops

7.3 Intuitionistic FOL: propositional problems based on Pelletier.

$\neg\neg 1$

lemma $\langle \neg \neg ((P \longrightarrow Q) \longleftrightarrow (\neg Q \longrightarrow \neg P)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 2$

lemma $\langle \neg \neg (\neg \neg P \longleftrightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

3

lemma $\langle \neg (P \longrightarrow Q) \longrightarrow (Q \longrightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 4$

lemma $\langle \neg \neg ((\neg P \longrightarrow Q) \longleftrightarrow (\neg Q \longrightarrow P)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 5$

lemma $\langle \neg \neg ((P \vee Q \longrightarrow P \vee R) \longrightarrow P \vee (Q \longrightarrow R)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 6$

lemma $\langle \neg \neg (P \vee \neg P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 7$

lemma $\langle \neg \neg (P \vee \neg \neg P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg 8$. Peirce's law

lemma $\langle \neg \neg (((P \longrightarrow Q) \longrightarrow P) \longrightarrow P) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

9

lemma $\langle ((P \vee Q) \wedge (\neg P \vee Q) \wedge (P \vee \neg Q)) \longrightarrow \neg (\neg P \vee \neg Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

10

lemma $\langle (Q \longrightarrow R) \longrightarrow (R \longrightarrow P \wedge Q) \longrightarrow (P \longrightarrow (Q \vee R)) \longrightarrow (P \longleftrightarrow Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

7.4 11. Proved in each direction (incorrectly, says Pelletier!!)

lemma $\langle P \longleftrightarrow P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg$ 12. Dijkstra's law

lemma $\langle \neg \neg (((P \longleftrightarrow Q) \longleftrightarrow R) \longleftrightarrow (P \longleftrightarrow (Q \longleftrightarrow R))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle ((P \longleftrightarrow Q) \longleftrightarrow R) \longrightarrow \neg \neg (P \longleftrightarrow (Q \longleftrightarrow R)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

13. Distributive law

lemma $\langle P \vee (Q \wedge R) \longleftrightarrow (P \vee Q) \wedge (P \vee R) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg$ 14

lemma $\langle \neg \neg ((P \longleftrightarrow Q) \longleftrightarrow ((Q \vee \neg P) \wedge (\neg Q \vee P))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg$ 15

lemma $\langle \neg \neg ((P \longrightarrow Q) \longleftrightarrow (\neg P \vee Q)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg$ 16

lemma $\langle \neg \neg ((P \longrightarrow Q) \vee (Q \longrightarrow P)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

$\neg\neg$ 17

lemma $\langle \neg \neg (((P \wedge (Q \longrightarrow R)) \longrightarrow S) \longleftrightarrow ((\neg P \vee Q \vee S) \wedge (\neg P \vee \neg R \vee S))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

Dijkstra's "Golden Rule"

lemma $\langle (P \wedge Q) \longleftrightarrow P \longleftrightarrow Q \longleftrightarrow (P \vee Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

8 Examples with quantifiers

8.1 The converse is classical in the following implications ...

lemma $\langle (\exists x. P(x) \longrightarrow Q) \longrightarrow (\forall x. P(x)) \longrightarrow Q \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle ((\forall x. P(x)) \longrightarrow Q) \longrightarrow \neg (\forall x. P(x) \wedge \neg Q) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle ((\forall x. \neg P(x)) \longrightarrow Q) \longrightarrow \neg (\forall x. \neg (P(x) \vee Q)) \rangle$

```

by (tactic ⟨IntPr.fast-tac context 1⟩)

lemma ⟨(∀ x. P(x)) ∨ Q ⟶ (∀ x. P(x) ∨ Q)⟩
  by (tactic ⟨IntPr.fast-tac context 1⟩)

lemma ⟨(∃ x. P ⟶ Q(x)) ⟶ (P ⟶ (∃ x. Q(x)))⟩
  by (tactic ⟨IntPr.fast-tac context 1⟩)

```

8.2 The following are not constructively valid!

The attempt to prove them terminates quickly!

```

lemma ⟨(∀ x. P(x)) ⟶ Q ⟶ (∃ x. P(x) ⟶ Q)⟩
  apply (tactic ⟨IntPr.fast-tac context 1⟩)?
  apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

```

```

lemma ⟨(P ⟶ (∃ x. Q(x))) ⟶ (∃ x. P ⟶ Q(x))⟩
  apply (tactic ⟨IntPr.fast-tac context 1⟩)?
  apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

```

```

lemma ⟨(∀ x. P(x) ∨ Q) ⟶ ((∀ x. P(x)) ∨ Q)⟩
  apply (tactic ⟨IntPr.fast-tac context 1⟩)?
  apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

```

```

lemma ⟨(∀ x. ¬ ¬ P(x)) ⟶ ¬ ¬ (∀ x. P(x))⟩
  apply (tactic ⟨IntPr.fast-tac context 1⟩)?
  apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

```

Classically but not intuitionistically valid. Proved by a bug in 1986!

```

lemma ⟨∃ x. Q(x) ⟶ (∀ x. Q(x))⟩
  apply (tactic ⟨IntPr.fast-tac context 1⟩)?
  apply (rule asm-rl) — Checks that subgoals remain: proof failed.
oops

```

8.3 Hard examples with quantifiers

The ones that have not been proved are not known to be valid! Some will require quantifier duplication – not currently available.

¬¬18

```

lemma ⟨¬ ¬ (∃ y. ∀ x. P(y) ⟶ P(x))⟩
  oops — NOT PROVED

```

¬¬19

```

lemma ⟨¬ ¬ (∃ x. ∀ y z. (P(y) ⟶ Q(z)) ⟶ (P(x) ⟶ Q(x)))⟩

```

oops — NOT PROVED

20

lemma

$\langle (\forall x y. \exists z. \forall w. (P(x) \wedge Q(y) \longrightarrow R(z) \wedge S(w)))$
 $\longrightarrow (\exists x y. P(x) \wedge Q(y)) \longrightarrow (\exists z. R(z)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

21

lemma $\langle (\exists x. P \longrightarrow Q(x)) \wedge (\exists x. Q(x) \longrightarrow P) \longrightarrow \neg \neg (\exists x. P \longleftrightarrow Q(x)) \rangle$
oops — NOT PROVED; needs quantifier duplication

22

lemma $\langle (\forall x. P \longleftrightarrow Q(x)) \longrightarrow (P \longleftrightarrow (\forall x. Q(x))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

$\neg \neg$ 23

lemma $\langle \neg \neg ((\forall x. P \vee Q(x)) \longleftrightarrow (P \vee (\forall x. Q(x)))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

24

lemma

$\langle \neg (\exists x. S(x) \wedge Q(x)) \wedge (\forall x. P(x) \longrightarrow Q(x) \vee R(x)) \wedge$
 $(\neg (\exists x. P(x)) \longrightarrow (\exists x. Q(x))) \wedge (\forall x. Q(x) \vee R(x) \longrightarrow S(x))$
 $\longrightarrow \neg \neg (\exists x. P(x) \wedge R(x)) \rangle$

Not clear why *fast-tac*, *best-tac*, *ASTAR* and *ITER-DEEPEN* all take forever.

apply (*tactic* $\langle \text{IntPr.safe-tac } \mathbf{context} \rangle$)
apply (*erule impE*)
apply (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)
apply (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)
done

25

lemma

$\langle (\exists x. P(x)) \wedge$
 $(\forall x. L(x) \longrightarrow \neg (M(x) \wedge R(x))) \wedge$
 $(\forall x. P(x) \longrightarrow (M(x) \wedge L(x))) \wedge$
 $((\forall x. P(x) \longrightarrow Q(x)) \vee (\exists x. P(x) \wedge R(x)))$
 $\longrightarrow (\exists x. Q(x) \wedge P(x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

$\neg \neg$ 26

lemma

$\langle (\neg \neg (\exists x. p(x)) \longleftrightarrow \neg \neg (\exists x. q(x))) \wedge$
 $(\forall x. \forall y. p(x) \wedge q(y) \longrightarrow (r(x) \longleftrightarrow s(y))) \rangle$

$\longrightarrow ((\forall x. p(x) \longrightarrow r(x)) \longleftrightarrow (\forall x. q(x) \longrightarrow s(x))) \rangle$
oops — NOT PROVED

27

lemma

$\langle (\exists x. P(x) \wedge \neg Q(x)) \wedge$
 $(\forall x. P(x) \longrightarrow R(x)) \wedge$
 $(\forall x. M(x) \wedge L(x) \longrightarrow P(x)) \wedge$
 $((\exists x. R(x) \wedge \neg Q(x)) \longrightarrow (\forall x. L(x) \longrightarrow \neg R(x)))$
 $\longrightarrow (\forall x. M(x) \longrightarrow \neg L(x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \textbf{context } 1 \rangle$)

$\neg\neg$ 28. AMENDED

lemma

$\langle (\forall x. P(x) \longrightarrow (\forall x. Q(x))) \wedge$
 $(\neg\neg (\forall x. Q(x) \vee R(x)) \longrightarrow (\exists x. Q(x) \wedge S(x))) \wedge$
 $(\neg\neg (\exists x. S(x)) \longrightarrow (\forall x. L(x) \longrightarrow M(x)))$
 $\longrightarrow (\forall x. P(x) \wedge L(x) \longrightarrow M(x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \textbf{context } 1 \rangle$)

29. Essentially the same as Principia Mathematica *11.71

lemma

$\langle (\exists x. P(x)) \wedge (\exists y. Q(y))$
 $\longrightarrow ((\forall x. P(x) \longrightarrow R(x)) \wedge (\forall y. Q(y) \longrightarrow S(y)) \longleftrightarrow$
 $(\forall x y. P(x) \wedge Q(y) \longrightarrow R(x) \wedge S(y))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \textbf{context } 1 \rangle$)

$\neg\neg$ 30

lemma

$\langle (\forall x. (P(x) \vee Q(x)) \longrightarrow \neg R(x)) \wedge$
 $(\forall x. (Q(x) \longrightarrow \neg S(x)) \longrightarrow P(x) \wedge R(x))$
 $\longrightarrow (\forall x. \neg\neg S(x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \textbf{context } 1 \rangle$)

31

lemma

$\langle \neg (\exists x. P(x) \wedge (Q(x) \vee R(x))) \wedge$
 $(\exists x. L(x) \wedge P(x)) \wedge$
 $(\forall x. \neg R(x) \longrightarrow M(x))$
 $\longrightarrow (\exists x. L(x) \wedge M(x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \textbf{context } 1 \rangle$)

32

lemma

$\langle (\forall x. P(x) \wedge (Q(x) \vee R(x)) \longrightarrow S(x)) \wedge$
 $(\forall x. S(x) \wedge R(x) \longrightarrow L(x)) \wedge$
 $(\forall x. M(x) \longrightarrow R(x))$
 $\longrightarrow (\forall x. P(x) \wedge M(x) \longrightarrow L(x)) \rangle$

by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

$\neg\neg 33$

lemma

$\langle (\forall x. \neg\neg (P(a) \wedge (P(x) \longrightarrow P(b)) \longrightarrow P(c))) \longleftrightarrow$
 $(\forall x. \neg\neg ((\neg P(a) \vee P(x) \vee P(c)) \wedge (\neg P(a) \vee \neg P(b) \vee P(c)))) \rangle$
apply (*tactic* $\langle \text{IntPr.best-tac } \mathbf{context} \ 1 \rangle$)
done

36

lemma

$\langle (\forall x. \exists y. J(x,y)) \wedge$
 $(\forall x. \exists y. G(x,y)) \wedge$
 $(\forall x y. J(x,y) \vee G(x,y) \longrightarrow (\forall z. J(y,z) \vee G(y,z) \longrightarrow H(x,z)))$
 $\longrightarrow (\forall x. \exists y. H(x,y)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

37

lemma

$\langle (\forall z. \exists w. \forall x. \exists y.$
 $\neg\neg (P(x,z) \longrightarrow P(y,w)) \wedge P(y,z) \wedge (P(y,w) \longrightarrow (\exists u. Q(u,w)))) \wedge$
 $(\forall x z. \neg P(x,z) \longrightarrow (\exists y. Q(y,z))) \wedge$
 $(\neg\neg (\exists x y. Q(x,y)) \longrightarrow (\forall x. R(x,x)))$
 $\longrightarrow \neg\neg (\forall x. \exists y. R(x,y)) \rangle$
oops — NOT PROVED

39

lemma $\langle \neg (\exists x. \forall y. F(y,x) \longleftrightarrow \neg F(y,y)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

40. AMENDED

lemma

$\langle (\exists y. \forall x. F(x,y) \longleftrightarrow F(x,x)) \longrightarrow$
 $\neg (\forall x. \exists y. \forall z. F(z,y) \longleftrightarrow \neg F(z,x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

44

lemma

$\langle (\forall x. f(x) \longrightarrow$
 $(\exists y. g(y) \wedge h(x,y) \wedge (\exists y. g(y) \wedge \neg h(x,y)))) \wedge$
 $(\exists x. j(x) \wedge (\forall y. g(y) \longrightarrow h(x,y)))$
 $\longrightarrow (\exists x. j(x) \wedge \neg f(x)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

48

lemma $\langle (a = b \vee c = d) \wedge (a = c \vee b = d) \longrightarrow a = d \vee b = c \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac } \mathbf{context} \ 1 \rangle$)

51

lemma

$\langle (\exists z w. \forall x y. P(x,y) \longleftrightarrow (x = z \wedge y = w)) \longrightarrow$
 $(\exists z. \forall x. \exists w. (\forall y. P(x,y) \longleftrightarrow y = w) \longleftrightarrow x = z) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

52

Almost the same as 51.

lemma

$\langle (\exists z w. \forall x y. P(x,y) \longleftrightarrow (x = z \wedge y = w)) \longrightarrow$
 $(\exists w. \forall y. \exists z. (\forall x. P(x,y) \longleftrightarrow x = z) \longleftrightarrow y = w) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

56

lemma $\langle (\forall x. (\exists y. P(y) \wedge x = f(y)) \longrightarrow P(x)) \longleftrightarrow (\forall x. P(x) \longrightarrow P(f(x))) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

57

lemma

$\langle P(f(a,b), f(b,c)) \wedge P(f(b,c), f(a,c)) \wedge$
 $(\forall x y z. P(x,y) \wedge P(y,z) \longrightarrow P(x,z)) \longrightarrow P(f(a,b), f(a,c)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

60

lemma $\langle \forall x. P(x, f(x)) \longleftrightarrow (\exists y. (\forall z. P(z, y) \longrightarrow P(z, f(x))) \wedge P(x, y)) \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

end

9 First-Order Logic: propositional examples (intuitionistic version)

theory *Propositional-Int*

imports *IFOL*

begin

commutative laws of \wedge and \vee

lemma $\langle P \wedge Q \longrightarrow Q \wedge P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

lemma $\langle P \vee Q \longrightarrow Q \vee P \rangle$
by (*tactic* $\langle \text{IntPr.fast-tac context } 1 \rangle$)

associative laws of \wedge and \vee

lemma $\langle (P \wedge Q) \wedge R \longrightarrow P \wedge (Q \wedge R) \rangle$

by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (P \vee Q) \vee R \longrightarrow P \vee (Q \vee R) \rangle$
by (*tactic IntPr.fast-tac context 1*)

distributive laws of \wedge and \vee

lemma $\langle (P \wedge Q) \vee R \longrightarrow (P \vee R) \wedge (Q \vee R) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (P \vee R) \wedge (Q \vee R) \longrightarrow (P \wedge Q) \vee R \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (P \vee Q) \wedge R \longrightarrow (P \wedge R) \vee (Q \wedge R) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (P \wedge R) \vee (Q \wedge R) \longrightarrow (P \vee Q) \wedge R \rangle$
by (*tactic IntPr.fast-tac context 1*)

Laws involving implication

lemma $\langle (P \longrightarrow R) \wedge (Q \longrightarrow R) \longleftrightarrow (P \vee Q \longrightarrow R) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (P \wedge Q \longrightarrow R) \longleftrightarrow (P \longrightarrow (Q \longrightarrow R)) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle ((P \longrightarrow R) \longrightarrow R) \longrightarrow ((Q \longrightarrow R) \longrightarrow R) \longrightarrow (P \wedge Q \longrightarrow R) \longrightarrow R \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle \neg (P \longrightarrow R) \longrightarrow \neg (Q \longrightarrow R) \longrightarrow \neg (P \wedge Q \longrightarrow R) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (P \longrightarrow Q \wedge R) \longleftrightarrow (P \longrightarrow Q) \wedge (P \longrightarrow R) \rangle$
by (*tactic IntPr.fast-tac context 1*)

Propositions-as-types

lemma $\langle P \longrightarrow (Q \longrightarrow P) \rangle$
by (*tactic IntPr.fast-tac context 1*)

— The combinator S

lemma $\langle (P \longrightarrow Q \longrightarrow R) \longrightarrow (P \longrightarrow Q) \longrightarrow (P \longrightarrow R) \rangle$
by (*tactic IntPr.fast-tac context 1*)

— Converse is classical

lemma $\langle (P \longrightarrow Q) \vee (P \longrightarrow R) \longrightarrow (P \longrightarrow Q \vee R) \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle (P \longrightarrow Q) \longrightarrow (\neg Q \longrightarrow \neg P) \rangle$
by (*tactic IntPr.fast-tac context 1*)

Schwichtenberg's examples (via T. Nipkow)

lemma *stab-imp*: $\langle (((Q \rightarrow R) \rightarrow R) \rightarrow Q) \rightarrow (((P \rightarrow Q) \rightarrow R) \rightarrow R) \rightarrow P \rightarrow Q \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma *stab-to-peirce*:
 $\langle (((P \rightarrow R) \rightarrow R) \rightarrow P) \rightarrow (((Q \rightarrow R) \rightarrow R) \rightarrow Q) \rightarrow ((P \rightarrow Q) \rightarrow P) \rightarrow P \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma *peirce-imp1*:
 $\langle (((Q \rightarrow R) \rightarrow Q) \rightarrow Q) \rightarrow (((P \rightarrow Q) \rightarrow R) \rightarrow P \rightarrow Q) \rightarrow P \rightarrow Q \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma *peirce-imp2*: $\langle (((P \rightarrow R) \rightarrow P) \rightarrow P) \rightarrow ((P \rightarrow Q \rightarrow R) \rightarrow P) \rightarrow P \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma *mints*: $\langle (((P \rightarrow Q) \rightarrow P) \rightarrow P) \rightarrow Q \rightarrow Q \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma *mints-solovev*: $\langle (P \rightarrow (Q \rightarrow R) \rightarrow Q) \rightarrow ((P \rightarrow Q) \rightarrow R) \rightarrow R \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma *tatsuta*:
 $\langle (((P7 \rightarrow P1) \rightarrow P10) \rightarrow P4 \rightarrow P5) \rightarrow (((P8 \rightarrow P2) \rightarrow P9) \rightarrow P3 \rightarrow P10) \rightarrow (P1 \rightarrow P8) \rightarrow P6 \rightarrow P7 \rightarrow (((P3 \rightarrow P2) \rightarrow P9) \rightarrow P4) \rightarrow (P1 \rightarrow P3) \rightarrow (((P6 \rightarrow P1) \rightarrow P2) \rightarrow P9) \rightarrow P5 \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma *tatsuta1*:
 $\langle (((P8 \rightarrow P2) \rightarrow P9) \rightarrow P3 \rightarrow P10) \rightarrow (((P3 \rightarrow P2) \rightarrow P9) \rightarrow P4) \rightarrow (((P6 \rightarrow P1) \rightarrow P2) \rightarrow P9) \rightarrow (((P7 \rightarrow P1) \rightarrow P10) \rightarrow P4 \rightarrow P5) \rightarrow (P1 \rightarrow P3) \rightarrow (P1 \rightarrow P8) \rightarrow P6 \rightarrow P7 \rightarrow P5 \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

end

10 First-Order Logic: quantifier examples (intuitionistic version)

theory *Quantifiers-Int*
imports *IFOL*

begin

lemma $\langle (\forall x y. P(x,y)) \longrightarrow (\forall y x. P(x,y)) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma $\langle (\exists x y. P(x,y)) \longrightarrow (\exists y x. P(x,y)) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

— Converse is false

lemma $\langle (\forall x. P(x)) \vee (\forall x. Q(x)) \longrightarrow (\forall x. P(x) \vee Q(x)) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma $\langle (\forall x. P \longrightarrow Q(x)) \longleftrightarrow (P \longrightarrow (\forall x. Q(x))) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma $\langle (\forall x. P(x) \longrightarrow Q) \longleftrightarrow ((\exists x. P(x)) \longrightarrow Q) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

Some harder ones

lemma $\langle (\exists x. P(x) \vee Q(x)) \longleftrightarrow (\exists x. P(x)) \vee (\exists x. Q(x)) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

— Converse is false

lemma $\langle (\exists x. P(x) \wedge Q(x)) \longrightarrow (\exists x. P(x)) \wedge (\exists x. Q(x)) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

Basic test of quantifier reasoning

lemma $\langle (\exists y. \forall x. Q(x,y)) \longrightarrow (\forall x. \exists y. Q(x,y)) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

lemma $\langle (\forall x. Q(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by (*tactic* *IntPr.fast-tac* **context** 1)

The following should fail, as they are false!

lemma $\langle (\forall x. \exists y. Q(x,y)) \longrightarrow (\exists y. \forall x. Q(x,y)) \rangle$
apply (*tactic* *IntPr.fast-tac* **context** 1)?
oops

lemma $\langle (\exists x. Q(x)) \longrightarrow (\forall x. Q(x)) \rangle$
apply (*tactic* *IntPr.fast-tac* **context** 1)?
oops

schematic-goal $\langle P(?a) \longrightarrow (\forall x. P(x)) \rangle$
apply (*tactic* *IntPr.fast-tac* **context** 1)?
oops

schematic-goal $\langle (P(?a) \longrightarrow (\forall x. Q(x))) \longrightarrow (\forall x. P(x) \longrightarrow Q(x)) \rangle$

```

apply (tactic IntPr.fast-tac context 1)?
oops

```

Back to things that are provable ...

```

lemma  $\langle (\forall x. P(x) \longrightarrow Q(x)) \wedge (\exists x. P(x)) \longrightarrow (\exists x. Q(x)) \rangle$ 
by (tactic IntPr.fast-tac context 1)

```

— An example of why exI should be delayed as long as possible

```

lemma  $\langle (P \longrightarrow (\exists x. Q(x))) \wedge P \longrightarrow (\exists x. Q(x)) \rangle$ 
by (tactic IntPr.fast-tac context 1)

```

```

schematic-goal  $\langle (\forall x. P(x) \longrightarrow Q(f(x))) \wedge (\forall x. Q(x) \longrightarrow R(g(x))) \wedge P(d) \longrightarrow R(?a) \rangle$ 
by (tactic IntPr.fast-tac context 1)

```

```

lemma  $\langle (\forall x. Q(x)) \longrightarrow (\exists x. Q(x)) \rangle$ 
by (tactic IntPr.fast-tac context 1)

```

Some slow ones

```

lemma  $\langle (\forall x y. P(x) \longrightarrow Q(y)) \longleftrightarrow ((\exists x. P(x)) \longrightarrow (\forall y. Q(y))) \rangle$ 
by (tactic IntPr.fast-tac context 1)

```

```

lemma  $\langle (\exists x y. P(x) \wedge Q(x,y)) \longleftrightarrow (\exists x. P(x) \wedge (\exists y. Q(x,y))) \rangle$ 
by (tactic IntPr.fast-tac context 1)

```

```

lemma  $\langle (\exists y. \forall x. P(x) \longrightarrow Q(x,y)) \longrightarrow (\forall x. P(x) \longrightarrow (\exists y. Q(x,y))) \rangle$ 
by (tactic IntPr.fast-tac context 1)

```

end

11 Classical Predicate Calculus Problems

```

theory Classical
imports FOL
begin

```

```

lemma  $\langle (P \longrightarrow Q \vee R) \longrightarrow (P \longrightarrow Q) \vee (P \longrightarrow R) \rangle$ 
by blast

```

11.0.1 If and only if

```

lemma  $\langle (P \longleftrightarrow Q) \longleftrightarrow (Q \longleftrightarrow P) \rangle$ 
by blast

```

```

lemma  $\langle \neg (P \longleftrightarrow \neg P) \rangle$ 
by blast

```

11.1 Pelletier's examples

Sample problems from

- F. J. Pelletier, Seventy-Five Problems for Testing Automatic Theorem Provers, J. Automated Reasoning 2 (1986), 191-216. Errata, JAR 4 (1988), 236-236.

The hardest problems – judging by experience with several theorem provers, including matrix ones – are 34 and 43.

1

lemma $\langle (P \longrightarrow Q) \longleftrightarrow (\neg Q \longrightarrow \neg P) \rangle$
by *blast*

2

lemma $\langle \neg \neg P \longleftrightarrow P \rangle$
by *blast*

3

lemma $\langle \neg (P \longrightarrow Q) \longrightarrow (Q \longrightarrow P) \rangle$
by *blast*

4

lemma $\langle (\neg P \longrightarrow Q) \longleftrightarrow (\neg Q \longrightarrow P) \rangle$
by *blast*

5

lemma $\langle ((P \vee Q) \longrightarrow (P \vee R)) \longrightarrow (P \vee (Q \longrightarrow R)) \rangle$
by *blast*

6

lemma $\langle P \vee \neg P \rangle$
by *blast*

7

lemma $\langle P \vee \neg \neg \neg P \rangle$
by *blast*

8. Peirce's law

lemma $\langle ((P \longrightarrow Q) \longrightarrow P) \longrightarrow P \rangle$
by *blast*

9

lemma $\langle ((P \vee Q) \wedge (\neg P \vee Q) \wedge (P \vee \neg Q)) \longrightarrow \neg (\neg P \vee \neg Q) \rangle$
by *blast*

10

lemma $\langle (Q \longrightarrow R) \wedge (R \longrightarrow P \wedge Q) \wedge (P \longrightarrow Q \vee R) \longrightarrow (P \longleftrightarrow Q) \rangle$
by *blast*

11. Proved in each direction (incorrectly, says Pelletier!!)

lemma $\langle P \longleftrightarrow P \rangle$
by *blast*

12. "Dijkstra's law"

lemma $\langle (P \longleftrightarrow Q) \longleftrightarrow R \longleftrightarrow (P \longleftrightarrow (Q \longleftrightarrow R)) \rangle$
by *blast*

13. Distributive law

lemma $\langle P \vee (Q \wedge R) \longleftrightarrow (P \vee Q) \wedge (P \vee R) \rangle$
by *blast*

14

lemma $\langle (P \longleftrightarrow Q) \longleftrightarrow ((Q \vee \neg P) \wedge (\neg Q \vee P)) \rangle$
by *blast*

15

lemma $\langle (P \longrightarrow Q) \longleftrightarrow (\neg P \vee Q) \rangle$
by *blast*

16

lemma $\langle (P \longrightarrow Q) \vee (Q \longrightarrow P) \rangle$
by *blast*

17

lemma $\langle ((P \wedge (Q \longrightarrow R)) \longrightarrow S) \longleftrightarrow ((\neg P \vee Q \vee S) \wedge (\neg P \vee \neg R \vee S)) \rangle$
by *blast*

11.2 Classical Logic: examples with quantifiers

lemma $\langle (\forall x. P(x) \wedge Q(x)) \longleftrightarrow (\forall x. P(x)) \wedge (\forall x. Q(x)) \rangle$
by *blast*

lemma $\langle (\exists x. P \longrightarrow Q(x)) \longleftrightarrow (P \longrightarrow (\exists x. Q(x))) \rangle$
by *blast*

lemma $\langle (\exists x. P(x) \longrightarrow Q) \longleftrightarrow (\forall x. P(x)) \longrightarrow Q \rangle$
by *blast*

lemma $\langle (\forall x. P(x)) \vee Q \longleftrightarrow (\forall x. P(x) \vee Q) \rangle$
by *blast*

Discussed in Avron, Gentzen-Type Systems, Resolution and Tableaux, JAR
10 (265-281), 1993. Proof is trivial!

lemma $\langle \neg ((\exists x. \neg P(x)) \wedge ((\exists x. P(x)) \vee (\exists x. P(x) \wedge Q(x))) \wedge \neg (\exists x. P(x))) \rangle$
by *blast*

11.3 Problems requiring quantifier duplication

Theorem B of Peter Andrews, Theorem Proving via General Matings, JACM 28 (1981).

lemma $\langle (\exists x. \forall y. P(x) \longleftrightarrow P(y)) \longrightarrow ((\exists x. P(x)) \longleftrightarrow (\forall y. P(y))) \rangle$
by *blast*

Needs multiple instantiation of ALL.

lemma $\langle (\forall x. P(x) \longrightarrow P(f(x))) \wedge P(d) \longrightarrow P(f(f(f(d)))) \rangle$
by *blast*

Needs double instantiation of the quantifier

lemma $\langle \exists x. P(x) \longrightarrow P(a) \wedge P(b) \rangle$
by *blast*

lemma $\langle \exists z. P(z) \longrightarrow (\forall x. P(x)) \rangle$
by *blast*

lemma $\langle \exists x. (\exists y. P(y)) \longrightarrow P(x) \rangle$
by *blast*

V. Lifschitz, What Is the Inverse Method?, JAR 5 (1989), 1–23. NOT PROVED.

lemma
 $\langle \exists x x'. \forall y. \exists z z'.$
 $(\neg P(y, y) \vee P(x, x) \vee \neg S(z, x)) \wedge$
 $(S(x, y) \vee \neg S(y, z) \vee Q(z', z')) \wedge$
 $(Q(x', y) \vee \neg Q(y, z') \vee S(x', x')) \rangle$
oops

11.4 Hard examples with quantifiers

18

lemma $\langle \exists y. \forall x. P(y) \longrightarrow P(x) \rangle$
by *blast*

19

lemma $\langle \exists x. \forall y z. (P(y) \longrightarrow Q(z)) \longrightarrow (P(x) \longrightarrow Q(x)) \rangle$
by *blast*

20

lemma $\langle (\forall x y. \exists z. \forall w. (P(x) \wedge Q(y) \longrightarrow R(z) \wedge S(w)))$
 $\longrightarrow (\exists x y. P(x) \wedge Q(y)) \longrightarrow (\exists z. R(z)) \rangle$

by *blast*

21

lemma $\langle (\exists x. P \longrightarrow Q(x)) \wedge (\exists x. Q(x) \longrightarrow P) \longrightarrow (\exists x. P \longleftrightarrow Q(x)) \rangle$
by *blast*

22

lemma $\langle (\forall x. P \longleftrightarrow Q(x)) \longrightarrow (P \longleftrightarrow (\forall x. Q(x))) \rangle$
by *blast*

23

lemma $\langle (\forall x. P \vee Q(x)) \longleftrightarrow (P \vee (\forall x. Q(x))) \rangle$
by *blast*

24

lemma
 $\langle \neg (\exists x. S(x) \wedge Q(x)) \wedge (\forall x. P(x) \longrightarrow Q(x) \vee R(x)) \wedge$
 $\neg (\exists x. P(x)) \longrightarrow (\exists x. Q(x)) \wedge (\forall x. Q(x) \vee R(x) \longrightarrow S(x))$
 $\longrightarrow (\exists x. P(x) \wedge R(x)) \rangle$
by *blast*

25

lemma
 $\langle (\exists x. P(x)) \wedge$
 $(\forall x. L(x) \longrightarrow \neg (M(x) \wedge R(x))) \wedge$
 $(\forall x. P(x) \longrightarrow (M(x) \wedge L(x))) \wedge$
 $((\forall x. P(x) \longrightarrow Q(x)) \vee (\exists x. P(x) \wedge R(x)))$
 $\longrightarrow (\exists x. Q(x) \wedge P(x)) \rangle$
by *blast*

26

lemma
 $\langle ((\exists x. p(x)) \longleftrightarrow (\exists x. q(x))) \wedge$
 $(\forall x. \forall y. p(x) \wedge q(y) \longrightarrow (r(x) \longleftrightarrow s(y)))$
 $\longrightarrow ((\forall x. p(x) \longrightarrow r(x)) \longleftrightarrow (\forall x. q(x) \longrightarrow s(x))) \rangle$
by *blast*

27

lemma
 $\langle (\exists x. P(x) \wedge \neg Q(x)) \wedge$
 $(\forall x. P(x) \longrightarrow R(x)) \wedge$
 $(\forall x. M(x) \wedge L(x) \longrightarrow P(x)) \wedge$
 $((\exists x. R(x) \wedge \neg Q(x)) \longrightarrow (\forall x. L(x) \longrightarrow \neg R(x)))$
 $\longrightarrow (\forall x. M(x) \longrightarrow \neg L(x)) \rangle$
by *blast*

28. AMENDED

lemma

$$\begin{aligned} &\langle (\forall x. P(x) \longrightarrow (\forall x. Q(x))) \wedge \\ &\quad ((\forall x. Q(x) \vee R(x)) \longrightarrow (\exists x. Q(x) \wedge S(x))) \wedge \\ &\quad ((\exists x. S(x)) \longrightarrow (\forall x. L(x) \longrightarrow M(x))) \\ &\longrightarrow (\forall x. P(x) \wedge L(x) \longrightarrow M(x)) \rangle \\ &\text{by } \textit{blast} \end{aligned}$$

29. Essentially the same as Principia Mathematica *11.71

lemma

$$\begin{aligned} &\langle (\exists x. P(x)) \wedge (\exists y. Q(y)) \\ &\longrightarrow ((\forall x. P(x) \longrightarrow R(x)) \wedge (\forall y. Q(y) \longrightarrow S(y)) \longleftrightarrow \\ &\quad (\forall x y. P(x) \wedge Q(y) \longrightarrow R(x) \wedge S(y))) \rangle \\ &\text{by } \textit{blast} \end{aligned}$$

30

lemma

$$\begin{aligned} &\langle (\forall x. P(x) \vee Q(x) \longrightarrow \neg R(x)) \wedge \\ &\quad (\forall x. (Q(x) \longrightarrow \neg S(x)) \longrightarrow P(x) \wedge R(x)) \\ &\longrightarrow (\forall x. S(x)) \rangle \\ &\text{by } \textit{blast} \end{aligned}$$

31

lemma

$$\begin{aligned} &\langle \neg (\exists x. P(x) \wedge (Q(x) \vee R(x))) \wedge \\ &\quad (\exists x. L(x) \wedge P(x)) \wedge \\ &\quad (\forall x. \neg R(x) \longrightarrow M(x)) \\ &\longrightarrow (\exists x. L(x) \wedge M(x)) \rangle \\ &\text{by } \textit{blast} \end{aligned}$$

32

lemma

$$\begin{aligned} &\langle (\forall x. P(x) \wedge (Q(x) \vee R(x)) \longrightarrow S(x)) \wedge \\ &\quad (\forall x. S(x) \wedge R(x) \longrightarrow L(x)) \wedge \\ &\quad (\forall x. M(x) \longrightarrow R(x)) \\ &\longrightarrow (\forall x. P(x) \wedge M(x) \longrightarrow L(x)) \rangle \\ &\text{by } \textit{blast} \end{aligned}$$

33

lemma

$$\begin{aligned} &\langle (\forall x. P(a) \wedge (P(x) \longrightarrow P(b)) \longrightarrow P(c)) \longleftrightarrow \\ &\quad (\forall x. (\neg P(a) \vee P(x) \vee P(c)) \wedge (\neg P(a) \vee \neg P(b) \vee P(c))) \rangle \\ &\text{by } \textit{blast} \end{aligned}$$

34. AMENDED (TWICE!!). Andrews's challenge.

lemma

$$\begin{aligned} &\langle ((\exists x. \forall y. p(x) \longleftrightarrow p(y)) \longleftrightarrow ((\exists x. q(x)) \longleftrightarrow (\forall y. p(y)))) \longleftrightarrow \\ &\quad ((\exists x. \forall y. q(x) \longleftrightarrow q(y)) \longleftrightarrow ((\exists x. p(x)) \longleftrightarrow (\forall y. q(y)))) \rangle \\ &\text{by } \textit{blast} \end{aligned}$$

35

lemma $\langle \exists x y. P(x,y) \longrightarrow (\forall u v. P(u,v)) \rangle$
by *blast*

36

lemma
 $\langle (\forall x. \exists y. J(x,y)) \wedge$
 $(\forall x. \exists y. G(x,y)) \wedge$
 $(\forall x y. J(x,y) \vee G(x,y) \longrightarrow (\forall z. J(y,z) \vee G(y,z) \longrightarrow H(x,z)))$
 $\longrightarrow (\forall x. \exists y. H(x,y)) \rangle$
by *blast*

37

lemma
 $\langle (\forall z. \exists w. \forall x. \exists y.$
 $(P(x,z) \longrightarrow P(y,w)) \wedge P(y,z) \wedge (P(y,w) \longrightarrow (\exists u. Q(u,w)))) \wedge$
 $(\forall x z. \neg P(x,z) \longrightarrow (\exists y. Q(y,z))) \wedge$
 $((\exists x y. Q(x,y)) \longrightarrow (\forall x. R(x,x)))$
 $\longrightarrow (\forall x. \exists y. R(x,y)) \rangle$
by *blast*

38

lemma
 $\langle (\forall x. p(a) \wedge (p(x) \longrightarrow (\exists y. p(y) \wedge r(x,y))) \longrightarrow$
 $(\exists z. \exists w. p(z) \wedge r(x,w) \wedge r(w,z))) \longleftrightarrow$
 $(\forall x. (\neg p(a) \vee p(x) \vee (\exists z. \exists w. p(z) \wedge r(x,w) \wedge r(w,z))) \wedge$
 $(\neg p(a) \vee \neg (\exists y. p(y) \wedge r(x,y)) \vee$
 $(\exists z. \exists w. p(z) \wedge r(x,w) \wedge r(w,z)))) \rangle$
by *blast*

39

lemma $\langle \neg (\exists x. \forall y. F(y,x) \longleftrightarrow \neg F(y,y)) \rangle$
by *blast*

40. AMENDED

lemma
 $\langle (\exists y. \forall x. F(x,y) \longleftrightarrow F(x,x)) \longrightarrow$
 $\neg (\forall x. \exists y. \forall z. F(z,y) \longleftrightarrow \neg F(z,x)) \rangle$
by *blast*

41

lemma
 $\langle (\forall z. \exists y. \forall x. f(x,y) \longleftrightarrow f(x,z) \wedge \neg f(x,x))$
 $\longrightarrow \neg (\exists z. \forall x. f(x,z)) \rangle$
by *blast*

42

lemma $\langle \neg (\exists y. \forall x. p(x,y) \longleftrightarrow \neg (\exists z. p(x,z) \wedge p(z,x))) \rangle$
by *blast*

43

lemma
 $\langle (\forall x. \forall y. q(x,y) \longleftrightarrow (\forall z. p(z,x) \longleftrightarrow p(z,y)))$
 $\longrightarrow (\forall x. \forall y. q(x,y) \longleftrightarrow q(y,x)) \rangle$
by *blast*

Other proofs: Can use *auto*, which cheats by using rewriting! *Deepen-tac* alone requires 253 secs. Or *by (mini-tac 1 THEN Deepen-tac 5 1)*.

44

lemma
 $\langle (\forall x. f(x) \longrightarrow (\exists y. g(y) \wedge h(x,y) \wedge (\exists y. g(y) \wedge \neg h(x,y)))) \wedge$
 $(\exists x. j(x) \wedge (\forall y. g(y) \longrightarrow h(x,y)))$
 $\longrightarrow (\exists x. j(x) \wedge \neg f(x)) \rangle$
by *blast*

45

lemma
 $\langle (\forall x. f(x) \wedge (\forall y. g(y) \wedge h(x,y) \longrightarrow j(x,y))$
 $\longrightarrow (\forall y. g(y) \wedge h(x,y) \longrightarrow k(y))) \wedge$
 $\neg (\exists y. l(y) \wedge k(y)) \wedge$
 $(\exists x. f(x) \wedge (\forall y. h(x,y) \longrightarrow l(y)) \wedge (\forall y. g(y) \wedge h(x,y) \longrightarrow j(x,y)))$
 $\longrightarrow (\exists x. f(x) \wedge \neg (\exists y. g(y) \wedge h(x,y))) \rangle$
by *blast*

46

lemma
 $\langle (\forall x. f(x) \wedge (\forall y. f(y) \wedge h(y,x) \longrightarrow g(y)) \longrightarrow g(x)) \wedge$
 $((\exists x. f(x) \wedge \neg g(x)) \longrightarrow$
 $(\exists x. f(x) \wedge \neg g(x) \wedge (\forall y. f(y) \wedge \neg g(y) \longrightarrow j(x,y)))) \wedge$
 $(\forall x y. f(x) \wedge f(y) \wedge h(x,y) \longrightarrow \neg j(y,x))$
 $\longrightarrow (\forall x. f(x) \longrightarrow g(x)) \rangle$
by *blast*

11.5 Problems (mainly) involving equality or functions

48

lemma $\langle (a = b \vee c = d) \wedge (a = c \vee b = d) \longrightarrow a = d \vee b = c \rangle$
by *blast*

49. NOT PROVED AUTOMATICALLY. Hard because it involves substitution for Vars; the type constraint ensures that x,y,z have the same type as a,b,u.

lemma

$\langle (\exists x y. : 'a. \forall z. z = x \vee z = y) \wedge P(a) \wedge P(b) \wedge a \neq b \longrightarrow (\forall u. : 'a. P(u)) \rangle$
apply *safe*
apply (*rule-tac* $x = \langle a \rangle$ **in** *allE, assumption*)
apply (*rule-tac* $x = \langle b \rangle$ **in** *allE, assumption*)
apply *fast* — blast's treatment of equality can't do it
done

50. (What has this to do with equality?)

lemma $\langle (\forall x. P(a, x) \vee (\forall y. P(x, y))) \longrightarrow (\exists x. \forall y. P(x, y)) \rangle$
by *blast*

51

lemma
 $\langle (\exists z w. \forall x y. P(x, y) \longleftrightarrow (x = z \wedge y = w)) \longrightarrow$
 $(\exists z. \forall x y. \exists w. (\forall y. P(x, y) \longleftrightarrow y = w) \longleftrightarrow x = z) \rangle$
by *blast*

52

Almost the same as 51.

lemma
 $\langle (\exists z w. \forall x y. P(x, y) \longleftrightarrow (x = z \wedge y = w)) \longrightarrow$
 $(\exists w. \forall y. \exists z. (\forall x. P(x, y) \longleftrightarrow x = z) \longleftrightarrow y = w) \rangle$
by *blast*

55

Non-equational version, from Manthey and Bry, CADE-9 (Springer, 1988).
fast DISCOVERS who killed Agatha.

schematic-goal
 $\langle \text{lives}(\text{agatha}) \wedge \text{lives}(\text{butler}) \wedge \text{lives}(\text{charles}) \wedge$
 $(\text{killed}(\text{agatha}, \text{agatha}) \vee \text{killed}(\text{butler}, \text{agatha}) \vee \text{killed}(\text{charles}, \text{agatha})) \wedge$
 $(\forall x y. \text{killed}(x, y) \longrightarrow \text{hates}(x, y) \wedge \neg \text{richer}(x, y)) \wedge$
 $(\forall x. \text{hates}(\text{agatha}, x) \longrightarrow \neg \text{hates}(\text{charles}, x)) \wedge$
 $(\text{hates}(\text{agatha}, \text{agatha}) \wedge \text{hates}(\text{agatha}, \text{charles})) \wedge$
 $(\forall x. \text{lives}(x) \wedge \neg \text{richer}(x, \text{agatha}) \longrightarrow \text{hates}(\text{butler}, x)) \wedge$
 $(\forall x. \text{hates}(\text{agatha}, x) \longrightarrow \text{hates}(\text{butler}, x)) \wedge$
 $(\forall x. \neg \text{hates}(x, \text{agatha}) \vee \neg \text{hates}(x, \text{butler}) \vee \neg \text{hates}(x, \text{charles})) \longrightarrow$
 $\text{killed}(\text{?who}, \text{agatha}) \rangle$
by *fast* — MUCH faster than blast

56

lemma $\langle (\forall x. (\exists y. P(y) \wedge x = f(y)) \longrightarrow P(x)) \longleftrightarrow (\forall x. P(x) \longrightarrow P(f(x))) \rangle$
by *blast*

57

lemma
 $\langle P(f(a, b), f(b, c)) \wedge P(f(b, c), f(a, c)) \wedge$

$(\forall x y z. P(x,y) \wedge P(y,z) \longrightarrow P(x,z)) \longrightarrow P(f(a,b), f(a,c))\rangle$
by *blast*

58 NOT PROVED AUTOMATICALLY

lemma $\langle (\forall x y. f(x) = g(y)) \longrightarrow (\forall x y. f(f(x)) = f(g(y))) \rangle$
by (*slow elim: subst-context*)

59

lemma $\langle (\forall x. P(x) \longleftrightarrow \neg P(f(x))) \longrightarrow (\exists x. P(x) \wedge \neg P(f(x))) \rangle$
by *blast*

60

lemma $\langle \forall x. P(x, f(x)) \longleftrightarrow (\exists y. (\forall z. P(z, y) \longrightarrow P(z, f(x))) \wedge P(x, y)) \rangle$
by *blast*

62 as corrected in JAR 18 (1997), page 135

lemma
 $\langle (\forall x. p(a) \wedge (p(x) \longrightarrow p(f(x))) \longrightarrow p(f(f(x)))) \longleftrightarrow$
 $(\forall x. (\neg p(a) \vee p(x) \vee p(f(f(x)))) \wedge$
 $(\neg p(a) \vee \neg p(f(x)) \vee p(f(f(x)))) \rangle$
by *blast*

From Davis, Obvious Logical Inferences, IJCAI-81, 530-531 fast indeed copes!

lemma
 $\langle (\forall x. F(x) \wedge \neg G(x) \longrightarrow (\exists y. H(x,y) \wedge J(y))) \wedge$
 $(\exists x. K(x) \wedge F(x) \wedge (\forall y. H(x,y) \longrightarrow K(y))) \wedge$
 $(\forall x. K(x) \longrightarrow \neg G(x)) \longrightarrow (\exists x. K(x) \wedge J(x)) \rangle$
by *fast*

From Rudnicki, Obvious Inferences, JAR 3 (1987), 383-393. It does seem obvious!

lemma
 $\langle (\forall x. F(x) \wedge \neg G(x) \longrightarrow (\exists y. H(x,y) \wedge J(y))) \wedge$
 $(\exists x. K(x) \wedge F(x) \wedge (\forall y. H(x,y) \longrightarrow K(y))) \wedge$
 $(\forall x. K(x) \longrightarrow \neg G(x)) \longrightarrow (\exists x. K(x) \longrightarrow \neg G(x)) \rangle$
by *fast*

Halting problem: Formulation of Li Dafa (AAR Newsletter 27, Oct 1994.)
author U. Egly.

lemma
 $\langle ((\exists x. A(x) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(x,y,z)))) \longrightarrow$
 $(\exists w. C(w) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(w,y,z)))) \wedge$
 $(\forall w. C(w) \wedge (\forall u. C(u) \longrightarrow (\forall v. D(w,u,v))) \longrightarrow$
 $(\forall y z.$
 $(C(y) \wedge P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,g)) \wedge$
 $(C(y) \wedge \neg P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,b)))) \rangle$

\wedge
 $(\forall w. C(w) \wedge$
 $(\forall y z.$
 $(C(y) \wedge P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,g)) \wedge$
 $(C(y) \wedge \neg P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,b))) \longrightarrow$
 $(\exists v. C(v) \wedge$
 $(\forall y. ((C(y) \wedge Q(w,y,y)) \wedge OO(w,g) \longrightarrow \neg P(v,y)) \wedge$
 $((C(y) \wedge Q(w,y,y)) \wedge OO(w,b) \longrightarrow P(v,y) \wedge OO(v,b))))$
 $\longrightarrow \neg (\exists x. A(x) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(x,y,z))))\rangle$
by (*blast 12*)
 — Needed because the search for depths below 12 is very slow.

Halting problem II: credited to M. Bruschi by Li Dafa in JAR 18(1), p. 105.

lemma

$\langle ((\exists x. A(x) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(x,y,z)))) \longrightarrow$
 $(\exists w. C(w) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(w,y,z))))$
 \wedge
 $(\forall w. C(w) \wedge (\forall u. C(u) \longrightarrow (\forall v. D(w,u,v))) \longrightarrow$
 $(\forall y z.$
 $(C(y) \wedge P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,g)) \wedge$
 $(C(y) \wedge \neg P(y,z) \longrightarrow Q(w,y,z) \wedge OO(w,b))))$
 \wedge
 $((\exists w. C(w) \wedge (\forall y. (C(y) \wedge P(y,y) \longrightarrow Q(w,y,y) \wedge OO(w,g)) \wedge$
 $(C(y) \wedge \neg P(y,y) \longrightarrow Q(w,y,y) \wedge OO(w,b))))$
 \longrightarrow
 $(\exists v. C(v) \wedge (\forall y. (C(y) \wedge P(y,y) \longrightarrow P(v,y) \wedge OO(v,g)) \wedge$
 $(C(y) \wedge \neg P(y,y) \longrightarrow P(v,y) \wedge OO(v,b))))$
 \longrightarrow
 $((\exists v. C(v) \wedge (\forall y. (C(y) \wedge P(y,y) \longrightarrow P(v,y) \wedge OO(v,g)) \wedge$
 $(C(y) \wedge \neg P(y,y) \longrightarrow P(v,y) \wedge OO(v,b))))$
 \longrightarrow
 $(\exists u. C(u) \wedge (\forall y. (C(y) \wedge P(y,y) \longrightarrow \neg P(u,y)) \wedge$
 $(C(y) \wedge \neg P(y,y) \longrightarrow P(u,y) \wedge OO(u,b))))$
 $\longrightarrow \neg (\exists x. A(x) \wedge (\forall y. C(y) \longrightarrow (\forall z. D(x,y,z))))\rangle$
by *blast*

Challenge found on info-hol.

lemma $\langle \forall x. \exists v w. \forall y z. P(x) \wedge Q(y) \longrightarrow (P(v) \vee R(w)) \wedge (R(z) \longrightarrow Q(v)) \rangle$
by *blast*

Attributed to Lewis Carroll by S. G. Pulman. The first or last assumption can be deleted.

lemma

$\langle (\forall x. \text{honest}(x) \wedge \text{industrious}(x) \longrightarrow \text{healthy}(x)) \wedge$
 $\neg (\exists x. \text{grocer}(x) \wedge \text{healthy}(x)) \wedge$
 $(\forall x. \text{industrious}(x) \wedge \text{grocer}(x) \longrightarrow \text{honest}(x)) \wedge$
 $(\forall x. \text{cyclist}(x) \longrightarrow \text{industrious}(x)) \wedge$
 $(\forall x. \neg \text{healthy}(x) \wedge \text{cyclist}(x) \longrightarrow \neg \text{honest}(x))$
 $\longrightarrow (\forall x. \text{grocer}(x) \longrightarrow \neg \text{cyclist}(x)) \rangle$

by *blast*

end

12 First-Order Logic: propositional examples (classical version)

theory *Propositional-Cla*
imports *FOL*
begin

commutative laws of \wedge and \vee

lemma $\langle P \wedge Q \longrightarrow Q \wedge P \rangle$
by (*tactic IntPr.fast-tac context 1*)

lemma $\langle P \vee Q \longrightarrow Q \vee P \rangle$
by *fast*

associative laws of \wedge and \vee

lemma $\langle (P \wedge Q) \wedge R \longrightarrow P \wedge (Q \wedge R) \rangle$
by *fast*

lemma $\langle (P \vee Q) \vee R \longrightarrow P \vee (Q \vee R) \rangle$
by *fast*

distributive laws of \wedge and \vee

lemma $\langle (P \wedge Q) \vee R \longrightarrow (P \vee R) \wedge (Q \vee R) \rangle$
by *fast*

lemma $\langle (P \vee R) \wedge (Q \vee R) \longrightarrow (P \wedge Q) \vee R \rangle$
by *fast*

lemma $\langle (P \vee Q) \wedge R \longrightarrow (P \wedge R) \vee (Q \wedge R) \rangle$
by *fast*

lemma $\langle (P \wedge R) \vee (Q \wedge R) \longrightarrow (P \vee Q) \wedge R \rangle$
by *fast*

Laws involving implication

lemma $\langle (P \longrightarrow R) \wedge (Q \longrightarrow R) \longleftrightarrow (P \vee Q \longrightarrow R) \rangle$
by *fast*

lemma $\langle (P \wedge Q \longrightarrow R) \longleftrightarrow (P \longrightarrow (Q \longrightarrow R)) \rangle$
by *fast*

lemma $\langle ((P \longrightarrow R) \longrightarrow R) \longrightarrow ((Q \longrightarrow R) \longrightarrow R) \longrightarrow (P \wedge Q \longrightarrow R) \longrightarrow R \rangle$
by *fast*

lemma $\langle \neg (P \longrightarrow R) \longrightarrow \neg (Q \longrightarrow R) \longrightarrow \neg (P \wedge Q \longrightarrow R) \rangle$
by *fast*

lemma $\langle (P \longrightarrow Q \wedge R) \longleftrightarrow (P \longrightarrow Q) \wedge (P \longrightarrow R) \rangle$
by *fast*

Propositions-as-types

lemma $\langle P \longrightarrow (Q \longrightarrow P) \rangle$
by *fast*

— The combinator S

lemma $\langle (P \longrightarrow Q \longrightarrow R) \longrightarrow (P \longrightarrow Q) \longrightarrow (P \longrightarrow R) \rangle$
by *fast*

— Converse is classical

lemma $\langle (P \longrightarrow Q) \vee (P \longrightarrow R) \longrightarrow (P \longrightarrow Q \vee R) \rangle$
by *fast*

lemma $\langle (P \longrightarrow Q) \longrightarrow (\neg Q \longrightarrow \neg P) \rangle$
by *fast*

Schwichtenberg’s examples (via T. Nipkow)

lemma *stab-imp*: $\langle (((Q \longrightarrow R) \longrightarrow R) \longrightarrow Q) \longrightarrow (((P \longrightarrow Q) \longrightarrow R) \longrightarrow R) \longrightarrow P \longrightarrow Q \rangle$
by *fast*

lemma *stab-to-peirce*:
 $\langle (((P \longrightarrow R) \longrightarrow R) \longrightarrow P) \longrightarrow (((Q \longrightarrow R) \longrightarrow R) \longrightarrow Q) \longrightarrow ((P \longrightarrow Q) \longrightarrow P) \longrightarrow P \rangle$
by *fast*

lemma *peirce-imp1*:
 $\langle (((Q \longrightarrow R) \longrightarrow Q) \longrightarrow Q) \longrightarrow (((P \longrightarrow Q) \longrightarrow R) \longrightarrow P \longrightarrow Q) \longrightarrow P \longrightarrow Q \rangle$
by *fast*

lemma *peirce-imp2*: $\langle (((P \longrightarrow R) \longrightarrow P) \longrightarrow P) \longrightarrow ((P \longrightarrow Q \longrightarrow R) \longrightarrow P) \longrightarrow P \rangle$
by *fast*

lemma *mits*: $\langle (((P \longrightarrow Q) \longrightarrow P) \longrightarrow P) \longrightarrow Q \longrightarrow Q \rangle$
by *fast*

lemma *mits-solovev*: $\langle (P \longrightarrow (Q \longrightarrow R) \longrightarrow Q) \longrightarrow ((P \longrightarrow Q) \longrightarrow R) \longrightarrow R \rangle$

by *fast*

lemma *tatsuta*:

$\langle (((P7 \rightarrow P1) \rightarrow P10) \rightarrow P4 \rightarrow P5)$
 $\rightarrow (((P8 \rightarrow P2) \rightarrow P9) \rightarrow P3 \rightarrow P10)$
 $\rightarrow (P1 \rightarrow P8) \rightarrow P6 \rightarrow P7$
 $\rightarrow (((P3 \rightarrow P2) \rightarrow P9) \rightarrow P4)$
 $\rightarrow (P1 \rightarrow P3) \rightarrow (((P6 \rightarrow P1) \rightarrow P2) \rightarrow P9) \rightarrow P5 \rangle$
by *fast*

lemma *tatsuta1*:

$\langle (((P8 \rightarrow P2) \rightarrow P9) \rightarrow P3 \rightarrow P10)$
 $\rightarrow (((P3 \rightarrow P2) \rightarrow P9) \rightarrow P4)$
 $\rightarrow (((P6 \rightarrow P1) \rightarrow P2) \rightarrow P9)$
 $\rightarrow (((P7 \rightarrow P1) \rightarrow P10) \rightarrow P4 \rightarrow P5)$
 $\rightarrow (P1 \rightarrow P3) \rightarrow (P1 \rightarrow P8) \rightarrow P6 \rightarrow P7 \rightarrow P5 \rangle$
by *fast*

end

13 First-Order Logic: quantifier examples (classical version)

theory *Quantifiers-Cla*

imports *FOL*

begin

lemma $\langle (\forall x y. P(x,y)) \rightarrow (\forall y x. P(x,y)) \rangle$
by *fast*

lemma $\langle (\exists x y. P(x,y)) \rightarrow (\exists y x. P(x,y)) \rangle$
by *fast*

Converse is false.

lemma $\langle (\forall x. P(x)) \vee (\forall x. Q(x)) \rightarrow (\forall x. P(x) \vee Q(x)) \rangle$
by *fast*

lemma $\langle (\forall x. P \rightarrow Q(x)) \longleftrightarrow (P \rightarrow (\forall x. Q(x))) \rangle$
by *fast*

lemma $\langle (\forall x. P(x) \rightarrow Q) \longleftrightarrow ((\exists x. P(x)) \rightarrow Q) \rangle$
by *fast*

Some harder ones.

lemma $\langle (\exists x. P(x) \vee Q(x)) \longleftrightarrow (\exists x. P(x)) \vee (\exists x. Q(x)) \rangle$
by *fast*

— Converse is false.

lemma $\langle (\exists x. P(x) \wedge Q(x)) \longrightarrow (\exists x. P(x)) \wedge (\exists x. Q(x)) \rangle$
by *fast*

Basic test of quantifier reasoning.

lemma $\langle (\exists y. \forall x. Q(x,y)) \longrightarrow (\forall x. \exists y. Q(x,y)) \rangle$
by *fast*

lemma $\langle (\forall x. Q(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by *fast*

The following should fail, as they are false!

lemma $\langle (\forall x. \exists y. Q(x,y)) \longrightarrow (\exists y. \forall x. Q(x,y)) \rangle$
apply *fast?*
oops

lemma $\langle (\exists x. Q(x)) \longrightarrow (\forall x. Q(x)) \rangle$
apply *fast?*
oops

schematic-goal $\langle P(?a) \longrightarrow (\forall x. P(x)) \rangle$
apply *fast?*
oops

schematic-goal $\langle (P(?a) \longrightarrow (\forall x. Q(x))) \longrightarrow (\forall x. P(x) \longrightarrow Q(x)) \rangle$
apply *fast?*
oops

Back to things that are provable ...

lemma $\langle (\forall x. P(x) \longrightarrow Q(x)) \wedge (\exists x. P(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by *fast*

An example of why *exI* should be delayed as long as possible.

lemma $\langle (P \longrightarrow (\exists x. Q(x))) \wedge P \longrightarrow (\exists x. Q(x)) \rangle$
by *fast*

schematic-goal $\langle (\forall x. P(x) \longrightarrow Q(f(x))) \wedge (\forall x. Q(x) \longrightarrow R(g(x))) \wedge P(d) \longrightarrow R(?a) \rangle$
by *fast*

lemma $\langle (\forall x. Q(x)) \longrightarrow (\exists x. Q(x)) \rangle$
by *fast*

Some slow ones

Principia Mathematica *11.53

lemma $\langle (\forall x y. P(x) \longrightarrow Q(y)) \longleftrightarrow ((\exists x. P(x)) \longrightarrow (\forall y. Q(y))) \rangle$
by *fast*

lemma $\langle (\exists x. y. P(x) \wedge Q(x,y)) \longleftrightarrow (\exists x. P(x) \wedge (\exists y. Q(x,y))) \rangle$
by *fast*

lemma $\langle (\exists y. \forall x. P(x) \longrightarrow Q(x,y)) \longrightarrow (\forall x. P(x) \longrightarrow (\exists y. Q(x,y))) \rangle$
by *fast*

end

theory *Miniscope*
imports *FOL*
begin

lemmas *ccontr* = *FalseE* [*THEN classical*]

13.1 Negation Normal Form

13.1.1 de Morgan laws

lemma *demorgans1*:
 $\langle \neg (P \wedge Q) \longleftrightarrow \neg P \vee \neg Q \rangle$
 $\langle \neg (P \vee Q) \longleftrightarrow \neg P \wedge \neg Q \rangle$
 $\langle \neg \neg P \longleftrightarrow P \rangle$
by *blast+*

lemma *demorgans2*:
 $\langle \bigwedge P. \neg (\forall x. P(x)) \longleftrightarrow (\exists x. \neg P(x)) \rangle$
 $\langle \bigwedge P. \neg (\exists x. P(x)) \longleftrightarrow (\forall x. \neg P(x)) \rangle$
by *blast+*

lemmas *demorgans* = *demorgans1 demorgans2*

lemma *nnf-simps*:
 $\langle (P \longrightarrow Q) \longleftrightarrow (\neg P \vee Q) \rangle$
 $\langle \neg (P \longrightarrow Q) \longleftrightarrow (P \wedge \neg Q) \rangle$
 $\langle (P \longleftrightarrow Q) \longleftrightarrow (\neg P \vee Q) \wedge (\neg Q \vee P) \rangle$
 $\langle \neg (P \longleftrightarrow Q) \longleftrightarrow (P \vee Q) \wedge (\neg P \vee \neg Q) \rangle$
by *blast+*

13.1.2 Pushing in the existential quantifiers

lemma *ex-simps*:
 $\langle (\exists x. P) \longleftrightarrow P \rangle$
 $\langle \bigwedge P. Q. (\exists x. P(x) \wedge Q) \longleftrightarrow (\exists x. P(x)) \wedge Q \rangle$
 $\langle \bigwedge P. Q. (\exists x. P \wedge Q(x)) \longleftrightarrow P \wedge (\exists x. Q(x)) \rangle$

```

⟨ $\bigwedge P Q. (\exists x. P(x) \vee Q(x)) \longleftrightarrow (\exists x. P(x)) \vee (\exists x. Q(x))$ ⟩
⟨ $\bigwedge P Q. (\exists x. P(x) \vee Q) \longleftrightarrow (\exists x. P(x)) \vee Q$ ⟩
⟨ $\bigwedge P Q. (\exists x. P \vee Q(x)) \longleftrightarrow P \vee (\exists x. Q(x))$ ⟩
by blast+

```

13.1.3 Pushing in the universal quantifiers

lemma *all-simps*:

```

⟨ $(\forall x. P) \longleftrightarrow P$ ⟩
⟨ $\bigwedge P Q. (\forall x. P(x) \wedge Q(x)) \longleftrightarrow (\forall x. P(x)) \wedge (\forall x. Q(x))$ ⟩
⟨ $\bigwedge P Q. (\forall x. P(x) \wedge Q) \longleftrightarrow (\forall x. P(x)) \wedge Q$ ⟩
⟨ $\bigwedge P Q. (\forall x. P \wedge Q(x)) \longleftrightarrow P \wedge (\forall x. Q(x))$ ⟩
⟨ $\bigwedge P Q. (\forall x. P(x) \vee Q) \longleftrightarrow (\forall x. P(x)) \vee Q$ ⟩
⟨ $\bigwedge P Q. (\forall x. P \vee Q(x)) \longleftrightarrow P \vee (\forall x. Q(x))$ ⟩
by blast+

```

lemmas *mini-simps = demorgans nnf-simps ex-simps all-simps*

ML <

```

val mini-ss = simpset-of (context |> Simplifier.add-simps @{thms mini-simps});
fun mini-tac ctxt =
  resolve-tac ctxt @{thms ccontr} THEN' asm-full-simp-tac (put-simpset mini-ss
  ctxt);
>

```

end

14 First-Order Logic: the 'if' example

theory *If*

imports *FOL*

begin

definition *if* :: $\langle [o, o, o] \Rightarrow o \rangle$

where $\langle if(P, Q, R) \equiv P \wedge Q \vee \neg P \wedge R \rangle$

lemma *ifI*: $\langle \llbracket P \Longrightarrow Q; \neg P \Longrightarrow R \rrbracket \Longrightarrow if(P, Q, R) \rangle$

unfolding *if-def* **by** *blast*

lemma *ifE*: $\langle \llbracket if(P, Q, R); \llbracket P; Q \rrbracket \Longrightarrow S; \llbracket \neg P; R \rrbracket \Longrightarrow S \rrbracket \Longrightarrow S \rangle$

unfolding *if-def* **by** *blast*

lemma *if-commute*: $\langle if(P, if(Q, A, B), if(Q, C, D)) \longleftrightarrow if(Q, if(P, A, C), if(P, B, D)) \rangle$

apply (*rule iffI*)

apply (*erule ifE*)

apply (*erule ifE*)

apply (*rule ifI*)

apply (*rule ifI*)

oops

Trying again from the beginning in order to use *blast*

declare *ifI* [*intro!*]

declare *ifE* [*elim!*]

lemma *if-commute*: $\langle \text{if}(P, \text{if}(Q, A, B), \text{if}(Q, C, D)) \longleftrightarrow \text{if}(Q, \text{if}(P, A, C), \text{if}(P, B, D)) \rangle$
by *blast*

lemma $\langle \text{if}(\text{if}(P, Q, R), A, B) \longleftrightarrow \text{if}(P, \text{if}(Q, A, B), \text{if}(R, A, B)) \rangle$
by *blast*

Trying again from the beginning in order to prove from the definitions

lemma $\langle \text{if}(\text{if}(P, Q, R), A, B) \longleftrightarrow \text{if}(P, \text{if}(Q, A, B), \text{if}(R, A, B)) \rangle$
unfolding *if-def* **by** *blast*

An invalid formula. High-level rules permit a simpler diagnosis.

lemma $\langle \text{if}(\text{if}(P, Q, R), A, B) \longleftrightarrow \text{if}(P, \text{if}(Q, A, B), \text{if}(R, B, A)) \rangle$
apply *auto*
— The next step will fail unless subgoals remain
apply (*tactic all-tac*)
oops

Trying again from the beginning in order to prove from the definitions.

lemma $\langle \text{if}(\text{if}(P, Q, R), A, B) \longleftrightarrow \text{if}(P, \text{if}(Q, A, B), \text{if}(R, B, A)) \rangle$
unfolding *if-def*
apply *auto*
— The next step will fail unless subgoals remain
apply (*tactic all-tac*)
oops

end