

Hoare Logic for Parallel Programs

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Abstract

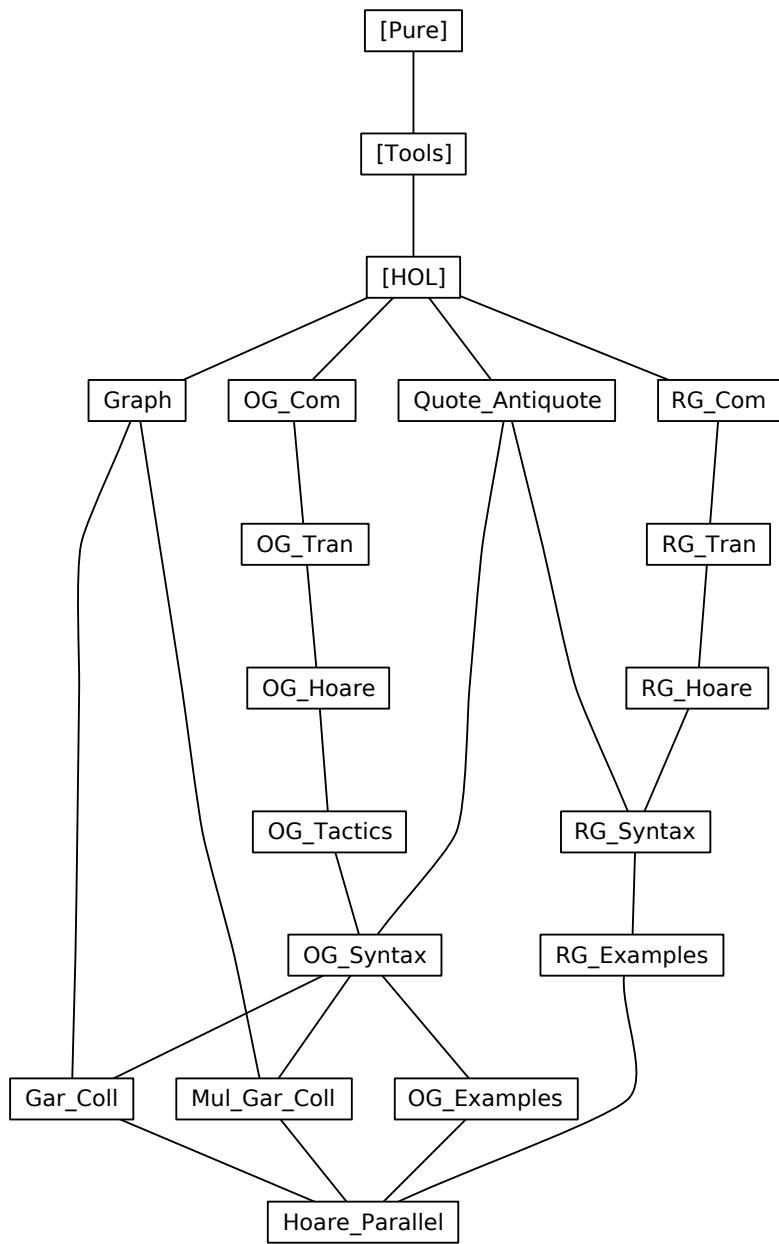
In the following theories a formalization of the Owicki-Gries and the rely-guarantee methods is presented. These methods are widely used for correctness proofs of parallel imperative programs with shared variables. We define syntax, semantics and proof rules in Isabelle/HOL. The proof rules also provide for programs parameterized in the number of parallel components. Their correctness w.r.t. the semantics is proven. Completeness proofs for both methods are extended to the new case of parameterized programs. (These proofs have not been formalized in Isabelle. They can be found in [1].) Using this formalizations we verify several non-trivial examples for parameterized and non-parameterized programs. For the automatic generation of verification conditions with the Owicki-Gries method we define a tactic based on the proof rules. The most involved examples are the verification of two garbage-collection algorithms, the second one parameterized in the number of mutators.

For excellent descriptions of this work see [2, 4, 1, 3].

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Chapter 1

The Owicki-Gries Method

1.1 Abstract Syntax

theory *OG-Com* **imports** *Main* **begin**

Type abbreviations for boolean expressions and assertions:

type-synonym *'a bexp* = *'a set*

type-synonym *'a assn* = *'a set*

The syntax of commands is defined by two mutually recursive datatypes: *'a ann-com* for annotated commands and *'a com* for non-annotated commands.

datatype *'a ann-com* =
 AnnBasic (*'a assn*) (*'a \Rightarrow 'a*)
 | *AnnSeq* (*'a ann-com*) (*'a ann-com*)
 | *AnnCond1* (*'a assn*) (*'a bexp*) (*'a ann-com*) (*'a ann-com*)
 | *AnnCond2* (*'a assn*) (*'a bexp*) (*'a ann-com*)
 | *AnnWhile* (*'a assn*) (*'a bexp*) (*'a assn*) (*'a ann-com*)
 | *AnnAwait* (*'a assn*) (*'a bexp*) (*'a com*)
and *'a com* =
 Parallel (*'a ann-com option* \times *'a assn*) *list*
 | *Basic* (*'a \Rightarrow 'a*)
 | *Seq* (*'a com*) (*'a com*)
 | *Cond* (*'a bexp*) (*'a com*) (*'a com*)
 | *While* (*'a bexp*) (*'a assn*) (*'a com*)

The function *pre* extracts the precondition of an annotated command:

primrec *pre* :: *'a ann-com \Rightarrow 'a assn* **where**
 pre (*AnnBasic* *r f*) = *r*
 | *pre* (*AnnSeq* *c1 c2*) = *pre c1*
 | *pre* (*AnnCond1* *r b c1 c2*) = *r*
 | *pre* (*AnnCond2* *r b c*) = *r*
 | *pre* (*AnnWhile* *r b i c*) = *r*
 | *pre* (*AnnAwait* *r b c*) = *r*

Well-formedness predicate for atomic programs:

primrec *atom-com* :: 'a com \Rightarrow bool **where**
 | *atom-com* (Parallel Ts) = False
 | *atom-com* (Basic f) = True
 | *atom-com* (Seq c1 c2) = (*atom-com* c1 \wedge *atom-com* c2)
 | *atom-com* (Cond b c1 c2) = (*atom-com* c1 \wedge *atom-com* c2)
 | *atom-com* (While b i c) = *atom-com* c

end

1.2 Operational Semantics

theory *OG-Tran* **imports** *OG-Com* **begin**

type-synonym 'a ann-com-op = ('a ann-com) option
type-synonym 'a ann-triple-op = ('a ann-com-op \times 'a assn)

primrec *com* :: 'a ann-triple-op \Rightarrow 'a ann-com-op **where**
com (c, q) = c

primrec *post* :: 'a ann-triple-op \Rightarrow 'a assn **where**
post (c, q) = q

definition *All-None* :: 'a ann-triple-op list \Rightarrow bool **where**
All-None Ts $\equiv \forall (c, q) \in \text{set } Ts. c = \text{None}$

1.2.1 The Transition Relation

inductive-set

ann-transition :: (('a ann-com-op \times 'a) \times ('a ann-com-op \times 'a)) set
and *transition* :: (('a com \times 'a) \times ('a com \times 'a)) set
and *ann-transition'* :: ('a ann-com-op \times 'a) \Rightarrow ('a ann-com-op \times 'a) \Rightarrow bool
 ($\langle \cdot \rangle -1 \rightarrow \rightarrow$ [81,81] 100)
and *transition'* :: ('a com \times 'a) \Rightarrow ('a com \times 'a) \Rightarrow bool
 ($\langle \cdot \rangle -P1 \rightarrow \rightarrow$ [81,81] 100)
and *transitions* :: ('a com \times 'a) \Rightarrow ('a com \times 'a) \Rightarrow bool
 ($\langle \cdot \rangle -P* \rightarrow \rightarrow$ [81,81] 100)

where

con-0 $-1 \rightarrow$ *con-1* $\equiv (con-0, con-1) \in \text{ann-transition}$
 | *con-0* $-P1 \rightarrow$ *con-1* $\equiv (con-0, con-1) \in \text{transition}$
 | *con-0* $-P* \rightarrow$ *con-1* $\equiv (con-0, con-1) \in \text{transition}^*$
 | *AnnBasic*: (Some (AnnBasic r f), s) $-1 \rightarrow$ (None, f s)
 | *AnnSeq1*: (Some c0, s) $-1 \rightarrow$ (None, t) \implies
 (Some (AnnSeq c0 c1), s) $-1 \rightarrow$ (Some c1, t)
 | *AnnSeq2*: (Some c0, s) $-1 \rightarrow$ (Some c2, t) \implies
 (Some (AnnSeq c0 c1), s) $-1 \rightarrow$ (Some (AnnSeq c2 c1), t)
 | *AnnCond1T*: s $\in b \implies$ (Some (AnnCond1 r b c1 c2), s) $-1 \rightarrow$ (Some c1, s)

| *AnnCond1F*: $s \notin b \implies (\text{Some } (\text{AnnCond1 } r \ b \ c1 \ c2), s) -1 \rightarrow (\text{Some } c2, s)$

| *AnnCond2T*: $s \in b \implies (\text{Some } (\text{AnnCond2 } r \ b \ c), s) -1 \rightarrow (\text{Some } c, s)$

| *AnnCond2F*: $s \notin b \implies (\text{Some } (\text{AnnCond2 } r \ b \ c), s) -1 \rightarrow (\text{None}, s)$

| *AnnWhileF*: $s \notin b \implies (\text{Some } (\text{AnnWhile } r \ b \ i \ c), s) -1 \rightarrow (\text{None}, s)$

| *AnnWhileT*: $s \in b \implies (\text{Some } (\text{AnnWhile } r \ b \ i \ c), s) -1 \rightarrow$
 $(\text{Some } (\text{AnnSeq } c \ (\text{AnnWhile } i \ b \ i \ c)), s)$

| *AnnAwait*: $\llbracket s \in b; \text{atom-com } c; (c, s) -P^* \rightarrow (\text{Parallel } [], t) \rrbracket \implies$
 $(\text{Some } (\text{AnnAwait } r \ b \ c), s) -1 \rightarrow (\text{None}, t)$

| *Parallel*: $\llbracket i < \text{length } Ts; Ts!i = (\text{Some } c, q); (\text{Some } c, s) -1 \rightarrow (r, t) \rrbracket$
 $\implies (\text{Parallel } Ts, s) -P1 \rightarrow (\text{Parallel } (Ts [i:= (r, q)]), t)$

| *Basic*: $(\text{Basic } f, s) -P1 \rightarrow (\text{Parallel } [], f \ s)$

| *Seq1*: $\text{All-None } Ts \implies (\text{Seq } (\text{Parallel } Ts) \ c, s) -P1 \rightarrow (c, s)$

| *Seq2*: $(c0, s) -P1 \rightarrow (c2, t) \implies (\text{Seq } c0 \ c1, s) -P1 \rightarrow (\text{Seq } c2 \ c1, t)$

| *CondT*: $s \in b \implies (\text{Cond } b \ c1 \ c2, s) -P1 \rightarrow (c1, s)$

| *CondF*: $s \notin b \implies (\text{Cond } b \ c1 \ c2, s) -P1 \rightarrow (c2, s)$

| *WhileF*: $s \notin b \implies (\text{While } b \ i \ c, s) -P1 \rightarrow (\text{Parallel } [], s)$

| *WhileT*: $s \in b \implies (\text{While } b \ i \ c, s) -P1 \rightarrow (\text{Seq } c \ (\text{While } b \ i \ c), s)$

monos *rtrancl-mono*

The corresponding abbreviations are:

abbreviation

ann-transition-n :: $('a \ \text{ann-com-op} \times 'a) \Rightarrow \text{nat} \Rightarrow ('a \ \text{ann-com-op} \times 'a)$
 $\Rightarrow \text{bool } (\leftarrow \dashrightarrow \rightarrow [81,81] \ 100) \ \mathbf{where}$
con-0 $-n \rightarrow$ *con-1* $\equiv (\text{con-0}, \text{con-1}) \in \text{ann-transition} \overset{\sim}{\sim} n$

abbreviation

ann-transitions :: $('a \ \text{ann-com-op} \times 'a) \Rightarrow ('a \ \text{ann-com-op} \times 'a) \Rightarrow \text{bool}$
 $(\leftarrow \dashrightarrow \rightarrow [81,81] \ 100) \ \mathbf{where}$
con-0 $-* \rightarrow$ *con-1* $\equiv (\text{con-0}, \text{con-1}) \in \text{ann-transition}^*$

abbreviation

transition-n :: $('a \ \text{com} \times 'a) \Rightarrow \text{nat} \Rightarrow ('a \ \text{com} \times 'a) \Rightarrow \text{bool}$
 $(\leftarrow -P \rightarrow \rightarrow [81,81,81] \ 100) \ \mathbf{where}$
con-0 $-Pn \rightarrow$ *con-1* $\equiv (\text{con-0}, \text{con-1}) \in \text{transition} \overset{\sim}{\sim} n$

1.2.2 Definition of Semantics

definition *ann-sem* :: $'a \ \text{ann-com} \Rightarrow 'a \Rightarrow 'a \ \text{set} \ \mathbf{where}$

ann-sem $c \equiv \lambda s. \{t. (\text{Some } c, s) -* \rightarrow (\text{None}, t)\}$

definition *ann-SEM* :: 'a ann-com \Rightarrow 'a set \Rightarrow 'a set **where**
ann-SEM c S $\equiv \bigcup$ (*ann-sem* c ' S)

definition *sem* :: 'a com \Rightarrow 'a \Rightarrow 'a set **where**
sem c $\equiv \lambda s. \{t. \exists Ts. (c, s) -P^* \rightarrow (Parallel\ Ts, t) \wedge AllNone\ Ts\}$

definition *SEM* :: 'a com \Rightarrow 'a set \Rightarrow 'a set **where**
SEM c S $\equiv \bigcup$ (*sem* c ' S)

abbreviation *Omega* :: 'a com $\langle \Omega \rangle$ 63
where $\Omega \equiv While\ UNIV\ UNIV$ (*Basic id*)

primrec *fwhile* :: 'a bexp \Rightarrow 'a com \Rightarrow nat \Rightarrow 'a com **where**
fwhile b c 0 = Ω
| *fwhile* b c (Suc n) = *Cond* b (*Seq* c (*fwhile* b c n)) (*Basic id*)

Proofs

declare *ann-transition-transition.intros* [intro]

inductive-cases *transition-cases*:

(*Parallel* T, s) -P1 \rightarrow t
(*Basic* f, s) -P1 \rightarrow t
(*Seq* c1 c2, s) -P1 \rightarrow t
(*Cond* b c1 c2, s) -P1 \rightarrow t
(*While* b i c, s) -P1 \rightarrow t

lemma *Parallel-empty-lemma* [rule-format (no-asm)]:

(*Parallel* [], s) -Pn \rightarrow (*Parallel* Ts, t) \longrightarrow Ts=[] \wedge n=0 \wedge s=t

apply(*induct* n)

apply(*simp* (no-asm))

apply *clarify*

apply(*drule* *relpow-Suc-D2*)

apply(*force* *elim:transition-cases*)

done

lemma *Parallel-AllNone-lemma* [rule-format (no-asm)]:

AllNone Ss \longrightarrow (*Parallel* Ss, s) -Pn \rightarrow (*Parallel* Ts, t) \longrightarrow Ts=Ss \wedge n=0 \wedge s=t

apply(*induct* n)

apply(*simp* (no-asm))

apply *clarify*

apply(*drule* *relpow-Suc-D2*)

apply *clarify*

apply(*erule* *transition-cases, simp-all*)

apply(*force* *dest:nth-mem simp add:AllNone-def*)

done

lemma *Parallel-AllNone*: *AllNone* Ts \Longrightarrow (*SEM* (*Parallel* Ts) X) = X

apply (*unfold* *SEM-def sem-def*)

apply *auto*

apply(*drule rtrancl-imp-UN-relpow*)
apply *clarify*
apply(*drule Parallel-AllNone-lemma*)
apply *auto*
done

lemma *Parallel-empty*: $Ts = [] \implies (SEM (Parallel Ts) X) = X$
apply(*rule Parallel-AllNone*)
apply(*simp add:All-None-def*)
done

Set of lemmas from Apt and Olderog "Verification of sequential and concurrent programs", page 63.

lemma *L3-5i*: $X \subseteq Y \implies SEM\ c\ X \subseteq SEM\ c\ Y$
apply (*unfold SEM-def*)
apply *force*
done

lemma *L3-5ii-lemma1*:

$$\begin{aligned} & \llbracket (c1, s1) -P^* \rightarrow (Parallel\ Ts, s2); All-None\ Ts; \\ & (c2, s2) -P^* \rightarrow (Parallel\ Ss, s3); All-None\ Ss \rrbracket \\ & \implies (Seq\ c1\ c2, s1) -P^* \rightarrow (Parallel\ Ss, s3) \end{aligned}$$

apply(*erule converse-rtrancl-induct2*)
apply(*force intro:converse-rtrancl-into-rtrancl*)
done

lemma *L3-5ii-lemma2* [*rule-format (no-asm)*]:

$$\begin{aligned} \forall c1\ c2\ s\ t. & (Seq\ c1\ c2, s) -Pn \rightarrow (Parallel\ Ts, t) \longrightarrow \\ & (All-None\ Ts) \longrightarrow (\exists y\ m\ Rs. (c1, s) -P^* \rightarrow (Parallel\ Rs, y) \wedge \\ & (All-None\ Rs) \wedge (c2, y) -Pm \rightarrow (Parallel\ Ts, t) \wedge m \leq n) \end{aligned}$$

apply(*induct n*)
apply(*force*)
apply(*safe dest!: relpow-Suc-D2*)
apply(*erule transition-cases, simp-all*)
apply (*fast intro!: le-SucI*)
apply (*fast intro!: le-SucI elim!: relpow-imp-rtrancl converse-rtrancl-into-rtrancl*)
done

lemma *L3-5ii-lemma3*:

$$\begin{aligned} & \llbracket (Seq\ c1\ c2, s) -P^* \rightarrow (Parallel\ Ts, t); All-None\ Ts \rrbracket \implies \\ & (\exists y\ Rs. (c1, s) -P^* \rightarrow (Parallel\ Rs, y) \wedge All-None\ Rs \\ & \wedge (c2, y) -P^* \rightarrow (Parallel\ Ts, t)) \end{aligned}$$

apply(*drule rtrancl-imp-UN-relpow*)
apply(*fast dest: L3-5ii-lemma2 relpow-imp-rtrancl*)
done

lemma *L3-5ii*: $SEM (Seq\ c1\ c2) X = SEM\ c2 (SEM\ c1\ X)$
apply (*unfold SEM-def sem-def*)
apply *auto*

```

apply(fast dest: L3-5ii-lemma3)
apply(fast elim: L3-5ii-lemma1)
done

```

```

lemma L3-5iii: SEM (Seq (Seq c1 c2) c3) X = SEM (Seq c1 (Seq c2 c3)) X
apply (simp (no-asm) add: L3-5ii)
done

```

```

lemma L3-5iv:
  SEM (Cond b c1 c2) X = (SEM c1 (X ∩ b)) Un (SEM c2 (X ∩ (-b)))
apply (unfold SEM-def sem-def)
apply auto
apply(erule converse-rtranclE)
prefer 2
apply (erule transition-cases,simp-all)
apply(fast intro: converse-rtrancl-into-rtrancl elim: transition-cases)+
done

```

```

lemma L3-5v-lemma1[rule-format]:
  (S,s) -Pn→ (T,t) → S=Ω → (¬(∃ Rs. T=(Parallel Rs) ∧ All-None Rs))
apply (unfold UNIV-def)
apply(rule nat-less-induct)
apply safe
apply(erule relpow-E2)
apply simp-all
apply(erule transition-cases)
apply simp-all
apply(erule relpow-E2)
apply(simp add: Id-def)
apply(erule transition-cases,simp-all)
apply clarify
apply(erule transition-cases,simp-all)
apply(erule relpow-E2,simp)
apply clarify
apply(erule transition-cases)
apply simp+
apply clarify
apply(erule transition-cases)
apply simp-all
done

```

```

lemma L3-5v-lemma2: [(Ω, s) -P*→ (Parallel Ts, t); All-None Ts] ⇒ False
apply(fast dest: rtrancl-imp-UN-relpow L3-5v-lemma1)
done

```

```

lemma L3-5v-lemma3: SEM (Ω) S = {}
apply (unfold SEM-def sem-def)
apply(fast dest: L3-5v-lemma2)

```

done

lemma *L3-5v-lemma4* [rule-format]:

$\forall s. (\text{While } b \ i \ c, \ s) -Pn \rightarrow (\text{Parallel } Ts, \ t) \longrightarrow \text{All-None } Ts \longrightarrow$
 $(\exists k. (\text{fwhile } b \ c \ k, \ s) -P* \rightarrow (\text{Parallel } Ts, \ t))$

apply(rule nat-less-induct)

apply safe

apply(erule relpow-E2)

apply safe

apply(erule transition-cases,simp-all)

apply (rule-tac $x = 1$ in exI)

apply(force dest: Parallel-empty-lemma intro: converse-rtrancl-into-rtrancl simp
add: Id-def)

apply safe

apply(erule L3-5ii-lemma2)

apply safe

apply(erule le-imp-less-Suc)

apply (erule allE , erule impE,assumption)

apply (erule allE , erule impE, assumption)

apply safe

apply (rule-tac $x = k+1$ in exI)

apply(simp (no-asm))

apply(rule converse-rtrancl-into-rtrancl)

apply fast

apply(fast elim: L3-5ii-lemma1)

done

lemma *L3-5v-lemma5* [rule-format]:

$\forall s. (\text{fwhile } b \ c \ k, \ s) -P* \rightarrow (\text{Parallel } Ts, \ t) \longrightarrow \text{All-None } Ts \longrightarrow$
 $(\text{While } b \ i \ c, \ s) -P* \rightarrow (\text{Parallel } Ts, \ t)$

apply(induct k)

apply(force dest: L3-5v-lemma2)

apply safe

apply(erule converse-rtranclE)

apply simp-all

apply(erule transition-cases,simp-all)

apply(rule converse-rtrancl-into-rtrancl)

apply(fast)

apply(fast elim!: L3-5ii-lemma1 dest: L3-5ii-lemma3)

apply(erule rtrancl-imp-UN-relpow)

apply clarify

apply(erule relpow-E2)

apply simp-all

apply(erule transition-cases,simp-all)

apply(force dest: Parallel-empty-lemma)

done

lemma *L3-5v*: SEM (While $b \ i \ c$) = $(\lambda x. (\bigcup k. \text{SEM } (\text{fwhile } b \ c \ k) \ x))$

apply(rule ext)

```

apply (simp add: SEM-def sem-def)
apply safe
apply(drule rtrancl-imp-UN-relpow,simp)
apply clarify
apply(fast dest:L3-5v-lemma4)
apply(fast intro: L3-5v-lemma5)
done

```

1.3 Validity of Correctness Formulas

definition *com-validity* :: 'a assn \Rightarrow 'a com \Rightarrow 'a assn \Rightarrow bool ($\langle \langle \mathcal{P} \models - // - // - \rangle \rangle$
 $[90,55,90] 50$) **where**
 $\models p c q \equiv SEM c p \subseteq q$

definition *ann-com-validity* :: 'a ann-com \Rightarrow 'a assn \Rightarrow bool ($\langle \models - - \rangle [60,90] 45$)
where
 $\models c q \equiv ann-SEM c (pre c) \subseteq q$

end

1.4 The Proof System

theory *OG-Hoare* **imports** *OG-Tran* **begin**

primrec *assertions* :: 'a ann-com \Rightarrow ('a assn) set **where**
 $assertions (AnnBasic r f) = \{r\}$
 $| assertions (AnnSeq c1 c2) = assertions c1 \cup assertions c2$
 $| assertions (AnnCond1 r b c1 c2) = \{r\} \cup assertions c1 \cup assertions c2$
 $| assertions (AnnCond2 r b c) = \{r\} \cup assertions c$
 $| assertions (AnnWhile r b i c) = \{r, i\} \cup assertions c$
 $| assertions (AnnAwait r b c) = \{r\}$

primrec *atomics* :: 'a ann-com \Rightarrow ('a assn \times 'a com) set **where**
 $atomics (AnnBasic r f) = \{(r, Basic f)\}$
 $| atomics (AnnSeq c1 c2) = atomics c1 \cup atomics c2$
 $| atomics (AnnCond1 r b c1 c2) = atomics c1 \cup atomics c2$
 $| atomics (AnnCond2 r b c) = atomics c$
 $| atomics (AnnWhile r b i c) = atomics c$
 $| atomics (AnnAwait r b c) = \{(r \cap b, c)\}$

primrec *com* :: 'a ann-triple-op \Rightarrow 'a ann-com-op **where**
 $com (c, q) = c$

primrec *post* :: 'a ann-triple-op \Rightarrow 'a assn **where**
 $post (c, q) = q$

definition *interfree-aux* :: ('a ann-com-op \times 'a assn \times 'a ann-com-op) \Rightarrow bool
where

$interfree\text{-aux} \equiv \lambda(co, q, co'). co' = None \vee$
 $(\forall (r, a) \in atomics (the\ co'). \models (q \cap r) a q \wedge$
 $(co = None \vee (\forall p \in assertions (the\ co). \models (p \cap r) a p)))$

definition $interfree :: (('a\ ann\ triple\ op)\ list) \Rightarrow bool$ **where**
 $interfree\ Ts \equiv \forall i\ j. i < length\ Ts \wedge j < length\ Ts \wedge i \neq j \longrightarrow$
 $interfree\text{-aux}\ (com\ (Ts!i), post\ (Ts!i), com\ (Ts!j))$

inductive

$oghoare :: 'a\ assn \Rightarrow 'a\ com \Rightarrow 'a\ assn \Rightarrow bool$ $(\langle (3 \parallel - \ / \ / \ / -) \rangle [90, 55, 90] 50)$
and $ann\ hoare :: 'a\ ann\ com \Rightarrow 'a\ assn \Rightarrow bool$ $(\langle (2 \vdash - \ / \ / -) \rangle [60, 90] 45)$

where

$AnnBasic: r \subseteq \{s. f\ s \in q\} \Longrightarrow \vdash (AnnBasic\ r\ f)\ q$

| $AnnSeq: \llbracket \vdash\ c0\ pre\ c1; \vdash\ c1\ q \rrbracket \Longrightarrow \vdash (AnnSeq\ c0\ c1)\ q$

| $AnnCond1: \llbracket r \cap b \subseteq pre\ c1; \vdash\ c1\ q; r \cap \neg b \subseteq pre\ c2; \vdash\ c2\ q \rrbracket$
 $\Longrightarrow \vdash (AnnCond1\ r\ b\ c1\ c2)\ q$

| $AnnCond2: \llbracket r \cap b \subseteq pre\ c; \vdash\ c\ q; r \cap \neg b \subseteq q \rrbracket \Longrightarrow \vdash (AnnCond2\ r\ b\ c)\ q$

| $AnnWhile: \llbracket r \subseteq i; i \cap b \subseteq pre\ c; \vdash\ c\ i; i \cap \neg b \subseteq q \rrbracket$
 $\Longrightarrow \vdash (AnnWhile\ r\ b\ i\ c)\ q$

| $AnnAwait: \llbracket atom\ com\ c; \parallel - (r \cap b)\ c\ q \rrbracket \Longrightarrow \vdash (AnnAwait\ r\ b\ c)\ q$

| $AnnConseq: \llbracket \vdash\ c\ q; q \subseteq q' \rrbracket \Longrightarrow \vdash\ c\ q'$

| $Parallel: \llbracket \forall i < length\ Ts. \exists c\ q. Ts!i = (Some\ c, q) \wedge \vdash\ c\ q; interfree\ Ts \rrbracket$
 $\Longrightarrow \parallel - (\bigcap i \in \{i. i < length\ Ts\}. pre(the(com(Ts!i))))$
 $Parallel\ Ts$
 $(\bigcap i \in \{i. i < length\ Ts\}. post(Ts!i))$

| $Basic: \parallel - \{s. f\ s \in q\} (Basic\ f)\ q$

| $Seq: \llbracket \parallel -\ p\ c1\ r; \parallel -\ r\ c2\ q \rrbracket \Longrightarrow \parallel -\ p\ (Seq\ c1\ c2)\ q$

| $Cond: \llbracket \parallel -\ (p \cap b)\ c1\ q; \parallel -\ (p \cap \neg b)\ c2\ q \rrbracket \Longrightarrow \parallel -\ p\ (Cond\ b\ c1\ c2)\ q$

| $While: \llbracket \parallel -\ (p \cap b)\ c\ p \rrbracket \Longrightarrow \parallel -\ p\ (While\ b\ i\ c)\ (p \cap \neg b)$

| $Conseq: \llbracket p' \subseteq p; \parallel -\ p\ c\ q; q \subseteq q' \rrbracket \Longrightarrow \parallel -\ p'\ c\ q'$

1.5 Soundness

lemmas $[cong\ del] = if\ weak\ cong$

lemmas $ann\ hoare\ induct = oghoare\ ann\ hoare. induct [THEN\ conjunct2]$

lemmas $oghoare\ induct = oghoare\ ann\ hoare. induct [THEN\ conjunct1]$

lemmas *AnnBasic* = *oghoare-ann-hoare.AnnBasic*
lemmas *AnnSeq* = *oghoare-ann-hoare.AnnSeq*
lemmas *AnnCond1* = *oghoare-ann-hoare.AnnCond1*
lemmas *AnnCond2* = *oghoare-ann-hoare.AnnCond2*
lemmas *AnnWhile* = *oghoare-ann-hoare.AnnWhile*
lemmas *AnnAwait* = *oghoare-ann-hoare.AnnAwait*
lemmas *AnnConseq* = *oghoare-ann-hoare.AnnConseq*

lemmas *Parallel* = *oghoare-ann-hoare.Parallel*
lemmas *Basic* = *oghoare-ann-hoare.Basic*
lemmas *Seq* = *oghoare-ann-hoare.Seq*
lemmas *Cond* = *oghoare-ann-hoare.Cond*
lemmas *While* = *oghoare-ann-hoare.While*
lemmas *Conseq* = *oghoare-ann-hoare.Conseq*

1.5.1 Soundness of the System for Atomic Programs

lemma *Basic-ntran* [*rule-format*]:
 $(\text{Basic } f, s) -Pn \rightarrow (\text{Parallel } Ts, t) \longrightarrow \text{All-None } Ts \longrightarrow t = f s$
apply (*induct n*)
apply (*simp (no-asm)*)
apply (*fast dest: relpow-Suc-D2 Parallel-empty-lemma elim: transition-cases*)
done

lemma *SEM-fwhile*: $SEM S (p \cap b) \subseteq p \implies SEM (fwhile b S k) p \subseteq (p \cap -b)$
apply (*induct k*)
apply (*simp (no-asm) add: L3-5v-lemma3*)
apply (*simp (no-asm) add: L3-5iv L3-5ii Parallel-empty*)
apply (*rule conjI*)
apply (*blast dest: L3-5i*)
apply (*simp add: SEM-def sem-def id-def*)
apply (*auto dest: Basic-ntran rtrancl-imp-UN-relpow*)
apply *blast*
done

lemma *atom-hoare-sound* [*rule-format*]:
 $\Vdash - p c q \longrightarrow \text{atom-com}(c) \longrightarrow \Vdash = p c q$
apply (*unfold com-validity-def*)
apply (*rule oghoare-induct*)
apply *simp-all*
— *Basic*
apply (*simp add: SEM-def sem-def*)
apply (*fast dest: rtrancl-imp-UN-relpow Basic-ntran*)
— *Seq*
apply (*rule impI*)
apply (*rule subset-trans*)
prefer 2 apply simp
apply (*simp add: L3-5ii L3-5i*)

— Cond
apply(*simp add: L3-5iv*)
— While
apply (*force simp add: L3-5v dest: SEM-fwhile*)
— Conseq
apply(*force simp add: SEM-def sem-def*)
done

1.5.2 Soundness of the System for Component Programs

inductive-cases *ann-transition-cases*:

(*None,s*) $-1 \rightarrow (c', s')$
(*Some (AnnBasic r f),s*) $-1 \rightarrow (c', s')$
(*Some (AnnSeq c1 c2), s*) $-1 \rightarrow (c', s')$
(*Some (AnnCond1 r b c1 c2), s*) $-1 \rightarrow (c', s')$
(*Some (AnnCond2 r b c), s*) $-1 \rightarrow (c', s')$
(*Some (AnnWhile r b I c), s*) $-1 \rightarrow (c', s')$
(*Some (AnnAwait r b c),s*) $-1 \rightarrow (c', s')$

Strong Soundness for Component Programs:

lemma *ann-hoare-case-analysis* [*rule-format*]: $\vdash C \ q' \longrightarrow$
 $((\forall r \ f. \ C = \text{AnnBasic } r \ f \longrightarrow (\exists q. \ r \subseteq \{s. \ f \ s \in q\} \wedge q \subseteq q')) \wedge$
 $(\forall c0 \ c1. \ C = \text{AnnSeq } c0 \ c1 \longrightarrow (\exists q. \ q \subseteq q' \wedge \vdash c0 \ \text{pre } c1 \wedge \vdash c1 \ q)) \wedge$
 $(\forall r \ b \ c1 \ c2. \ C = \text{AnnCond1 } r \ b \ c1 \ c2 \longrightarrow (\exists q. \ q \subseteq q' \wedge$
 $r \cap b \subseteq \text{pre } c1 \wedge \vdash c1 \ q \wedge r \cap -b \subseteq \text{pre } c2 \wedge \vdash c2 \ q)) \wedge$
 $(\forall r \ b \ c. \ C = \text{AnnCond2 } r \ b \ c \longrightarrow$
 $(\exists q. \ q \subseteq q' \wedge r \cap b \subseteq \text{pre } c \wedge \vdash c \ q \wedge r \cap -b \subseteq q)) \wedge$
 $(\forall r \ i \ b \ c. \ C = \text{AnnWhile } r \ b \ i \ c \longrightarrow$
 $(\exists q. \ q \subseteq q' \wedge r \subseteq i \wedge i \cap b \subseteq \text{pre } c \wedge \vdash c \ i \wedge i \cap -b \subseteq q)) \wedge$
 $(\forall r \ b \ c. \ C = \text{AnnAwait } r \ b \ c \longrightarrow (\exists q. \ q \subseteq q' \wedge \Vdash (r \cap b) \ c \ q)))$
apply(*rule ann-hoare-induct*)
apply *simp-all*
apply(*rule-tac x=q in exI, simp*)+
apply(*rule conjI, clarify, simp, clarify, rule-tac x=qa in exI, fast*)+
apply(*clarify, simp, clarify, rule-tac x=qa in exI, fast*)
done

lemma *Help*: $(\text{transition} \cap \{(x,y). \ \text{True}\}) = (\text{transition})$
apply *force*
done

lemma *Strong-Soundness-aux-aux* [*rule-format*]:
 $(co, s) -1 \rightarrow (co', t) \longrightarrow (\forall c. \ co = \text{Some } c \longrightarrow s \in \text{pre } c \longrightarrow$
 $(\forall q. \ \vdash c \ q \longrightarrow (\text{if } co' = \text{None} \text{ then } t \in q \text{ else } t \in \text{pre}(\text{the } co') \wedge \vdash (\text{the } co') \ q)))$
apply(*rule ann-transition-transition.induct [THEN conjunct1]*)
apply *simp-all*
— Basic
apply *clarify*
apply(*frule ann-hoare-case-analysis*)


```

      apply force
— Seq
  apply clarify
  apply(frule ann-hoare-case-analysis,simp)
  apply(fast intro: AnnConseq)
  apply clarify
  apply(frule ann-hoare-case-analysis,simp)
  apply clarify
  apply(rule conjI)
  apply force
  apply(rule AnnSeq,simp)
  apply(fast intro: AnnConseq)
— Cond1
  apply clarify
  apply(frule ann-hoare-case-analysis,simp)
  apply(fast intro: AnnConseq)
  apply clarify
  apply(frule ann-hoare-case-analysis,simp)
  apply(fast intro: AnnConseq)
— Cond2
  apply clarify
  apply(frule ann-hoare-case-analysis,simp)
  apply(fast intro: AnnConseq)
  apply clarify
  apply(frule ann-hoare-case-analysis,simp)
  apply(fast intro: AnnConseq)
— While
  apply clarify
  apply(frule ann-hoare-case-analysis,simp)
  apply force
  apply clarify
  apply(frule ann-hoare-case-analysis,simp)
  apply auto
  apply(rule AnnSeq)
  apply simp
  apply(rule AnnWhile)
  apply simp-all
— Await
apply(frule ann-hoare-case-analysis,simp)
apply clarify
apply(drule atom-hoare-sound)
  apply simp
apply(simp add: com-validity-def SEM-def sem-def)
apply(simp add: Help All-None-def)
apply force
done

```

lemma Strong-Soundness-aux: $\llbracket (\text{Some } c, s) \text{--}^* \text{--} \rightarrow (co, t); s \in \text{pre } c; \vdash c \ q \rrbracket$
 \implies if $co = \text{None}$ then $t \in q$ else $t \in \text{pre } (the \ co) \wedge \vdash (the \ co) \ q$

```

apply(erule rtrancl-induct2)
apply simp
apply(case-tac a)
apply(fast elim: ann-transition-cases)
apply(erule Strong-Soundness-aux-aux)
apply simp
apply simp-all
done

```

```

lemma Strong-Soundness:  $\llbracket (\text{Some } c, s) \text{--}^* \text{--} \rightarrow (co, t); s \in \text{pre } c; \vdash c \ q \rrbracket$ 
 $\implies$  if  $co = \text{None}$  then  $t \in q$  else  $t \in \text{pre } (the \ co)$ 
apply(force dest:Strong-Soundness-aux)
done

```

```

lemma ann-hoare-sound:  $\vdash c \ q \implies \models c \ q$ 
apply (unfold ann-com-validity-def ann-SEM-def ann-sem-def)
apply clarify
apply(erule Strong-Soundness)
apply simp-all
done

```

1.5.3 Soundness of the System for Parallel Programs

```

lemma Parallel-length-post-P1:  $(\text{Parallel } Ts, s) \text{--}P1 \rightarrow (R', t) \implies$ 
 $(\exists Rs. R' = (\text{Parallel } Rs) \wedge (\text{length } Rs) = (\text{length } Ts) \wedge$ 
 $(\forall i. i < \text{length } Ts \longrightarrow \text{post}(Rs \ ! \ i) = \text{post}(Ts \ ! \ i)))$ 
apply(erule transition-cases)
apply simp
apply clarify
apply(case-tac i=ia)
apply simp+
done

```

```

lemma Parallel-length-post-PStar:  $(\text{Parallel } Ts, s) \text{--}P^* \rightarrow (R', t) \implies$ 
 $(\exists Rs. R' = (\text{Parallel } Rs) \wedge (\text{length } Rs) = (\text{length } Ts) \wedge$ 
 $(\forall i. i < \text{length } Ts \longrightarrow \text{post}(Ts \ ! \ i) = \text{post}(Rs \ ! \ i)))$ 
apply(erule rtrancl-induct2)
apply(simp-all)
apply clarify
apply simp
apply(erule Parallel-length-post-P1)
apply auto
done

```

```

lemma assertions-lemma:  $\text{pre } c \in \text{assertions } c$ 
apply(rule ann-com-com.induct [THEN conjunct1])
apply auto
done

```

lemma *interfree-aux1* [rule-format]:
 $(c,s) -1 \rightarrow (r,t) \longrightarrow (\text{interfree-aux}(c1, q1, c) \longrightarrow \text{interfree-aux}(c1, q1, r))$
apply (rule *ann-transition-transition.induct* [THEN *conjunct1*])
apply (*safe*)
prefer 13
apply (rule *TrueI*)
apply (*simp-all add:interfree-aux-def*)
apply *force+*
done

lemma *interfree-aux2* [rule-format]:
 $(c,s) -1 \rightarrow (r,t) \longrightarrow (\text{interfree-aux}(c, q, a) \longrightarrow \text{interfree-aux}(r, q, a))$
apply (rule *ann-transition-transition.induct* [THEN *conjunct1*])
apply (*force simp add:interfree-aux-def*)
done

lemma *interfree-lemma*: $\llbracket (\text{Some } c, s) -1 \rightarrow (r, t); \text{interfree } Ts ; i < \text{length } Ts ;$
 $Ts!i = (\text{Some } c, q) \rrbracket \Longrightarrow \text{interfree } (Ts[i := (r, q)])$
apply (*simp add: interfree-def*)
apply *clarify*
apply (*case-tac i=j*)
apply (*drule-tac t = ia in not-sym*)
apply *simp-all*
apply (*force elim: interfree-aux1*)
apply (*force elim: interfree-aux2 simp add:nth-list-update*)
done

Strong Soundness Theorem for Parallel Programs:

lemma *Parallel-Strong-Soundness-Seq-aux*:
 $\llbracket \text{interfree } Ts ; i < \text{length } Ts ; \text{com}(Ts ! i) = \text{Some}(\text{AnnSeq } c0 \ c1) \rrbracket$
 $\Longrightarrow \text{interfree } (Ts[i := (\text{Some } c0, \text{pre } c1)])$
apply (*simp add: interfree-def*)
apply *clarify*
apply (*case-tac i=j*)
apply (*force simp add: nth-list-update interfree-aux-def*)
apply (*case-tac i=ia*)
apply (*erule-tac x=ia in allE*)
apply (*force simp add:interfree-aux-def assertions-lemma*)
apply *simp*
done

lemma *Parallel-Strong-Soundness-Seq* [rule-format (*no-asm*):
 $\llbracket \forall i < \text{length } Ts. (\text{if } \text{com}(Ts!i) = \text{None} \text{ then } b \in \text{post}(Ts!i)$
 $\text{ else } b \in \text{pre}(\text{the}(\text{com}(Ts!i))) \wedge \vdash \text{the}(\text{com}(Ts!i)) \text{ post}(Ts!i);$
 $\text{com}(Ts ! i) = \text{Some}(\text{AnnSeq } c0 \ c1); i < \text{length } Ts ; \text{interfree } Ts \rrbracket \Longrightarrow$
 $(\forall ia < \text{length } Ts. (\text{if } \text{com}(Ts[i := (\text{Some } c0, \text{pre } c1)]! ia) = \text{None}$
 $\text{ then } b \in \text{post}(Ts[i := (\text{Some } c0, \text{pre } c1)]! ia)$
 $\text{ else } b \in \text{pre}(\text{the}(\text{com}(Ts[i := (\text{Some } c0, \text{pre } c1)]! ia))) \wedge$
 $\vdash \text{the}(\text{com}(Ts[i := (\text{Some } c0, \text{pre } c1)]! ia)) \text{ post}(Ts[i := (\text{Some } c0, \text{pre } c1)]! ia)))$

```

   $\wedge$  interfree ( $Ts[i:= (Some\ c0, pre\ c1)]$ )
apply(rule conjI)
apply safe
apply(case-tac i=ia)
apply simp
apply(force dest: ann-hoare-case-analysis)
apply simp
apply(fast elim: Parallel-Strong-Soundness-Seq-aux)
done

lemma Parallel-Strong-Soundness-aux-aux [rule-format]:
  ( $Some\ c, b$ )  $-1 \rightarrow (co, t) \rightarrow$ 
  ( $\forall Ts. i < length\ Ts \rightarrow com(Ts\ !\ i) = Some\ c \rightarrow$ 
  ( $\forall i < length\ Ts. (if\ com(Ts\ !\ i) = None\ then\ b \in post(Ts!\ i)$ 
   $else\ b \in pre(the(com(Ts!\ i))) \wedge \vdash the(com(Ts!\ i))\ post(Ts!\ i)) \rightarrow$ 
  interfree  $Ts \rightarrow$ 
  ( $\forall j. j < length\ Ts \wedge i \neq j \rightarrow (if\ com(Ts!\ j) = None\ then\ t \in post(Ts!\ j)$ 
   $else\ t \in pre(the(com(Ts!\ j))) \wedge \vdash the(com(Ts!\ j))\ post(Ts!\ j))$  )
apply(rule ann-transition-transition.induct [THEN conjunct1])
apply safe
prefer 11
apply(rule TrueI)
apply simp-all
— Basic
apply(erule-tac x = i in all-dupE, erule (1) notE impE)
apply(erule-tac x = j in allE, erule (1) notE impE)
apply(simp add: interfree-def)
apply(erule-tac x = j in allE, simp)
apply(erule-tac x = i in allE, simp)
apply(drule-tac t = i in not-sym)
apply(case-tac com(Ts ! j)=None)
apply(force intro: converse-rtrancl-into-rtrancl
  simp add: interfree-aux-def com-validity-def SEM-def sem-def All-None-def)
apply(simp add: interfree-aux-def)
apply clarify
apply simp
apply(erule-tac x=pre y in ballE)
apply(force intro: converse-rtrancl-into-rtrancl
  simp add: com-validity-def SEM-def sem-def All-None-def)
apply(simp add: assertions-lemma)
— Seqs
apply(erule-tac x = Ts[i:=(Some c0, pre c1)] in allE)
apply(drule Parallel-Strong-Soundness-Seq, simp+)
apply(erule-tac x = Ts[i:=(Some c0, pre c1)] in allE)
apply(drule Parallel-Strong-Soundness-Seq, simp+)
— Await
apply(rule-tac x = i in allE, assumption, erule (1) notE impE)
apply(erule-tac x = j in allE, erule (1) notE impE)
apply(simp add: interfree-def)

```

```

apply(erule-tac x = j in allE,simp)
apply(erule-tac x = i in allE,simp)
apply(erule-tac t = i in not-sym)
apply(case-tac com(Ts ! j)=None)
  apply(force intro: converse-rtrancl-into-rtrancl simp add: interfree-aux-def
    com-validity-def SEM-def sem-def All-None-def Help)
apply(simp add:interfree-aux-def)
apply clarify
apply simp
apply(erule-tac x=pre y in ballE)
  apply(force intro: converse-rtrancl-into-rtrancl
    simp add: com-validity-def SEM-def sem-def All-None-def Help)
apply(simp add:assertions-lemma)
done

```

lemma *Parallel-Strong-Soundness-aux* [rule-format]:

$$\llbracket (Ts',s) -P*\rightarrow (Rs',t); Ts' = (\text{Parallel } Ts); \text{interfree } Ts; \\ \forall i. i < \text{length } Ts \rightarrow (\exists c q. (Ts ! i) = (\text{Some } c, q) \wedge s \in (\text{pre } c) \wedge \vdash c q) \rrbracket \implies \\ \forall Rs. Rs' = (\text{Parallel } Rs) \rightarrow (\forall j. j < \text{length } Rs \rightarrow \\ (\text{if } \text{com}(Rs ! j) = \text{None} \text{ then } t \in \text{post}(Ts ! j) \\ \text{else } t \in \text{pre}(\text{the}(\text{com}(Rs ! j))) \wedge \vdash \text{the}(\text{com}(Rs ! j)) \text{ post}(Ts ! j))) \wedge \text{interfree } Rs$$

```

apply(erule rtrancl-induct2)
apply clarify
— Base
  apply force
— Induction step
apply clarify
apply(erule Parallel-length-post-PStar)
apply clarify
apply(ind-cases (Parallel Ts, s) -P1 → (Parallel Rs, t) for Ts s Rs t)
apply(rule conjI)
  apply clarify
  apply(case-tac i=j)
  apply(simp split del:if-split)
  apply(erule Strong-Soundness-aux-aux,simp+)
  apply force
  apply force
  apply(simp split del: if-split)
  apply(erule Parallel-Strong-Soundness-aux-aux)
  apply(simp-all add: split del:if-split)
  apply force
apply(rule interfree-lemma)
apply simp-all
done

```

lemma *Parallel-Strong-Soundness*:

$$\llbracket (\text{Parallel } Ts, s) -P*\rightarrow (\text{Parallel } Rs, t); \text{interfree } Ts; j < \text{length } Rs; \\ \forall i. i < \text{length } Ts \rightarrow (\exists c q. Ts ! i = (\text{Some } c, q) \wedge s \in \text{pre } c \wedge \vdash c q) \rrbracket \implies \\ \text{if } \text{com}(Rs ! j) = \text{None} \text{ then } t \in \text{post}(Ts ! j) \text{ else } t \in \text{pre}(\text{the}(\text{com}(Rs ! j)))$$

```

apply(drule Parallel-Strong-Soundness-aux)
apply simp+
done

lemma oghoare-sound [rule-format]:  $\|- p c q \longrightarrow \|= p c q$ 
apply (unfold com-validity-def)
apply(rule oghoare-induct)
apply(rule TrueI)+
— Parallel
  apply(simp add: SEM-def sem-def)
  apply(clarify, rename-tac x y i Ts')
  apply(frule Parallel-length-post-PStar)
  apply clarify
  apply(drule-tac j=i in Parallel-Strong-Soundness)
  apply clarify
  apply simp
  apply force
  apply simp
  apply(erule-tac V =  $\forall i. P i$  for P in thin-rl)
  apply(drule-tac s = length Rs in sym)
  apply(erule allE, erule impE, assumption)
  apply(force dest: nth-mem simp add: All-None-def)
— Basic
  apply(simp add: SEM-def sem-def)
  apply(force dest: rtrancl-imp-UN-relpow Basic-ntran)
— Seq
  apply(rule subset-trans)
  prefer 2 apply assumption
  apply(simp add: L3-5ii L3-5i)
— Cond
  apply(simp add: L3-5iv)
— While
  apply(simp add: L3-5v)
  apply (blast dest: SEM-fwhile)
— Conseq
apply(auto simp add: SEM-def sem-def)
done

end

```

1.6 Generation of Verification Conditions

```

theory OG-Tactics
imports OG-Hoare
begin

```

```

lemmas ann-hoare-intros=AnnBasic AnnSeq AnnCond1 AnnCond2 AnnWhile An-
nAwait AnnConseq

```

```

lemmas oghoare-intros=Parallel Basic Seq Cond While Conseq

```

lemma *ParallelConseqRule*:

$\llbracket p \subseteq (\bigcap i \in \{i. i < \text{length } Ts\}. \text{pre}(\text{the}(\text{com}(Ts ! i)))));$
 $\llbracket - (\bigcap i \in \{i. i < \text{length } Ts\}. \text{pre}(\text{the}(\text{com}(Ts ! i))))$
 (Parallel Ts)
 $(\bigcap i \in \{i. i < \text{length } Ts\}. \text{post}(Ts ! i));$
 $(\bigcap i \in \{i. i < \text{length } Ts\}. \text{post}(Ts ! i)) \subseteq q \rrbracket$
 $\implies \llbracket - p \text{ (Parallel Ts) } q$
apply *(rule Conseq)*
prefer 2
 apply *fast*
apply *assumption+*
done

lemma *SkipRule*: $p \subseteq q \implies \llbracket - p \text{ (Basic id) } q$

apply*(rule oghoare-intros)*
 prefer 2 **apply***(rule Basic)*
 prefer 2 **apply***(rule subset-refl)*
apply*(simp add:Id-def)*
done

lemma *BasicRule*: $p \subseteq \{s. (f s) \in q\} \implies \llbracket - p \text{ (Basic f) } q$

apply*(rule oghoare-intros)*
 prefer 2 **apply***(rule oghoare-intros)*
 prefer 2 **apply***(rule subset-refl)*
apply *assumption*
done

lemma *SeqRule*: $\llbracket \llbracket - p \text{ } c1 \text{ } r; \llbracket - r \text{ } c2 \text{ } q \rrbracket \implies \llbracket - p \text{ (Seq } c1 \text{ } c2) \text{ } q$

apply*(rule Seq)*
apply *fast+*
done

lemma *CondRule*:

$\llbracket p \subseteq \{s. (s \in b \implies s \in w) \wedge (s \notin b \implies s \in w')\}; \llbracket - w \text{ } c1 \text{ } q; \llbracket - w' \text{ } c2 \text{ } q \rrbracket$
 $\implies \llbracket - p \text{ (Cond } b \text{ } c1 \text{ } c2) \text{ } q$
apply*(rule Cond)*
 apply*(rule Conseq)*
 prefer 4 **apply***(rule Conseq)*
apply *simp-all*
apply *force+*
done

lemma *WhileRule*: $\llbracket p \subseteq i; \llbracket - (i \cap b) \text{ } c \text{ } i; (i \cap (-b)) \subseteq q \rrbracket$

$\implies \llbracket - p \text{ (While } b \text{ } i \text{ } c) \text{ } q$
apply*(rule Conseq)*
 prefer 2 **apply***(rule While)*
apply *assumption+*
done

Three new proof rules for special instances of the *AnnBasic* and the *AnnAwait* commands when the transformation performed on the state is the identity, and for an *AnnAwait* command where the boolean condition is $\{s. True\}$:

lemma *AnnatomRule*:

```

[[ atom-com(c); ||- r c q ]] ==> ⊢ (AnnAwait r {s. True} c) q
apply(rule AnnAwait)
apply simp-all
done

```

lemma *AnnskipRule*:

```

r ⊆ q ==> ⊢ (AnnBasic r id) q
apply(rule AnnBasic)
apply simp
done

```

lemma *AnnwaitRule*:

```

[[ (r ∩ b) ⊆ q ]] ==> ⊢ (AnnAwait r b (Basic id)) q
apply(rule AnnAwait)
apply simp
apply(rule BasicRule)
apply simp
done

```

Lemmata to avoid using the definition of *map-ann-hoare*, *interfree-aux*, *interfree-swap* and *interfree* by splitting it into different cases:

lemma *interfree-aux-rule1*: *interfree-aux*(*co*, *q*, *None*)
by(*simp add:interfree-aux-def*)

lemma *interfree-aux-rule2*:

```

∀(R,r)∈(atomics a). ||- (q ∩ R) r q ==> interfree-aux(None, q, Some a)
apply(simp add:interfree-aux-def)
apply(force elim:oghoare-sound)
done

```

lemma *interfree-aux-rule3*:

```

(∀(R, r)∈(atomics a). ||- (q ∩ R) r q ∧ (∀p∈(assertions c). ||- (p ∩ R) r p))
==> interfree-aux(Some c, q, Some a)
apply(simp add:interfree-aux-def)
apply(force elim:oghoare-sound)
done

```

lemma *AnnBasic-assertions*:

```

[[interfree-aux(None, r, Some a); interfree-aux(None, q, Some a)]] ==>
interfree-aux(Some (AnnBasic r f), q, Some a)
apply(simp add: interfree-aux-def)
by force

```


lemma *AnnSeq-assertions*:
 $\llbracket \text{interfree-aux}(\text{Some } c1, q, \text{Some } a); \text{interfree-aux}(\text{Some } c2, q, \text{Some } a) \rrbracket \Longrightarrow$
 $\text{interfree-aux}(\text{Some } (\text{AnnSeq } c1 \ c2), q, \text{Some } a)$
apply(*simp add: interfree-aux-def*)
by *force*

lemma *AnnCond1-assertions*:
 $\llbracket \text{interfree-aux}(\text{None}, r, \text{Some } a); \text{interfree-aux}(\text{Some } c1, q, \text{Some } a);$
 $\text{interfree-aux}(\text{Some } c2, q, \text{Some } a) \rrbracket \Longrightarrow$
 $\text{interfree-aux}(\text{Some } (\text{AnnCond1 } r \ b \ c1 \ c2), q, \text{Some } a)$
apply(*simp add: interfree-aux-def*)
by *force*

lemma *AnnCond2-assertions*:
 $\llbracket \text{interfree-aux}(\text{None}, r, \text{Some } a); \text{interfree-aux}(\text{Some } c, q, \text{Some } a) \rrbracket \Longrightarrow$
 $\text{interfree-aux}(\text{Some } (\text{AnnCond2 } r \ b \ c), q, \text{Some } a)$
apply(*simp add: interfree-aux-def*)
by *force*

lemma *AnnWhile-assertions*:
 $\llbracket \text{interfree-aux}(\text{None}, r, \text{Some } a); \text{interfree-aux}(\text{None}, i, \text{Some } a);$
 $\text{interfree-aux}(\text{Some } c, q, \text{Some } a) \rrbracket \Longrightarrow$
 $\text{interfree-aux}(\text{Some } (\text{AnnWhile } r \ b \ i \ c), q, \text{Some } a)$
apply(*simp add: interfree-aux-def*)
by *force*

lemma *AnnAwait-assertions*:
 $\llbracket \text{interfree-aux}(\text{None}, r, \text{Some } a); \text{interfree-aux}(\text{None}, q, \text{Some } a) \rrbracket \Longrightarrow$
 $\text{interfree-aux}(\text{Some } (\text{AnnAwait } r \ b \ c), q, \text{Some } a)$
apply(*simp add: interfree-aux-def*)
by *force*

lemma *AnnBasic-atomics*:
 $\llbracket - (q \cap r) \ (\text{Basic } f) \ q \Longrightarrow \text{interfree-aux}(\text{None}, q, \text{Some } (\text{AnnBasic } r \ f)) \rrbracket$
by(*simp add: interfree-aux-def oghoare-sound*)

lemma *AnnSeq-atomics*:
 $\llbracket \text{interfree-aux}(\text{Any}, q, \text{Some } a1); \text{interfree-aux}(\text{Any}, q, \text{Some } a2) \rrbracket \Longrightarrow$
 $\text{interfree-aux}(\text{Any}, q, \text{Some } (\text{AnnSeq } a1 \ a2))$
apply(*simp add: interfree-aux-def*)
by *force*

lemma *AnnCond1-atomics*:
 $\llbracket \text{interfree-aux}(\text{Any}, q, \text{Some } a1); \text{interfree-aux}(\text{Any}, q, \text{Some } a2) \rrbracket \Longrightarrow$
 $\text{interfree-aux}(\text{Any}, q, \text{Some } (\text{AnnCond1 } r \ b \ a1 \ a2))$
apply(*simp add: interfree-aux-def*)
by *force*

lemma *AnnCond2-atomics*:

$interfree\text{-}aux\ (Any, q, Some\ a) \implies interfree\text{-}aux\ (Any, q, Some\ (AnnCond2\ r\ b\ a))$

by(simp add: interfree-aux-def)

lemma *AnnWhile-atomics*: $interfree\text{-}aux\ (Any, q, Some\ a)$

$\implies interfree\text{-}aux\ (Any, q, Some\ (AnnWhile\ r\ b\ i\ a))$

by(simp add: interfree-aux-def)

lemma *Annatom-atomics*:

$\llbracket - \ (q \cap r)\ a\ q \implies interfree\text{-}aux\ (None, q, Some\ (AnnAwait\ r\ \{x.\ True\}\ a))$

by(simp add: interfree-aux-def oghoare-sound)

lemma *AnnAwait-atomics*:

$\llbracket - \ (q \cap (r \cap b))\ a\ q \implies interfree\text{-}aux\ (None, q, Some\ (AnnAwait\ r\ b\ a))$

by(simp add: interfree-aux-def oghoare-sound)

definition *interfree-swap* :: $('a\ ann\text{-}triple\text{-}op * ('a\ ann\text{-}triple\text{-}op)\ list) \Rightarrow bool$ **where**

$interfree\text{-}swap == \lambda(x, xs). \forall y \in set\ xs. interfree\text{-}aux\ (com\ x, post\ x, com\ y)$

$\wedge interfree\text{-}aux\ (com\ y, post\ y, com\ x)$

lemma *interfree-swap-Empty*: $interfree\text{-}swap\ (x, [])$

by(simp add: interfree-swap-def)

lemma *interfree-swap-List*:

$\llbracket interfree\text{-}aux\ (com\ x, post\ x, com\ y);$

$interfree\text{-}aux\ (com\ y, post\ y, com\ x); interfree\text{-}swap\ (x, xs) \rrbracket$

$\implies interfree\text{-}swap\ (x, y\#\!xs)$

by(simp add: interfree-swap-def)

lemma *interfree-swap-Map*: $\forall k. i \leq k \wedge k < j \longrightarrow interfree\text{-}aux\ (com\ x, post\ x, c\ k)$

$\wedge interfree\text{-}aux\ (c\ k, Q\ k, com\ x)$

$\implies interfree\text{-}swap\ (x, map\ (\lambda k. (c\ k, Q\ k))\ [i..<j])$

by(force simp add: interfree-swap-def less-diff-conv)

lemma *interfree-Empty*: $interfree\ []$

by(simp add: interfree-def)

lemma *interfree-List*:

$\llbracket interfree\text{-}swap(x, xs); interfree\ xs \rrbracket \implies interfree\ (x\#\!xs)$

apply(simp add: interfree-def interfree-swap-def)

apply *clarify*

apply(case-tac *i*)

apply(case-tac *j*)

apply *simp-all*

apply(case-tac *j*, simp+)

done

lemma *interfree-Map*:

$(\forall i\ j. a \leq i \wedge i < b \wedge a \leq j \wedge j < b \wedge i \neq j \longrightarrow interfree\text{-}aux\ (c\ i, Q\ i, c\ j))$

$\implies \text{interfree } (\text{map } (\lambda k. (c\ k, Q\ k)) [a..<b])$
by(*force simp add: interfree-def less-diff-conv*)

definition *map-ann-hoare* :: (('a ann-com-op * 'a assn) list) \Rightarrow bool ($\langle [\vdash] \rightarrow [0]$
45) **where**
 $[\vdash] Ts == (\forall i < \text{length } Ts. \exists c\ q. Ts!i = (\text{Some } c, q) \wedge \vdash c\ q)$

lemma *MapAnnEmpty*: $[\vdash] []$
by(*simp add: map-ann-hoare-def*)

lemma *MapAnnList*: $[[\vdash c\ q ; [\vdash] xs]] \implies [\vdash] (\text{Some } c, q) \# xs$
apply(*simp add: map-ann-hoare-def*)
apply *clarify*
apply(*case-tac i, simp+*)
done

lemma *MapAnnMap*:
 $\forall k. i < k \wedge k < j \longrightarrow \vdash (c\ k) (Q\ k) \implies [\vdash] \text{map } (\lambda k. (\text{Some } (c\ k), Q\ k)) [i..<j]$
apply(*simp add: map-ann-hoare-def less-diff-conv*)
done

lemma *ParallelRule*: $[[[\vdash] Ts ; \text{interfree } Ts]]$
 $\implies \|\vdash - (\bigcap i \in \{i. i < \text{length } Ts\}. \text{pre}(\text{the}(\text{com}(Ts!i))))$
Parallel Ts
 $(\bigcap i \in \{i. i < \text{length } Ts\}. \text{post}(Ts!i))$
apply(*rule Parallel*)
apply(*simp add: map-ann-hoare-def*)
apply *simp*
done

The following are some useful lemmas and simplification tactics to control which theorems are used to simplify at each moment, so that the original input does not suffer any unexpected transformation.

lemma *Compl-Collect*: $\neg(\text{Collect } b) = \{x. \neg(b\ x)\}$
by *fast*

lemma *list-length*: $\text{length } [] = 0$ $\text{length } (x \# xs) = \text{Suc}(\text{length } xs)$
by *simp-all*

lemma *list-lemmas*: $\text{length } [] = 0$ $\text{length } (x \# xs) = \text{Suc}(\text{length } xs)$
 $(x \# xs) ! 0 = x$ $(x \# xs) ! \text{Suc } n = xs ! n$
by *simp-all*

lemma *le-Suc-eq-insert*: $\{i. i < \text{Suc } n\} = \text{insert } n \{i. i < n\}$
by *auto*

lemmas *primrecdef-list* = *pre.simps assertions.simps atomics.simps atom-com.simps*

lemmas *my-simp-list* = *list-lemmas fst-conv snd-conv*
not-less0 refl le-Suc-eq-insert Suc-not-Zero Zero-not-Suc nat.inject
Collect-mem-eq ball-simps option.simps primrecdef-list

lemmas *ParallelConseq-list* = *INTER-eq Collect-conj-eq length-map length-upt length-append*

```

ML <
fun before-interfree-simp-tac ctxt =
  simp-tac (put-simpset HOL-basic-ss ctxt addsimps [@{thm com.simps}, @{thm
post.simps}])

fun interfree-simp-tac ctxt =
  asm-simp-tac (put-simpset HOL-ss ctxt
  addsimps [@{thm split}, @{thm ball-Un}, @{thm ball-empty}] @ @{thms my-simp-list})

fun ParallelConseq ctxt =
  simp-tac (put-simpset HOL-basic-ss ctxt
  addsimps @@{thms ParallelConseq-list} @ @{thms my-simp-list})
>

```

The following tactic applies *tac* to each conjunct in a subgoal of the form $A \implies a1 \wedge a2 \wedge \dots \wedge an$ returning n subgoals, one for each conjunct:

```

ML <
fun conjI-Tac ctxt tac i st = st |>
  ( (EVERY [resolve-tac ctxt [conjI] i,
  conjI-Tac ctxt tac (i+1),
  tac i]) ORELSE (tac i) )
>

```

Tactic for the generation of the verification conditions

The tactic basically uses two subtactics:

HoareRuleTac is called at the level of parallel programs, it uses the **ParallelTac** to solve parallel composition of programs. This verification has two parts, namely, (1) all component programs are correct and (2) they are interference free. *HoareRuleTac* is also called at the level of atomic regions, i.e. $\langle \rangle$ and *AWAIT b THEN - END*, and at each interference freedom test.

AnnHoareRuleTac is for component programs which are annotated programs and so, there are not unknown assertions (no need to use the parameter *precond*, see NOTE).

NOTE: *precond*(::bool) informs if the subgoal has the form $\| - ?p \ c \ q$, in this case we have *precond*=False and the generated verification condition would have the form $?p \subseteq \dots$ which can be solved by *rtac subset-refl*, if True we proceed to simplify it using the simplification tactics above.

```

ML <
fun WlpTac ctxt i = resolve-tac ctxt @@{thms SeqRule} i THEN HoareRuleTac ctxt
false (i + 1)

```

and HoareRuleTac ctxt precondition i st = st |>
 ((WlpTac ctxt i THEN HoareRuleTac ctxt precondition i)
 ORELSE
 (FIRST[resolve-tac ctxt @{thms SkipRule} i,
 resolve-tac ctxt @{thms BasicRule} i,
 EVERY[resolve-tac ctxt @{thms ParallelConseqRule} i,
 ParallelConseq ctxt (i+2),
 ParallelTac ctxt (i+1),
 ParallelConseq ctxt i],
 EVERY[resolve-tac ctxt @{thms CondRule} i,
 HoareRuleTac ctxt false (i+2),
 HoareRuleTac ctxt false (i+1)],
 EVERY[resolve-tac ctxt @{thms WhileRule} i,
 HoareRuleTac ctxt true (i+1)],
 K all-tac i]
 THEN (if precondition then (K all-tac i) else resolve-tac ctxt @{thms subset-refl}
 i)))

and AnnWlpTac ctxt i = resolve-tac ctxt @{thms AnnSeq} i THEN AnnHoareRule-
Tac ctxt (i + 1)

and AnnHoareRuleTac ctxt i st = st |>
 ((AnnWlpTac ctxt i THEN AnnHoareRuleTac ctxt i)
 ORELSE
 (FIRST[(resolve-tac ctxt @{thms AnnskipRule} i),
 EVERY[resolve-tac ctxt @{thms AnnatomRule} i,
 HoareRuleTac ctxt true (i+1)],
 (resolve-tac ctxt @{thms AnnwaitRule} i),
 resolve-tac ctxt @{thms AnnBasic} i,
 EVERY[resolve-tac ctxt @{thms AnnCond1} i,
 AnnHoareRuleTac ctxt (i+3),
 AnnHoareRuleTac ctxt (i+1)],
 EVERY[resolve-tac ctxt @{thms AnnCond2} i,
 AnnHoareRuleTac ctxt (i+1)],
 EVERY[resolve-tac ctxt @{thms AnnWhile} i,
 AnnHoareRuleTac ctxt (i+2)],
 EVERY[resolve-tac ctxt @{thms AnnAwait} i,
 HoareRuleTac ctxt true (i+1)],
 K all-tac i]))

*and ParallelTac ctxt i = EVERY[resolve-tac ctxt @{thms ParallelRule} i,
 interfree-Tac ctxt (i+1),
 MapAnn-Tac ctxt i]*

and MapAnn-Tac ctxt i st = st |>
 (FIRST[resolve-tac ctxt @{thms MapAnnEmpty} i,
 EVERY[resolve-tac ctxt @{thms MapAnnList} i,
 MapAnn-Tac ctxt (i+1),
 AnnHoareRuleTac ctxt i],
 EVERY[resolve-tac ctxt @{thms MapAnnMap} i,

resolve-tac ctxt @{*thms allI*} *i*,
resolve-tac ctxt @{*thms impI*} *i*,
AnnHoareRuleTac ctxt i])

and interfree-swap-Tac ctxt i st = st |>
 (FIRST[*resolve-tac ctxt* @{*thms interfree-swap-Empty*} *i*,
 EVERY[*resolve-tac ctxt* @{*thms interfree-swap-List*} *i*,
interfree-swap-Tac ctxt (i+2),
interfree-aux-Tac ctxt (i+1),
interfree-aux-Tac ctxt i],
 EVERY[*resolve-tac ctxt* @{*thms interfree-swap-Map*} *i*,
resolve-tac ctxt @{*thms allI*} *i*,
resolve-tac ctxt @{*thms impI*} *i*,
conjI-Tac ctxt (interfree-aux-Tac ctxt i)]])

and interfree-Tac ctxt i st = st |>
 (FIRST[*resolve-tac ctxt* @{*thms interfree-Empty*} *i*,
 EVERY[*resolve-tac ctxt* @{*thms interfree-List*} *i*,
interfree-Tac ctxt (i+1),
interfree-swap-Tac ctxt i],
 EVERY[*resolve-tac ctxt* @{*thms interfree-Map*} *i*,
resolve-tac ctxt @{*thms allI*} *i*,
resolve-tac ctxt @{*thms allI*} *i*,
resolve-tac ctxt @{*thms impI*} *i*,
interfree-aux-Tac ctxt i]])

and interfree-aux-Tac ctxt i = (before-interfree-simp-tac ctxt i) THEN
 (FIRST[*resolve-tac ctxt* @{*thms interfree-aux-rule1*} *i*,
dest-assertions-Tac ctxt i])

and dest-assertions-Tac ctxt i st = st |>
 (FIRST[EVERY[*resolve-tac ctxt* @{*thms AnnBasic-assertions*} *i*,
dest-atomics-Tac ctxt (i+1),
dest-atomics-Tac ctxt i],
 EVERY[*resolve-tac ctxt* @{*thms AnnSeq-assertions*} *i*,
dest-assertions-Tac ctxt (i+1),
dest-assertions-Tac ctxt i],
 EVERY[*resolve-tac ctxt* @{*thms AnnCond1-assertions*} *i*,
dest-assertions-Tac ctxt (i+2),
dest-assertions-Tac ctxt (i+1),
dest-atomics-Tac ctxt i],
 EVERY[*resolve-tac ctxt* @{*thms AnnCond2-assertions*} *i*,
dest-assertions-Tac ctxt (i+1),
dest-atomics-Tac ctxt i],
 EVERY[*resolve-tac ctxt* @{*thms AnnWhile-assertions*} *i*,
dest-assertions-Tac ctxt (i+2),
dest-atomics-Tac ctxt (i+1),
dest-atomics-Tac ctxt i],
 EVERY[*resolve-tac ctxt* @{*thms AnnAwait-assertions*} *i*,

```

      dest-atomics-Tac ctxt (i+1),
      dest-atomics-Tac ctxt i],
    dest-atomics-Tac ctxt i])

and dest-atomics-Tac ctxt i st = st |>
  (FIRST[EVERY[resolve-tac ctxt @{thms AnnBasic-atomics} i,
    HoareRuleTac ctxt true i],
    EVERY[resolve-tac ctxt @{thms AnnSeq-atomics} i,
      dest-atomics-Tac ctxt (i+1),
      dest-atomics-Tac ctxt i],
    EVERY[resolve-tac ctxt @{thms AnnCond1-atomics} i,
      dest-atomics-Tac ctxt (i+1),
      dest-atomics-Tac ctxt i],
    EVERY[resolve-tac ctxt @{thms AnnCond2-atomics} i,
      dest-atomics-Tac ctxt i],
    EVERY[resolve-tac ctxt @{thms AnnWhile-atomics} i,
      dest-atomics-Tac ctxt i],
    EVERY[resolve-tac ctxt @{thms Annatom-atomics} i,
      HoareRuleTac ctxt true i],
    EVERY[resolve-tac ctxt @{thms AnnAwait-atomics} i,
      HoareRuleTac ctxt true i],
    K all-tac i])
  >

```

The final tactic is given the name *oghoare*:

```

ML <
fun oghoare-tac ctxt = SUBGOAL (fn (-, i) => HoareRuleTac ctxt true i)
  >

```

Notice that the tactic for parallel programs *oghoare-tac* is initially invoked with the value *true* for the parameter *precond*.

Parts of the tactic can be also individually used to generate the verification conditions for annotated sequential programs and to generate verification conditions out of interference freedom tests:

```

ML <
fun annhoare-tac ctxt = SUBGOAL (fn (-, i) => AnnHoareRuleTac ctxt i)

fun interfree-aux-tac ctxt = SUBGOAL (fn (-, i) => interfree-aux-Tac ctxt i)
  >

```

The so defined ML tactics are then “exported” to be used in Isabelle proofs.

```

method-setup oghoare = <
  Scan.succeed (SIMPLE-METHOD' o oghoare-tac)>
  verification condition generator for the oghoare logic

```

```

method-setup annhoare = <
  Scan.succeed (SIMPLE-METHOD' o annhoare-tac)>
  verification condition generator for the ann-hoare logic

```

```

method-setup interfree-aux = ⟨
  Scan.succeed (SIMPLE-METHOD' o interfree-aux-tac)⟩
  verification condition generator for interference freedom tests

```

Tactics useful for dealing with the generated verification conditions:

```

method-setup conjI-tac = ⟨
  Scan.succeed (fn ctxt => SIMPLE-METHOD' (conjI-Tac ctxt (K all-tac)))⟩
  verification condition generator for interference freedom tests

```

```

ML ⟨
  fun disjE-Tac ctxt tac i st = st |>
    ( (EVERY [eresolve-tac ctxt [disjE] i,
           disjE-Tac ctxt tac (i+1),
           tac i] ) ORELSE (tac i) )
  >

```

```

method-setup disjE-tac = ⟨
  Scan.succeed (fn ctxt => SIMPLE-METHOD' (disjE-Tac ctxt (K all-tac)))⟩
  verification condition generator for interference freedom tests

```

end

1.7 Concrete Syntax

```

theory Quote-Antiquote imports Main begin

```

```

syntax
  -quote    :: 'b ⇒ ('a ⇒ 'b)           (⟨«-»⟩ [0] 1000)
  -antiquote :: ('a ⇒ 'b) ⇒ 'b          (⟨'-> [1000] 1000)
  -Assert   :: 'a ⇒ 'a set              (⟨{!-!}⟩ [0] 1000)

```

```

translations
  {!b} → CONST Collect «b»

```

```

parse-translation ⟨
  let
    fun quote-tr [t] = Syntax-Trans.quote-tr syntax-const ⟨-antiquote⟩ t
      | quote-tr ts = raise TERM (quote-tr, ts);
  in [(syntax-const ⟨-quote⟩, K quote-tr)] end
  >

```

```

end
theory OG-Syntax
imports OG-Tactics Quote-Antiquote
begin

```

Syntax for commands and for assertions and boolean expressions in commands *com* and annotated commands *ann-com*.

abbreviation $Skip :: 'a\ com\ (\langle SKIP \rangle\ 63)$
where $SKIP \equiv Basic\ id$

abbreviation $AnnSkip :: 'a\ assn \Rightarrow 'a\ ann-com\ (\langle -//SKIP \rangle\ [90]\ 63)$
where $r\ SKIP \equiv AnnBasic\ r\ id$

notation

$Seq\ (\langle -, / \rightarrow [55, 56]\ 55)$ **and**
 $AnnSeq\ (\langle -; / \rightarrow [60, 61]\ 60)$

syntax

$-Assign :: idt \Rightarrow 'b \Rightarrow 'a\ com\ (\langle (- := / -) \rangle\ [70, 65]\ 61)$
 $-AnnAssign :: 'a\ assn \Rightarrow idt \Rightarrow 'b \Rightarrow 'a\ com\ (\langle (- ' := / -) \rangle\ [90, 70, 65]\ 61)$

translations

$'x := a \rightarrow CONST\ Basic\ \langle\langle (-update-name\ x\ (\lambda-. a)) \rangle\rangle$
 $r\ 'x := a \rightarrow CONST\ AnnBasic\ r\ \langle\langle (-update-name\ x\ (\lambda-. a)) \rangle\rangle$

syntax

$-AnnCond1 :: 'a\ assn \Rightarrow 'a\ bexp \Rightarrow 'a\ ann-com \Rightarrow 'a\ ann-com \Rightarrow 'a\ ann-com$
 $(\langle - // IF - / THEN - / ELSE - / FI \rangle\ [90, 0, 0, 0]\ 61)$
 $-AnnCond2 :: 'a\ assn \Rightarrow 'a\ bexp \Rightarrow 'a\ ann-com \Rightarrow 'a\ ann-com$
 $(\langle - // IF - / THEN - / FI \rangle\ [90, 0, 0]\ 61)$
 $-AnnWhile :: 'a\ assn \Rightarrow 'a\ bexp \Rightarrow 'a\ assn \Rightarrow 'a\ ann-com \Rightarrow 'a\ ann-com$
 $(\langle - // WHILE - / INV - // DO - // OD \rangle\ [90, 0, 0, 0]\ 61)$
 $-AnnAwait :: 'a\ assn \Rightarrow 'a\ bexp \Rightarrow 'a\ com \Rightarrow 'a\ ann-com$
 $(\langle - // AWAIT - / THEN - / END \rangle\ [90, 0, 0]\ 61)$
 $-AnnAtom :: 'a\ assn \Rightarrow 'a\ com \Rightarrow 'a\ ann-com\ (\langle - // \langle - \rangle \rangle\ [90, 0]\ 61)$
 $-AnnWait :: 'a\ assn \Rightarrow 'a\ bexp \Rightarrow 'a\ ann-com\ (\langle - // WAIT - END \rangle\ [90, 0]\ 61)$

$-Cond :: 'a\ bexp \Rightarrow 'a\ com \Rightarrow 'a\ com \Rightarrow 'a\ com$
 $(\langle (0IF - / THEN - / ELSE - / FI) \rangle\ [0, 0, 0]\ 61)$
 $-Cond2 :: 'a\ bexp \Rightarrow 'a\ com \Rightarrow 'a\ com\ (\langle IF - THEN - FI \rangle\ [0, 0]\ 56)$
 $-While-inv :: 'a\ bexp \Rightarrow 'a\ assn \Rightarrow 'a\ com \Rightarrow 'a\ com$
 $(\langle (0WHILE - / INV - // DO - / OD) \rangle\ [0, 0, 0]\ 61)$
 $-While :: 'a\ bexp \Rightarrow 'a\ com \Rightarrow 'a\ com$
 $(\langle (0WHILE - // DO - / OD) \rangle\ [0, 0]\ 61)$

translations

$IF\ b\ THEN\ c1\ ELSE\ c2\ FI \rightarrow CONST\ Cond\ \{\!\{b\}\!\} c1\ c2$
 $IF\ b\ THEN\ c\ FI \Leftrightarrow IF\ b\ THEN\ c\ ELSE\ SKIP\ FI$
 $WHILE\ b\ INV\ i\ DO\ c\ OD \rightarrow CONST\ While\ \{\!\{b\}\!\} i\ c$
 $WHILE\ b\ DO\ c\ OD \Leftrightarrow WHILE\ b\ INV\ CONST\ undefined\ DO\ c\ OD$

$r\ IF\ b\ THEN\ c1\ ELSE\ c2\ FI \rightarrow CONST\ AnnCond1\ r\ \{\!\{b\}\!\} c1\ c2$
 $r\ IF\ b\ THEN\ c\ FI \rightarrow CONST\ AnnCond2\ r\ \{\!\{b\}\!\} c$
 $r\ WHILE\ b\ INV\ i\ DO\ c\ OD \rightarrow CONST\ AnnWhile\ r\ \{\!\{b\}\!\} i\ c$
 $r\ AWAIT\ b\ THEN\ c\ END \rightarrow CONST\ AnnAwait\ r\ \{\!\{b\}\!\} c$
 $r\ \langle c \rangle \Leftrightarrow r\ AWAIT\ CONST\ True\ THEN\ c\ END$

$r \text{ WAIT } b \text{ END} \rightleftharpoons r \text{ AWAIT } b \text{ THEN SKIP END}$

nonterminal *prgs*

syntax

$\text{-PAR} :: \text{prgs} \Rightarrow 'a \quad (\langle \text{COBEGIN} // - // \text{COEND} \rangle [57] 56)$

$\text{-prg} :: ['a, 'a] \Rightarrow \text{prgs} \quad (\langle - // - \rangle [60, 90] 57)$

$\text{-prgs} :: ['a, 'a, \text{prgs}] \Rightarrow \text{prgs} \quad (\langle - // - // || // - \rangle [60, 90, 57] 57)$

$\text{-prg-scheme} :: ['a, 'a, 'a, 'a, 'a] \Rightarrow \text{prgs}$
 $(\langle \text{SCHEME} [- \leq - < -] - // - \rangle [0, 0, 0, 60, 90] 57)$

translations

$\text{-prg } c \ q \rightleftharpoons [(\text{CONST } \text{Some } c, q)]$

$\text{-prgs } c \ q \ ps \rightleftharpoons (\text{CONST } \text{Some } c, q) \# ps$

$\text{-PAR } ps \rightleftharpoons \text{CONST Parallel } ps$

$\text{-prg-scheme } j \ i \ k \ c \ q \rightleftharpoons \text{CONST map } (\lambda i. (\text{CONST } \text{Some } c, q)) [j..<k]$

print-translation \langle

let

$\text{fun quote-tr}' f (t :: ts) =$

$\text{Term.list-comb } (f \ \$ \ \text{Syntax-Trans.quote-tr}' \ \mathbf{syntax-const} \ \langle \text{-antiquote} \rangle \ t,$

$ts)$

$| \ \text{quote-tr}' \ - \ - = \text{raise Match};$

$\text{fun annquote-tr}' f (r :: t :: ts) =$

$\text{Term.list-comb } (f \ \$ \ r \ \$ \ \text{Syntax-Trans.quote-tr}' \ \mathbf{syntax-const} \ \langle \text{-antiquote} \rangle$

$t, ts)$

$| \ \text{annquote-tr}' \ - \ - = \text{raise Match};$

$\text{val assert-tr}' = \text{quote-tr}' (\text{Syntax.const } \mathbf{syntax-const} \ \langle \text{-Assert} \rangle);$

$\text{fun bexp-tr}' \ \text{name} ((\text{Const } (\mathbf{const-syntax} \ \langle \text{Collect} \rangle, -) \ \$ \ t) :: ts) =$

$\text{quote-tr}' (\text{Syntax.const } \text{name}) (t :: ts)$

$| \ \text{bexp-tr}' \ - \ - = \text{raise Match};$

$\text{fun annbexp-tr}' \ \text{name} (r :: (\text{Const } (\mathbf{const-syntax} \ \langle \text{Collect} \rangle, -) \ \$ \ t) :: ts) =$

$\text{annquote-tr}' (\text{Syntax.const } \text{name}) (r :: t :: ts)$

$| \ \text{annbexp-tr}' \ - \ - = \text{raise Match};$

$\text{fun assign-tr}' (\text{Abs } (x, -, f \ \$ \ k \ \$ \ \text{Bound } 0) :: ts) =$

$\text{quote-tr}' (\text{Syntax.const } \mathbf{syntax-const} \ \langle \text{-Assign} \rangle \ \$ \ \text{Syntax-Trans.update-name-tr}'$

$f)$

$(\text{Abs } (x, \text{dummyT}, \text{Syntax-Trans.const-abs-tr}' \ k) :: ts)$

$| \ \text{assign-tr}' \ - \ - = \text{raise Match};$

$\text{fun annassign-tr}' (r :: \text{Abs } (x, -, f \ \$ \ k \ \$ \ \text{Bound } 0) :: ts) =$

$\text{quote-tr}' (\text{Syntax.const } \mathbf{syntax-const} \ \langle \text{-AnnAssign} \rangle \ \$ \ r \ \$ \ \text{Syntax-Trans.update-name-tr}'$

```

f)
  (Abs (x, dummyT, Syntax-Trans.const-abs-tr' k) :: ts)
  | annassign-tr' - = raise Match;

  fun Parallel-PAR [(Const (const-syntax⟨Cons⟩, -) $
    (Const (const-syntax⟨Pair⟩, -) $ (Const (const-syntax⟨Some⟩, -) $ t1
  ) $ t2) $
    Const (const-syntax⟨Nil⟩, -))] = Syntax.const syntax-const⟨-prg⟩ $
t1 $ t2
  | Parallel-PAR [(Const (const-syntax⟨Cons⟩, -) $
    (Const (const-syntax⟨Pair⟩, -) $ (Const (const-syntax⟨Some⟩, -) $
t1) $ t2) $ ts)] =
    Syntax.const syntax-const⟨-prgs⟩ $ t1 $ t2 $ Parallel-PAR [ts]
  | Parallel-PAR - = raise Match;

  fun Parallel-tr' ts = Syntax.const syntax-const⟨-PAR⟩ $ Parallel-PAR ts;
in
  [(const-syntax⟨Collect⟩, K assert-tr'),
   (const-syntax⟨Basic⟩, K assign-tr'),
   (const-syntax⟨Cond⟩, K (bexp-tr' syntax-const⟨-Cond⟩)),
   (const-syntax⟨While⟩, K (bexp-tr' syntax-const⟨-While-inv⟩)),
   (const-syntax⟨AnnBasic⟩, K annassign-tr'),
   (const-syntax⟨AnnWhile⟩, K (annbexp-tr' syntax-const⟨-AnnWhile⟩)),
   (const-syntax⟨AnnAwait⟩, K (annbexp-tr' syntax-const⟨-AnnAwait⟩)),
   (const-syntax⟨AnnCond1⟩, K (annbexp-tr' syntax-const⟨-AnnCond1⟩)),
   (const-syntax⟨AnnCond2⟩, K (annbexp-tr' syntax-const⟨-AnnCond2⟩))]
end

```

1.8 Examples

theory *OG-Examples* **imports** *OG-Syntax* **begin**

1.8.1 Mutual Exclusion

Peterson's Algorithm I

Eike Best. "Semantics of Sequential and Parallel Programs", page 217.

record *Petersons-mutex-1* =

```

pr1 :: nat
pr2 :: nat
in1 :: bool
in2 :: bool
hold :: nat

```

lemma *Petersons-mutex-1*:

```

||- { 'pr1=0 ∧ ¬ in1 ∧ 'pr2=0 ∧ ¬ in2 }

```

```

COBEGIN { 'pr1=0 ∧ ¬ 'in1 }
WHILE True INV { 'pr1=0 ∧ ¬ 'in1 }
DO
  { 'pr1=0 ∧ ¬ 'in1 } ⟨ 'in1:=True,, 'pr1:=1 ⟩;;
  { 'pr1=1 ∧ 'in1 } ⟨ 'hold:=1,, 'pr1:=2 ⟩;;
  { 'pr1=2 ∧ 'in1 ∧ ('hold=1 ∨ 'hold=2 ∧ 'pr2=2) }
  AWAIT (¬ 'in2 ∨ ¬('hold=1)) THEN 'pr1:=3 END;;
  { 'pr1=3 ∧ 'in1 ∧ ('hold=1 ∨ 'hold=2 ∧ 'pr2=2) }
  ⟨ 'in1:=False,, 'pr1:=0 ⟩
OD { 'pr1=0 ∧ ¬ 'in1 }
||
{ 'pr2=0 ∧ ¬ 'in2 }
WHILE True INV { 'pr2=0 ∧ ¬ 'in2 }
DO
  { 'pr2=0 ∧ ¬ 'in2 } ⟨ 'in2:=True,, 'pr2:=1 ⟩;;
  { 'pr2=1 ∧ 'in2 } ⟨ 'hold:=2,, 'pr2:=2 ⟩;;
  { 'pr2=2 ∧ 'in2 ∧ ('hold=2 ∨ ('hold=1 ∧ 'pr1=2)) }
  AWAIT (¬ 'in1 ∨ ¬('hold=2)) THEN 'pr2:=3 END;;
  { 'pr2=3 ∧ 'in2 ∧ ('hold=2 ∨ ('hold=1 ∧ 'pr1=2)) }
  ⟨ 'in2:=False,, 'pr2:=0 ⟩
OD { 'pr2=0 ∧ ¬ 'in2 }
COEND
{ 'pr1=0 ∧ ¬ 'in1 ∧ 'pr2=0 ∧ ¬ 'in2 }

```

apply *oghoare*

— 104 verification conditions.

apply *auto*

done

Peterson's Algorithm II: A Busy Wait Solution

Apt and Olderog. "Verification of sequential and concurrent Programs",
page 282.

record *Busy-wait-mutex* =

flag1 :: bool

flag2 :: bool

turn :: nat

after1 :: bool

after2 :: bool

lemma *Busy-wait-mutex*:

||- { True }

'flag1:=False,, 'flag2:=False,,

COBEGIN { ¬ 'flag1 }

WHILE True

INV { ¬ 'flag1 }

DO { ¬ 'flag1 } ⟨ 'flag1:=True,, 'after1:=False ⟩;;

{ 'flag1 ∧ ¬ 'after1 } ⟨ 'turn:=1,, 'after1:=True ⟩;;

{ 'flag1 ∧ 'after1 ∧ ('turn=1 ∨ 'turn=2) }

WHILE ¬('flag2 → 'turn=2)

```

        INV { 'flag1 ∧ 'after1 ∧ ('turn=1 ∨ 'turn=2) }
        DO { 'flag1 ∧ 'after1 ∧ ('turn=1 ∨ 'turn=2) } SKIP OD;;
        { 'flag1 ∧ 'after1 ∧ ('flag2 ∧ 'after2 → 'turn=2) }
        'flag1:=False
    OD
    {False}
||
    {¬'flag2}
    WHILE True
    INV {¬'flag2}
    DO {¬'flag2} { 'flag2:=True,, 'after2:=False };;
        { 'flag2 ∧ ¬'after2 } { 'turn:=2,, 'after2:=True };;
        { 'flag2 ∧ 'after2 ∧ ('turn=1 ∨ 'turn=2) }
        WHILE ¬('flag1 → 'turn=1)
        INV { 'flag2 ∧ 'after2 ∧ ('turn=1 ∨ 'turn=2) }
        DO { 'flag2 ∧ 'after2 ∧ ('turn=1 ∨ 'turn=2) } SKIP OD;;
        { 'flag2 ∧ 'after2 ∧ ('flag1 ∧ 'after1 → 'turn=1) }
        'flag2:=False
    OD
    {False}
COEND
{False}
apply oghoare
— 122 vc
apply auto
done

```

Peterson's Algorithm III: A Solution using Semaphores

```

record Semaphores-mutex =
  out :: bool
  who :: nat

```

lemma Semaphores-mutex:

```

||- {i≠j}
  'out:=True ,,
  COBEGIN {i≠j}
    WHILE True INV {i≠j}
    DO {i≠j} AWAIT 'out THEN 'out:=False,, 'who:=i END;;
        {¬'out ∧ 'who=i ∧ i≠j} 'out:=True OD
    {False}
  ||
    {i≠j}
    WHILE True INV {i≠j}
    DO {i≠j} AWAIT 'out THEN 'out:=False,, 'who:=j END;;
        {¬'out ∧ 'who=j ∧ i≠j} 'out:=True OD
    {False}
  COEND
  {False}

```

apply *oghoare*
 — 38 vc
apply *auto*
done

Peterson's Algorithm III: Parameterized version:

lemma *Semaphores-parameterized-mutex:*

$0 < n \implies \parallel - \{ \text{True} \}$
 $\text{'out} := \text{True} \text{ , ,}$
COBEGIN
SCHEME $[0 \leq i < n]$
 $\{ \text{True} \}$
WHILE *True INV* $\{ \text{True} \}$
DO $\{ \text{True} \}$ *AWAIT* $\text{'out} := \text{False} \text{ , , 'who} := i \text{ END};$
 $\{ \neg \text{'out} \wedge \text{'who} = i \}$ $\text{'out} := \text{True} \text{ OD}$
 $\{ \text{False} \}$
COEND
 $\{ \text{False} \}$
apply *oghoare*
 — 20 vc
apply *auto*
done

The Ticket Algorithm

record *Ticket-mutex* =

num :: nat
nextv :: nat
turn :: nat list
index :: nat

lemma *Ticket-mutex:*

$\llbracket 0 < n; I = \langle n = \text{length } \text{'turn} \wedge 0 < \text{'nextv} \wedge (\forall k \ l. k < n \wedge l < n \wedge k \neq l$
 $\rightarrow \text{'turn}!k < \text{'num} \wedge (\text{'turn}!k = 0 \vee \text{'turn}!k \neq \text{'turn}!l)) \rangle \rrbracket$
 $\implies \parallel - \{ n = \text{length } \text{'turn} \}$
 $\text{'index} := 0 \text{ , ,}$
WHILE $\text{'index} < n \text{ INV}$ $\{ n = \text{length } \text{'turn} \wedge (\forall i < \text{'index}. \text{'turn}!i = 0) \}$
DO $\text{'turn} := \text{'turn}[\text{'index} := 0] \text{ , , 'index} := \text{'index} + 1 \text{ OD} \text{ , ,}$
 $\text{'num} := 1 \text{ , , 'nextv} := 1 \text{ , ,}$
COBEGIN
SCHEME $[0 \leq i < n]$
 $\{ \text{'I} \}$
WHILE *True INV* $\{ \text{'I} \}$
DO $\{ \text{'I} \}$ $\langle \text{'turn} := \text{'turn}[i := \text{'num}] \text{ , , 'num} := \text{'num} + 1 \rangle;$
 $\{ \text{'I} \}$ *WAIT* $\text{'turn}!i = \text{'nextv} \text{ END};$
 $\{ \text{'I} \wedge \text{'turn}!i = \text{'nextv} \}$ $\text{'nextv} := \text{'nextv} + 1$
OD
 $\{ \text{False} \}$
COEND

$\{\{False\}\}$
apply *oghoare*
— 35 vc
apply *simp-all*
— 16 vc
apply(*tactic* $\langle ALLGOALS$ (*clarify-tac* *context*) \rangle)
— 11 vc
apply *simp-all*
apply(*tactic* $\langle ALLGOALS$ (*clarify-tac* *context*) \rangle)
— 10 subgoals left
apply(*erule* *less-SucE*)
 apply *simp*
 apply *simp*
— 9 subgoals left
apply(*case-tac* *i=k*)
 apply *force*
 apply *simp*
apply(*case-tac* *i=l*)
 apply *force*
apply *force*
— 8 subgoals left
prefer 8
apply *force*
apply *force*
— 6 subgoals left
prefer 6
apply(*erule-tac* *x=j* **in** *allE*)
apply *fastforce*
— 5 subgoals left
prefer 5
apply(*case-tac* [!] *j=k*)
— 10 subgoals left
apply *simp-all*
apply(*erule-tac* *x=k* **in** *allE*)
apply *force*
— 9 subgoals left
apply(*case-tac* *j=l*)
 apply *simp*
 apply(*erule-tac* *x=k* **in** *allE*)
 apply(*erule-tac* *x=k* **in** *allE*)
 apply(*erule-tac* *x=l* **in** *allE*)
 apply *force*
apply(*erule-tac* *x=k* **in** *allE*)
apply(*erule-tac* *x=k* **in** *allE*)
apply(*erule-tac* *x=l* **in** *allE*)
apply *force*
— 8 subgoals left
apply *force*
apply(*case-tac* *j=l*)

```

apply simp
apply(erule-tac x=k in allE)
apply(erule-tac x=l in allE)
apply force
apply force
apply force
— 5 subgoals left
apply(erule-tac x=k in allE)
apply(erule-tac x=l in allE)
apply(case-tac j=l)
  apply force
apply force
apply force
— 3 subgoals left
apply(erule-tac x=k in allE)
apply(erule-tac x=l in allE)
apply(case-tac j=l)
  apply force
apply force
apply force
— 1 subgoals left
apply(erule-tac x=k in allE)
apply(erule-tac x=l in allE)
apply(case-tac j=l)
  apply force
apply force
done

```

1.8.2 Parallel Zero Search

Synchronized Zero Search. Zero-6

Apt and Olderog. "Verification of sequential and concurrent Programs" page 294:

```

record Zero-search =
  turn :: nat
  found :: bool
  x :: nat
  y :: nat

```

lemma *Zero-search*:

```

[[I1 = «  $a \leq 'x \wedge ('found \longrightarrow (a < 'x \wedge f('x)=0) \vee ('y \leq a \wedge f('y)=0))$ 
   $\wedge (\neg 'found \wedge a < 'x \longrightarrow f('x) \neq 0)$  » ;
  I2 = «  $'y \leq a+1 \wedge ('found \longrightarrow (a < 'x \wedge f('x)=0) \vee ('y \leq a \wedge f('y)=0))$ 
   $\wedge (\neg 'found \wedge 'y \leq a \longrightarrow f('y) \neq 0)$  » ]]  $\implies$ 
||- { $\exists u. f(u)=0$ }
'turn:=1,, 'found:= False,,
'x:=a,, 'y:=a+1 ,,
COBEGIN { 'I1 }

```



```

    WHILE  $\neg$ 'found
    INV  $\{I1\}$ 
    DO  $\{a \leq 'x \wedge ('found \longrightarrow 'y \leq a \wedge f('y)=0) \wedge (a < 'x \longrightarrow f('x) \neq 0)\}$ 
      WAIT 'turn=1 END;;
       $\{a \leq 'x \wedge ('found \longrightarrow 'y \leq a \wedge f('y)=0) \wedge (a < 'x \longrightarrow f('x) \neq 0)\}$ 
      'turn:=2;;
       $\{a \leq 'x \wedge ('found \longrightarrow 'y \leq a \wedge f('y)=0) \wedge (a < 'x \longrightarrow f('x) \neq 0)\}$ 
       $\langle 'x := 'x + 1, ,$ 
      IF  $f('x)=0$  THEN 'found:=True ELSE SKIP FI)
    OD;;
     $\{I1 \wedge 'found\}$ 
    'turn:=2
     $\{I1 \wedge 'found\}$ 
  ||
   $\{I2\}$ 
  WHILE  $\neg$ 'found
  INV  $\{I2\}$ 
  DO  $\{'y \leq a+1 \wedge ('found \longrightarrow a < 'x \wedge f('x)=0) \wedge ('y \leq a \longrightarrow f('y) \neq 0)\}$ 
    WAIT 'turn=2 END;;
     $\{'y \leq a+1 \wedge ('found \longrightarrow a < 'x \wedge f('x)=0) \wedge ('y \leq a \longrightarrow f('y) \neq 0)\}$ 
    'turn:=1;;
     $\{'y \leq a+1 \wedge ('found \longrightarrow a < 'x \wedge f('x)=0) \wedge ('y \leq a \longrightarrow f('y) \neq 0)\}$ 
     $\langle 'y := ('y - 1), ,$ 
    IF  $f('y)=0$  THEN 'found:=True ELSE SKIP FI)
  OD;;
   $\{I2 \wedge 'found\}$ 
  'turn:=1
   $\{I2 \wedge 'found\}$ 
COEND
 $\{f('x)=0 \vee f('y)=0\}$ 
apply oghoare
— 98 verification conditions
apply auto
— auto takes about 3 minutes !!
done

```

Easier Version: without AWAIT. Apt and Olderog. page 256:

lemma Zero-Search-2:

```

 $\llbracket I1 = \langle a \leq 'x \wedge ('found \longrightarrow (a < 'x \wedge f('x)=0) \vee ('y \leq a \wedge f('y)=0))$ 
   $\wedge (\neg 'found \wedge a < 'x \longrightarrow f('x) \neq 0) \rangle \rrbracket;$ 
 $I2 = \langle 'y \leq a+1 \wedge ('found \longrightarrow (a < 'x \wedge f('x)=0) \vee ('y \leq a \wedge f('y)=0))$ 
   $\wedge (\neg 'found \wedge 'y \leq a \longrightarrow f('y) \neq 0) \rangle \rrbracket \implies$ 
 $\llbracket - \{ \exists u. f(u)=0 \}$ 
  'found:= False,,
  'x:=a,, 'y:=a+1,,
COBEGIN  $\{I1\}$ 
  WHILE  $\neg$ 'found
  INV  $\{I1\}$ 
  DO  $\{a \leq 'x \wedge ('found \longrightarrow 'y \leq a \wedge f('y)=0) \wedge (a < 'x \longrightarrow f('x) \neq 0)\}$ 

```

```

      < 'x:='x+1,,IF f('x)=0 THEN 'found:=True ELSE SKIP FI>
    OD
    {'I1 ∧ 'found}
  ||
    {'I2}
    WHILE ¬'found
    INV {'I2}
    DO {'y≤a+1 ∧ ('found → a<'x ∧ f('x)=0) ∧ ('y≤a → f('y)≠0)}
      < 'y:=('y - 1),,IF f('y)=0 THEN 'found:=True ELSE SKIP FI>
    OD
    {'I2 ∧ 'found}
  COEND
  {f('x)=0 ∨ f('y)=0}
apply oghoare
— 20 vc
apply auto
— auto takes approx. 2 minutes.
done

```

1.8.3 Producer/Consumer

Previous lemmas

lemma *nat-lemma2*: $\llbracket b = m*(n::nat) + t; a = s*n + u; t=u; b-a < n \rrbracket \implies m \leq s$

proof —

```

  assume  $b = m*(n::nat) + t$   $a = s*n + u$   $t=u$ 
  hence  $(m - s) * n = b - a$  by (simp add: diff-mult-distrib)
  also assume  $\dots < n$ 
  finally have  $m - s < 1$  by simp
  thus ?thesis by arith

```

qed

lemma *mod-lemma*: $\llbracket (c::nat) \leq a; a < b; b - c < n \rrbracket \implies b \bmod n \neq a \bmod n$

```

apply(subgoal-tac  $b=b \text{ div } n*n + b \bmod n$ )
  prefer 2 apply (simp add: div-mult-mod-eq [symmetric])
apply(subgoal-tac  $a=a \text{ div } n*n + a \bmod n$ )
  prefer 2
  apply(simp add: div-mult-mod-eq [symmetric])
apply(subgoal-tac  $b - a \leq b - c$ )
  prefer 2 apply arith
apply(drule le-less-trans)
back
  apply assumption
apply(frule less-not-refl2)
apply(drule less-imp-le)
apply (drule-tac  $m = a$  and  $k = n$  in div-le-mono)
apply(safe)
apply(frule-tac  $b = b$  and  $a = a$  and  $n = n$  in nat-lemma2, assumption, assumption)

```

```

apply assumption
apply(drule order-antisym, assumption)
apply(rotate-tac -3)
apply(simp)
done

```

Producer/Consumer Algorithm

```

record Producer-consumer =
  ins :: nat
  outs :: nat
  li :: nat
  lj :: nat
  vx :: nat
  vy :: nat
  buffer :: nat list
  b :: nat list

```

The whole proof takes aprox. 4 minutes.

lemma *Producer-consumer*:

```

[[INIT = « $0 < \text{length } a \wedge 0 < \text{length } 'buffer \wedge \text{length } 'b = \text{length } a$ » ;
  I = « $(\forall k < 'ins. 'outs \leq k \longrightarrow (a ! k) = 'buffer ! (k \bmod (\text{length } 'buffer))) \wedge$ 
   $'outs \leq 'ins \wedge 'ins - 'outs \leq \text{length } 'buffer$ » ;
  I1 = « $'I \wedge 'li \leq \text{length } a$ » ;
  p1 = « $'I1 \wedge 'li = 'ins$ » ;
  I2 = « $'I \wedge (\forall k < 'lj. (a ! k) = ('b ! k)) \wedge 'lj \leq \text{length } a$ » ;
  p2 = « $'I2 \wedge 'lj = 'outs$ » ]]  $\implies$ 
||- {INIT}
'ins:=0,, 'outs:=0,, 'li:=0,, 'lj:=0,,
COBEGIN {p1  $\wedge$  INIT}
  WHILE 'li < length a
    INV {p1  $\wedge$  INIT}
    DO {p1  $\wedge$  INIT  $\wedge$  'li < length a}
      'vx := (a ! 'li);;
      {p1  $\wedge$  INIT  $\wedge$  'li < length a  $\wedge$  'vx = (a ! 'li)}
      WAIT 'ins - 'outs < length 'buffer END;
      {p1  $\wedge$  INIT  $\wedge$  'li < length a  $\wedge$  'vx = (a ! 'li)
       $\wedge$  'ins - 'outs < length 'buffer}
      'buffer := (list-update 'buffer ('ins mod (length 'buffer)) 'vx);;
      {p1  $\wedge$  INIT  $\wedge$  'li < length a
       $\wedge$  (a ! 'li) = ('buffer ! ('ins mod (length 'buffer)))
       $\wedge$  'ins - 'outs < length 'buffer}
      'ins := 'ins + 1;;
      {I1  $\wedge$  INIT  $\wedge$  ('li + 1) = 'ins  $\wedge$  'li < length a}
      'li := 'li + 1
    OD
  {p1  $\wedge$  INIT  $\wedge$  'li = length a}
||
{p2  $\wedge$  INIT}

```

```

WHILE 'lj < length a
  INV {'p2 ∧ 'INIT}
DO {'p2 ∧ 'lj < length a ∧ 'INIT}
  WAIT 'outs < 'ins END;;
  {'p2 ∧ 'lj < length a ∧ 'outs < 'ins ∧ 'INIT}
  'vy := ('buffer ! ('outs mod (length 'buffer)));;
  {'p2 ∧ 'lj < length a ∧ 'outs < 'ins ∧ 'vy = (a ! 'lj) ∧ 'INIT}
  'outs := 'outs + 1;;
  {'I2 ∧ ('lj + 1) = 'outs ∧ 'lj < length a ∧ 'vy = (a ! 'lj) ∧ 'INIT}
  'b := (list-update 'b 'lj 'vy);;
  {'I2 ∧ ('lj + 1) = 'outs ∧ 'lj < length a ∧ (a ! 'lj) = ('b ! 'lj) ∧ 'INIT}
  'lj := 'lj + 1
OD
{'p2 ∧ 'lj = length a ∧ 'INIT}
COEND
{'∀ k < length a. (a ! k) = ('b ! k)}
apply oghoare
— 138 vc
apply(tactic ⟨ALLGOALS (clarify-tac context)⟩)
— 112 subgoals left
apply(simp-all (no-asm))
— 43 subgoals left
apply(tactic ⟨ALLGOALS (conjI-Tac context (K all-tac))⟩)
— 419 subgoals left
apply(tactic ⟨ALLGOALS (clarify-tac context)⟩)
— 99 subgoals left
apply(simp-all only:length-0-conv [THEN sym])
— 20 subgoals left
apply (simp-all del:length-0-conv length-greater-0-conv add: nth-list-update mod-lemma)
— 9 subgoals left
apply (force simp add:less-Suc-eq)
apply(hypsubst-thin, drule sym)
apply (force simp add:less-Suc-eq)+
done

```

1.8.4 Parameterized Examples

Set Elements of an Array to Zero

```

record Example1 =
  a :: nat ⇒ nat

lemma Example1:
  ||- {'True}
  COBEGIN SCHEME [0 ≤ i < n] {'True} 'a := 'a (i := 0) {'a i = 0} COEND
  {'∀ i < n. 'a i = 0}
apply oghoare
apply simp-all
done

```

Same example with lists as auxiliary variables.

```

record Example1-list =
  A :: nat list
lemma Example1-list:
  ||- {n < length `A}
  COBEGIN
    SCHEME [0 ≤ i < n] {n < length `A} `A := `A[i:=0] {`A!i=0}
  COEND
  {∀ i < n. `A!i = 0}
apply oghoare
apply force+
done

```

Increment a Variable in Parallel

First some lemmas about summation properties.

```

lemma Example2-lemma2-aux: !!b. j < n ==>
  (∑ i=0..<n. (b i::nat)) =
  (∑ i=0..<j. b i) + b j + (∑ i=0..<n-(Suc j) . b (Suc j + i))
apply(induct n)
apply simp-all
apply(simp add:less-Suc-eq)
apply(auto)
apply(subgoal-tac n - j = Suc(n- Suc j))
apply simp
apply arith
done

```

```

lemma Example2-lemma2-aux2:
  !!b. j ≤ s ==> (∑ i::nat=0..<j. (b (s:=t)) i) = (∑ i=0..<j. b i)
apply(induct j)
apply simp-all
done

```

```

lemma Example2-lemma2:
  !!b. [j < n; b j = 0] ==> Suc (∑ i::nat=0..<n. b i) = (∑ i=0..<n. (b (j := Suc 0))
  i)
apply(frule-tac b=(b (j:=(Suc 0))) in Example2-lemma2-aux)
apply(erule-tac t=sum (b(j := (Suc 0))) {0..<n} in ssubst)
apply(frule-tac b=b in Example2-lemma2-aux)
apply(erule-tac t=sum b {0..<n} in ssubst)
apply(subgoal-tac Suc (sum b {0..<j} + b j + (∑ i=0..<n - Suc j. b (Suc j +
  i)))=(sum b {0..<j} + Suc (b j) + (∑ i=0..<n - Suc j. b (Suc j + i))))
apply(rotate-tac -1)
apply(erule ssubst)
apply(subgoal-tac j ≤ j)
apply(drule-tac b=b and t=(Suc 0) in Example2-lemma2-aux2)
apply(rotate-tac -1)

```

```

apply(erule ssubst)
apply simp-all
done

```

```

record Example2 =
  c :: nat ⇒ nat
  x :: nat

```

```

lemma Example2:  $0 < n \implies$ 
 $\| - \{ \text{'x} = 0 \wedge (\sum_{i=0..<n. \text{'c } i} = 0) \}$ 
COBEGIN
  SCHEME [ $0 < i < n$ ]
   $\{ \text{'x} = (\sum_{i=0..<n. \text{'c } i} \wedge \text{'c } i = 0) \}$ 
   $\langle \text{'x} := \text{'x} + (\text{Suc } 0),, \text{'c} := \text{'c } (i := (\text{Suc } 0)) \rangle$ 
   $\{ \text{'x} = (\sum_{i=0..<n. \text{'c } i} \wedge \text{'c } i = (\text{Suc } 0)) \}$ 
COEND
 $\{ \text{'x} = n \}$ 
apply oghoare
apply (simp-all cong del: sum.cong-simp)
apply (tactic  $\langle \text{ALLGOALS } (\text{clarify-tac } \textbf{context}) \rangle$ )
apply (simp-all cong del: sum.cong-simp)
  apply(erule (1) Example2-lemma2)
  apply(erule (1) Example2-lemma2)
  apply(erule (1) Example2-lemma2)
apply(simp)
done

end

```

Chapter 2

Case Study: Single and Multi-Mutator Garbage Collection Algorithms

2.1 Formalization of the Memory

theory *Graph* imports *Main* begin

datatype *node* = *Black* | *White*

type-synonym *nodes* = *node list*

type-synonym *edge* = *nat* × *nat*

type-synonym *edges* = *edge list*

consts *Roots* :: *nat set*

definition *Proper-Roots* :: *nodes* ⇒ *bool* **where**

$Proper-Roots\ M \equiv Roots \neq \{\} \wedge Roots \subseteq \{i. i < length\ M\}$

definition *Proper-Edges* :: (*nodes* × *edges*) ⇒ *bool* **where**

$Proper-Edges \equiv (\lambda(M, E). \forall i < length\ E. fst(E!i) < length\ M \wedge snd(E!i) < length\ M)$

definition *BtoW* :: (*edge* × *nodes*) ⇒ *bool* **where**

$BtoW \equiv (\lambda(e, M). (M!fst\ e) = Black \wedge (M!snd\ e) \neq Black)$

definition *Blacks* :: *nodes* ⇒ *nat set* **where**

$Blacks\ M \equiv \{i. i < length\ M \wedge M!i = Black\}$

definition *Reach* :: *edges* ⇒ *nat set* **where**

$Reach\ E \equiv \{x. (\exists path. 1 < length\ path \wedge path!(length\ path - 1) \in Roots \wedge x = path!0 \wedge (\forall i < length\ path - 1. (\exists j < length\ E. E!j = (path!(i+1), path!i)))) \vee x \in Roots\}$

Reach: the set of reachable nodes is the set of Roots together with the nodes reachable from some Root by a path represented by a list of nodes (at least two since we traverse at least one edge), where two consecutive nodes correspond to an edge in E.

2.1.1 Proofs about Graphs

lemmas *Graph-defs= Blacks-def Proper-Roots-def Proper-Edges-def BtoW-def*
declare *Graph-defs [simp]*

Graph 1

lemma *Graph1-aux [rule-format]:*

$\llbracket \text{Roots} \subseteq \text{Blacks } M; \forall i < \text{length } E. \neg \text{BtoW}(E!i, M) \rrbracket$
 $\implies 1 < \text{length path} \longrightarrow (\text{path}!(\text{length path} - 1)) \in \text{Roots} \longrightarrow$
 $(\forall i < \text{length path} - 1. (\exists j. j < \text{length } E \wedge E!j = (\text{path}!(\text{Suc } i), \text{path}!i)))$
 $\longrightarrow M!(\text{path}!0) = \text{Black}$

apply (*induct-tac path*)

apply *force*

apply *clarify*

apply *simp*

apply (*case-tac list*)

apply *force*

apply *simp*

apply (*rename-tac lista*)

apply (*rotate-tac -2*)

apply (*erule-tac x = 0 in all-dupE*)

apply *simp*

apply *clarify*

apply (*erule allE , erule (1) notE impE*)

apply *simp*

apply (*erule mp*)

apply (*case-tac lista*)

apply *force*

apply *simp*

apply (*erule mp*)

apply *clarify*

apply (*erule-tac x = Suc i in allE*)

apply *force*

done

lemma *Graph1:*

$\llbracket \text{Roots} \subseteq \text{Blacks } M; \text{Proper-Edges}(M, E); \forall i < \text{length } E. \neg \text{BtoW}(E!i, M) \rrbracket$
 $\implies \text{Reach } E \subseteq \text{Blacks } M$

apply (*unfold Reach-def*)

apply *simp*

apply *clarify*

apply (*erule disjE*)

apply *clarify*


```

apply(rule conjI)
apply(subgoal-tac 0 < length path - Suc 0)
apply(erule allE , erule (1) notE impE)
apply force
apply simp
apply(rule Graph1-aux)
apply auto
done

```

Graph 2

```

lemma Ex-first-occurrence [rule-format]:
  P (n::nat)  $\longrightarrow$  ( $\exists m. P m \wedge (\forall i. i < m \longrightarrow \neg P i)$ )
apply(rule nat-less-induct)
apply clarify
apply(case-tac  $\forall m. m < n \longrightarrow \neg P m$ )
apply auto
done

```

```

lemma Compl-lemma: (n::nat)  $\leq l \implies (\exists m. m \leq l \wedge n = l - m)$ 
apply(rule-tac x = l - n in exI)
apply arith
done

```

```

lemma Ex-last-occurrence:
   $\llbracket P (n::nat); n \leq l \rrbracket \implies (\exists m. P (l - m) \wedge (\forall i. i < m \longrightarrow \neg P (l - i)))$ 
apply(drule Compl-lemma)
apply clarify
apply(erule Ex-first-occurrence)
done

```

```

lemma Graph2:
   $\llbracket T \in \text{Reach } E; R < \text{length } E \rrbracket \implies T \in \text{Reach } (E[R := (\text{fst}(E!R), T)])$ 
apply (unfold Reach-def)
apply clarify
apply simp
apply(case-tac  $\forall z < \text{length path}. \text{fst}(E!R) \neq \text{path}!z$ )
apply(rule-tac x = path in exI)
apply simp
apply clarify
apply(erule allE , erule (1) notE impE)
apply clarify
apply(rule-tac x = j in exI)
apply(case-tac j=R)
apply(erule-tac x = Suc i in allE)
apply simp
apply (force simp add:nth-list-update)
apply simp
apply(erule exE)

```

```

apply(subgoal-tac  $z \leq \text{length path} - \text{Suc } 0$ )
  prefer 2 apply arith
apply(drule-tac  $P = \lambda m. m < \text{length path} \wedge \text{fst}(E!R) = \text{path}!m$  in Ex-last-occurrence)
  apply assumption
apply clarify
apply simp
apply(rule-tac  $x = (\text{path}!0) \# (\text{drop } (\text{length path} - \text{Suc } m) \text{ path})$  in exI)
apply simp
apply(case-tac  $\text{length path} - (\text{length path} - \text{Suc } m)$ )
  apply arith
apply simp
apply(subgoal-tac  $(\text{length path} - \text{Suc } m) + \text{nat} \leq \text{length path}$ )
  prefer 2 apply arith
apply(subgoal-tac  $\text{length path} - \text{Suc } m + \text{nat} = \text{length path} - \text{Suc } 0$ )
  prefer 2 apply arith
apply clarify
apply(case-tac i)
  apply(force simp add: nth-list-update)
apply simp
apply(subgoal-tac  $(\text{length path} - \text{Suc } m) + \text{nata} \leq \text{length path}$ )
  prefer 2 apply arith
apply(subgoal-tac  $(\text{length path} - \text{Suc } m) + (\text{Suc } \text{nata}) \leq \text{length path}$ )
  prefer 2 apply arith
apply simp
apply(erule-tac  $x = \text{length path} - \text{Suc } m + \text{nata}$  in allE)
apply simp
apply clarify
apply(rule-tac  $x = j$  in exI)
apply(case-tac  $R=j$ )
  prefer 2 apply force
apply simp
apply(drule-tac  $t = \text{path} ! (\text{length path} - \text{Suc } m)$  in sym)
apply simp
apply(case-tac  $\text{length path} - \text{Suc } 0 < m$ )
  apply(subgoal-tac  $(\text{length path} - \text{Suc } m) = 0$ )
    prefer 2 apply arith
  apply(simp del: diff-is-0-eq)
  apply(subgoal-tac  $\text{Suc } \text{nata} \leq \text{nat}$ )
  prefer 2 apply arith
  apply(drule-tac  $n = \text{Suc } \text{nata}$  in Compl-lemma)
  apply clarify
  subgoal using  $[[\text{linarith-split-limit} = 0]]$  by force
apply(drule leI)
apply(subgoal-tac  $\text{Suc } (\text{length path} - \text{Suc } m + \text{nata}) = (\text{length path} - \text{Suc } 0) - (m - \text{Suc } \text{nata})$ )
  apply(erule-tac  $x = m - (\text{Suc } \text{nata})$  in allE)
  apply(case-tac m)
    apply simp
  apply simp

```

apply *simp*
done

Graph 3

declare *min.absorb1* [*simp*] *min.absorb2* [*simp*]

lemma *Graph3*:

$\llbracket T \in \text{Reach } E; R < \text{length } E \rrbracket \implies \text{Reach}(E[R := (\text{fst}(E!R), T)]) \subseteq \text{Reach } E$

apply (*unfold Reach-def*)

apply *clarify*

apply *simp*

apply (*case-tac* $\exists i < \text{length path} - 1. (\text{fst}(E!R), T) = (\text{path}!(\text{Suc } i), \text{path}!i)$)

— the changed edge is part of the path

apply (*erule exE*)

apply (*drule-tac* $P = \lambda i. i < \text{length path} - 1 \wedge (\text{fst}(E!R), T) = (\text{path}!\text{Suc } i, \text{path}!i)$)

in *Ex-first-occurrence*)

apply *clarify*

apply (*erule disjE*)

— T is NOT a root

apply *clarify*

apply (*rule-tac* $x = (\text{take } m \text{ path})@\text{patha}$ **in** *exI*)

apply (*subgoal-tac* $\neg(\text{length path} \leq m)$)

prefer 2 **apply** *arith*

apply (*simp*)

apply (*rule conjI*)

apply (*subgoal-tac* $\neg(m + \text{length patha} - 1 < m)$)

prefer 2 **apply** *arith*

apply (*simp add: nth-append*)

apply (*rule conjI*)

apply (*case-tac* m)

apply *force*

apply (*case-tac* path)

apply *force*

apply *force*

apply *clarify*

apply (*case-tac* $\text{Suc } i \leq m$)

apply (*erule-tac* $x = i$ **in** *allE*)

apply *simp*

apply *clarify*

apply (*rule-tac* $x = j$ **in** *exI*)

apply (*case-tac* $\text{Suc } i < m$)

apply (*simp add: nth-append*)

apply (*case-tac* $R = j$)

apply (*simp add: nth-list-update*)

apply (*case-tac* $i = m$)

apply *force*

apply (*erule-tac* $x = i$ **in** *allE*)

apply *force*

```

  apply(force simp add: nth-list-update)
  apply(simp add: nth-append)
  apply(subgoal-tac i=m - 1)
  prefer 2 apply arith
  apply(case-tac R=j)
  apply(erule-tac x = m - 1 in allE)
  apply(simp add: nth-list-update)
  apply(force simp add: nth-list-update)
  apply(simp add: nth-append)
  apply(rotate-tac -4)
  apply(erule-tac x = i - m in allE)
  apply(subgoal-tac Suc (i - m)=(Suc i - m) )
  prefer 2 apply arith
  apply simp
— T is a root
  apply(case-tac m=0)
  apply force
  apply(rule-tac x = take (Suc m) path in exI)
  apply(subgoal-tac ¬(length path≤Suc m) )
  prefer 2 apply arith
  apply clarsimp
  apply(erule-tac x = i in allE)
  apply simp
  apply clarify
  apply(case-tac R=j)
  apply(force simp add: nth-list-update)
  apply(force simp add: nth-list-update)
— the changed edge is not part of the path
  apply(rule-tac x = path in exI)
  apply simp
  apply clarify
  apply(erule-tac x = i in allE)
  apply clarify
  apply(case-tac R=j)
  apply(erule-tac x = i in allE)
  apply simp
  apply(force simp add: nth-list-update)
done

```

Graph 4

lemma *Graph4*:

$\llbracket T \in \text{Reach } E; \text{Roots} \subseteq \text{Blacks } M; I \leq \text{length } E; T < \text{length } M; R < \text{length } E;$
 $\forall i < I. \neg \text{BtoW}(E!i, M); R < I; M!fst(E!R) = \text{Black}; M!T \neq \text{Black} \rrbracket \implies$
 $(\exists r. I \leq r \wedge r < \text{length } E \wedge \text{BtoW}(E[R := (fst(E!R), T)]!r, M))$

```

  apply (unfold Reach-def)
  apply simp
  apply(erule disjE)
  prefer 2 apply force

```

```

apply clarify
— there exist a black node in the path to T
apply(case-tac  $\exists m < \text{length path. } M!(\text{path!}m) = \text{Black}$ )
apply(erule exE)
apply(drule-tac  $P = \lambda m. m < \text{length path} \wedge M!(\text{path!}m) = \text{Black}$  in Ex-first-occurrence)
apply clarify
apply(case-tac ma)
apply force
apply simp
apply(case-tac length path)
apply force
apply simp
apply(erule-tac  $P = \lambda i. i < \text{nat} \longrightarrow P i$  and  $x = \text{nat}$  for  $P$  in allE)
apply simp
apply clarify
apply(erule-tac  $P = \lambda i. i < \text{Suc nat} \longrightarrow P i$  and  $x = \text{nat}$  for  $P$  in allE)
apply simp
apply(case-tac  $j < I$ )
apply(erule-tac  $x = j$  in allE)
apply force
apply(rule-tac  $x = j$  in exI)
apply(force simp add: nth-list-update)
apply simp
apply(rotate-tac  $-1$ )
apply(erule-tac  $x = \text{length path} - 1$  in allE)
apply(case-tac length path)
apply force
apply force
done

```

```

declare min.absorb1 [simp del] min.absorb2 [simp del]

```

Graph 5

lemma *Graph5*:

```

[[  $T \in \text{Reach } E$  ;  $\text{Roots} \subseteq \text{Blacks } M$ ;  $\forall i < R. \neg \text{BtoW}(E!i, M)$ ;  $T < \text{length } M$ ;
 $R < \text{length } E$ ;  $M!\text{fst}(E!R) = \text{Black}$ ;  $M!\text{snd}(E!R) = \text{Black}$ ;  $M!T \neq \text{Black}$ ]
 $\implies (\exists r. R < r \wedge r < \text{length } E \wedge \text{BtoW}(E[R := (\text{fst}(E!R), T)]!r, M))$ 

```

```

apply (unfold Reach-def)

```

```

apply simp

```

```

apply(erule disjE)

```

```

prefer 2 apply force

```

```

apply clarify

```

— there exist a black node in the path to T

```

apply(case-tac  $\exists m < \text{length path. } M!(\text{path!}m) = \text{Black}$ )

```

```

apply(erule exE)

```

```

apply(drule-tac  $P = \lambda m. m < \text{length path} \wedge M!(\text{path!}m) = \text{Black}$  in Ex-first-occurrence)

```

```

apply clarify

```

```

apply(case-tac ma)

```

```

apply force
apply simp
apply(case-tac length path)
  apply force
  apply simp
  apply(erule-tac  $P = \lambda i. i < \text{nat} \longrightarrow P\ i$  and  $x = \text{nat}$  for  $P$  in  $\text{all}E$ )
  apply simp
  apply clarify
  apply(erule-tac  $P = \lambda i. i < \text{Suc nat} \longrightarrow P\ i$  and  $x = \text{nat}$  for  $P$  in  $\text{all}E$ )
  apply simp
  apply(case-tac  $j \leq R$ )
  apply(drule le-imp-less-or-eq [of -  $R$ ])
  apply(erule  $\text{disj}E$ )
  apply(erule  $\text{all}E$  , erule (1)  $\text{not}E\ \text{imp}E$ )
  apply force
  apply force
  apply(rule-tac  $x = j$  in  $\text{ex}I$ )
  apply(force simp add: nth-list-update)
apply simp
apply(rotate-tac -1)
apply(erule-tac  $x = \text{length path} - 1$  in  $\text{all}E$ )
apply(case-tac length path)
  apply force
apply force
done

```

Other lemmas about graphs

lemma *Graph6*:

```

[[Proper-Edges( $M, E$ );  $R < \text{length } E$  ;  $T < \text{length } M$ ]]  $\implies$  Proper-Edges( $M, E[R := (\text{fst}(E!R), T)]$ )
apply (unfold Proper-Edges-def)
apply(force simp add: nth-list-update)
done

```

lemma *Graph7*:

```

[[Proper-Edges( $M, E$ )]]  $\implies$  Proper-Edges( $M[T := a], E$ )
apply (unfold Proper-Edges-def)
apply force
done

```

lemma *Graph8*:

```

[[Proper-Roots( $M$ )]]  $\implies$  Proper-Roots( $M[T := a]$ )
apply (unfold Proper-Roots-def)
apply force
done

```

Some specific lemmata for the verification of garbage collection algorithms.

```

lemma Graph9:  $j < \text{length } M \implies \text{Blacks } M \subseteq \text{Blacks } (M[j := \text{Black}])$ 
apply (unfold Blacks-def)

```

```

apply(force simp add: nth-list-update)
done

```

```

lemma Graph10 [rule-format (no-asm)]:  $\forall i. M!i=a \longrightarrow M[i:=a]=M$ 
apply(induct-tac M)
apply auto
apply(case-tac i)
apply auto
done

```

```

lemma Graph11 [rule-format (no-asm)]:
   $\llbracket M!j \neq \text{Black}; j < \text{length } M \rrbracket \implies \text{Blacks } M \subset \text{Blacks } (M[j := \text{Black}])$ 
apply (unfold Blacks-def)
apply(rule psubsetI)
apply(force simp add: nth-list-update)
apply safe
apply(erule-tac c = j in equalityCE)
apply auto
done

```

```

lemma Graph12:  $\llbracket a \subseteq \text{Blacks } M; j < \text{length } M \rrbracket \implies a \subseteq \text{Blacks } (M[j := \text{Black}])$ 
apply (unfold Blacks-def)
apply(force simp add: nth-list-update)
done

```

```

lemma Graph13:  $\llbracket a \subset \text{Blacks } M; j < \text{length } M \rrbracket \implies a \subset \text{Blacks } (M[j := \text{Black}])$ 
apply (unfold Blacks-def)
apply(erule psubset-subset-trans)
apply(force simp add: nth-list-update)
done

```

```

declare Graph-defs [simp del]

```

```

end

```

2.2 The Single Mutator Case

```

theory Gar-Coll imports Graph OG-Syntax begin

```

```

declare psubsetE [rule del]

```

Declaration of variables:

```

record gar-coll-state =
  M :: nodes
  E :: edges
  bc :: nat set
  obc :: nat set
  Ma :: nodes
  ind :: nat

```

$k :: nat$
 $z :: bool$

2.2.1 The Mutator

The mutator first redirects an arbitrary edge R from an arbitrary accessible node towards an arbitrary accessible node T . It then colors the new target T black.

We declare the arbitrarily selected node and edge as constants:

consts $R :: nat$ $T :: nat$

The following predicate states, given a list of nodes m and a list of edges e , the conditions under which the selected edge R and node T are valid:

definition $Mut-init :: gar-coll-state \Rightarrow bool$ **where**

$Mut-init \equiv \ll T \in Reach \ 'E \wedge R < length \ 'E \wedge T < length \ 'M \gg$

For the mutator we consider two modules, one for each action. An auxiliary variable $'z$ is set to false if the mutator has already redirected an edge but has not yet colored the new target.

definition $Redirect-Edge :: gar-coll-state \text{ ann-com}$ **where**

$Redirect-Edge \equiv \{ \ 'Mut-init \wedge \ 'z \} \langle \ 'E := \ 'E[R := (fst(\ 'E!R), T)], \ 'z := (\neg \ 'z) \rangle$

definition $Color-Target :: gar-coll-state \text{ ann-com}$ **where**

$Color-Target \equiv \{ \ 'Mut-init \wedge \neg \ 'z \} \langle \ 'M := \ 'M[T := Black], \ 'z := (\neg \ 'z) \rangle$

definition $Mutator :: gar-coll-state \text{ ann-com}$ **where**

$Mutator \equiv$
 $\{ \ 'Mut-init \wedge \ 'z \}$
 $WHILE \ True \ INV \ \{ \ 'Mut-init \wedge \ 'z \}$
 $DO \ Redirect-Edge \ ; \ Color-Target \ OD$

Correctness of the mutator

lemmas $mutator-defs = Mut-init-def \ Redirect-Edge-def \ Color-Target-def$

lemma $Redirect-Edge$:

$\vdash Redirect-Edge \ pre(\ Color-Target)$

apply $(unfold \ mutator-defs)$

apply $annhoare$

apply $(simp-all)$

apply $(force \ elim:Graph2)$

done

lemma $Color-Target$:

$\vdash Color-Target \ \{ \ 'Mut-init \wedge \ 'z \}$

apply $(unfold \ mutator-defs)$

apply $annhoare$

apply(*simp-all*)
done

lemma *Mutator*:
 $\vdash \text{Mutator } \{\{False\}\}$
apply(*unfold Mutator-def*)
apply *annhoare*
apply(*simp-all add:Redirect-Edge Color-Target*)
apply(*simp add:mutator-defs*)
done

2.2.2 The Collector

A constant *M-init* is used to give *'Ma* a suitable first value, defined as a list of nodes where only the *Roots* are black.

consts *M-init* :: *nodes*

definition *Proper-M-init* :: *gar-coll-state* \Rightarrow *bool* **where**
Proper-M-init \equiv « *Blacks M-init=Roots* \wedge *length M-init=length 'M* »

definition *Proper* :: *gar-coll-state* \Rightarrow *bool* **where**
Proper \equiv « *Proper-Roots 'M* \wedge *Proper-Edges('M, 'E)* \wedge *'Proper-M-init* »

definition *Safe* :: *gar-coll-state* \Rightarrow *bool* **where**
Safe \equiv « *Reach 'E* \subseteq *Blacks 'M* »

lemmas *collector-defs* = *Proper-M-init-def Proper-def Safe-def*

Blackening the roots

definition *Blacken-Roots* :: *gar-coll-state ann-com* **where**
Blacken-Roots \equiv
 $\{\{ 'Proper \}$
 $'ind:=0;;$
 $\{\{ 'Proper \wedge 'ind=0 \}$
 $WHILE 'ind < length 'M$
 $INV \{\{ 'Proper \wedge (\forall i < 'ind. i \in Roots \longrightarrow 'M!i=Black) \wedge 'ind \leq length 'M \}$
 $DO \{\{ 'Proper \wedge (\forall i < 'ind. i \in Roots \longrightarrow 'M!i=Black) \wedge 'ind < length 'M \}$
 $IF 'ind \in Roots THEN$
 $\{\{ 'Proper \wedge (\forall i < 'ind. i \in Roots \longrightarrow 'M!i=Black) \wedge 'ind < length 'M \wedge$
 $'ind \in Roots \}$
 $'M := 'M['ind:=Black] FI;;$
 $\{\{ 'Proper \wedge (\forall i < 'ind+1. i \in Roots \longrightarrow 'M!i=Black) \wedge 'ind < length 'M \}$
 $'ind := 'ind+1$
 OD

lemma *Blacken-Roots*:
 $\vdash \text{Blacken-Roots } \{\{ 'Proper \wedge Roots \subseteq Blacks 'M \}$
apply (*unfold Blacken-Roots-def*)

```

apply annhoare
apply(simp-all add:collector-defs Graph-defs)
apply safe
apply(simp-all add:nth-list-update)
  apply (erule less-SucE)
  apply simp+
  apply force
apply force
done

```

Propagating black

definition *PBInv* :: *gar-coll-state* \Rightarrow *nat* \Rightarrow *bool* **where**

$$PBInv \equiv \llbracket \lambda ind. 'obc < Blacks\ 'M \vee (\forall i < ind. \neg BtoW\ ('E!i, 'M) \vee (\neg 'z \wedge i=R \wedge (snd('E!R)) = T \wedge (\exists r. ind \leq r \wedge r < length\ 'E \wedge BtoW('E!r, 'M)))) \rrbracket$$

definition *Propagate-Black-aux* :: *gar-coll-state* *ann-com* **where**

Propagate-Black-aux \equiv

$$\{\{ 'Proper \wedge Roots \subseteq Blacks\ 'M \wedge 'obc \subseteq Blacks\ 'M \wedge 'bc \subseteq Blacks\ 'M \}$$

$$'ind := 0;;$$

$$\{\{ 'Proper \wedge Roots \subseteq Blacks\ 'M \wedge 'obc \subseteq Blacks\ 'M \wedge 'bc \subseteq Blacks\ 'M \wedge 'ind = 0 \}$$

WHILE $'ind < length\ 'E$

$$INV\ \{\{ 'Proper \wedge Roots \subseteq Blacks\ 'M \wedge 'obc \subseteq Blacks\ 'M \wedge 'bc \subseteq Blacks\ 'M$$

$$\wedge 'PBInv\ 'ind \wedge 'ind \leq length\ 'E \}$$

DO $\{\{ 'Proper \wedge Roots \subseteq Blacks\ 'M \wedge 'obc \subseteq Blacks\ 'M \wedge 'bc \subseteq Blacks\ 'M$

$$\wedge 'PBInv\ 'ind \wedge 'ind < length\ 'E \}$$

IF $'M!(fst\ ('E!\ 'ind)) = Black$ *THEN*

$$\{\{ 'Proper \wedge Roots \subseteq Blacks\ 'M \wedge 'obc \subseteq Blacks\ 'M \wedge 'bc \subseteq Blacks\ 'M$$

$$\wedge 'PBInv\ 'ind \wedge 'ind < length\ 'E \wedge 'M!(fst('E!\ 'ind)) = Black \}$$

$$'M := 'M[snd('E!\ 'ind) := Black];;$$

$$\{\{ 'Proper \wedge Roots \subseteq Blacks\ 'M \wedge 'obc \subseteq Blacks\ 'M \wedge 'bc \subseteq Blacks\ 'M$$

$$\wedge 'PBInv\ ('ind + 1) \wedge 'ind < length\ 'E \}$$

$$'ind := 'ind + 1$$

FI

OD

lemma *Propagate-Black-aux*:

\vdash *Propagate-Black-aux*

$$\{\{ 'Proper \wedge Roots \subseteq Blacks\ 'M \wedge 'obc \subseteq Blacks\ 'M \wedge 'bc \subseteq Blacks\ 'M$$

$$\wedge ('obc < Blacks\ 'M \vee 'Safe) \}$$

apply (*unfold Propagate-Black-aux-def PBInv-def collector-defs*)

apply *annhoare*

apply(*simp-all add:Graph6 Graph7 Graph8 Graph12*)

apply *force*

apply *force*

apply *force*

— 4 subgoals left

apply *clarify*

apply(*simp add:Proper-Edges-def Proper-Roots-def Graph6 Graph7 Graph8 Graph12*)

```

apply (erule disjE)
apply(rule disjI1)
apply(erule Graph13)
apply force
apply (case-tac M x ! snd (E x ! ind x)=Black)
apply (simp add: Graph10 BtoW-def)
apply (rule disjI2)
apply clarify
apply (erule less-SucE)
apply (erule-tac x=i in allE , erule (1) notE impE)
apply simp
apply clarify
apply (erule-tac y = r in le-imp-less-or-eq)
apply (erule disjE)
apply (subgoal-tac Suc (ind x)≤r)
apply fast
apply arith
apply fast
apply fast
apply(rule disjI1)
apply(erule subset-psubset-trans)
apply(erule Graph11)
apply fast
— 3 subgoals left
apply force
apply force
— last
apply clarify
apply simp
apply(subgoal-tac ind x = length (E x))
apply (simp)
apply(erule Graph1)
apply simp
apply clarify
apply(erule allE, erule impE, assumption)
apply force
apply force
apply arith
done

```

Refining propagating black

definition Auxk :: gar-coll-state ⇒ bool **where**

$$\begin{aligned}
 \text{Auxk} \equiv & \ll 'k < \text{length } 'M \wedge ('M! 'k \neq \text{Black} \vee \neg \text{BtoW}('E! 'ind, 'M) \vee \\
 & 'obc < \text{Blacks } 'M \vee (\neg 'z \wedge 'ind = R \wedge \text{snd}('E!R) = T \\
 & \wedge (\exists r. 'ind < r \wedge r < \text{length } 'E \wedge \text{BtoW}('E!r, 'M))) \gg
 \end{aligned}$$

definition Propagate-Black :: gar-coll-state ann-com **where**

$$\text{Propagate-Black} \equiv$$

```

  { ' Proper  $\wedge$  Roots $\subseteq$ Blacks ' M  $\wedge$  ' obc $\subseteq$ Blacks ' M  $\wedge$  ' bc $\subseteq$ Blacks ' M }
  ' ind:=0;;
  { ' Proper  $\wedge$  Roots $\subseteq$ Blacks ' M  $\wedge$  ' obc $\subseteq$ Blacks ' M  $\wedge$  ' bc $\subseteq$ Blacks ' M  $\wedge$  ' ind=0 }
  WHILE ' ind<length ' E
    INV { ' Proper  $\wedge$  Roots $\subseteq$ Blacks ' M  $\wedge$  ' obc $\subseteq$ Blacks ' M  $\wedge$  ' bc $\subseteq$ Blacks ' M
       $\wedge$  ' PBIInv ' ind  $\wedge$  ' ind $\leq$ length ' E }
    DO { ' Proper  $\wedge$  Roots $\subseteq$ Blacks ' M  $\wedge$  ' obc $\subseteq$ Blacks ' M  $\wedge$  ' bc $\subseteq$ Blacks ' M
       $\wedge$  ' PBIInv ' ind  $\wedge$  ' ind<length ' E }
    IF ( ' M!(fst ( ' E! ' ind))=Black THEN
      { ' Proper  $\wedge$  Roots $\subseteq$ Blacks ' M  $\wedge$  ' obc $\subseteq$ Blacks ' M  $\wedge$  ' bc $\subseteq$ Blacks ' M
         $\wedge$  ' PBIInv ' ind  $\wedge$  ' ind<length ' E  $\wedge$  ( ' M!fst(' E! ' ind))=Black }
      ' k:=(snd(' E! ' ind));
      { ' Proper  $\wedge$  Roots $\subseteq$ Blacks ' M  $\wedge$  ' obc $\subseteq$ Blacks ' M  $\wedge$  ' bc $\subseteq$ Blacks ' M
         $\wedge$  ' PBIInv ' ind  $\wedge$  ' ind<length ' E  $\wedge$  ( ' M!fst(' E! ' ind))=Black
           $\wedge$  ' Auxk }
      < ' M:= ' M[ ' k:=Black],, ' ind:= ' ind+1 >
    ELSE { ' Proper  $\wedge$  Roots $\subseteq$ Blacks ' M  $\wedge$  ' obc $\subseteq$ Blacks ' M  $\wedge$  ' bc $\subseteq$ Blacks ' M
       $\wedge$  ' PBIInv ' ind  $\wedge$  ' ind<length ' E }
      < IF ( ' M!(fst ( ' E! ' ind)) $\neq$ Black THEN ' ind:= ' ind+1 FI )
    FI
  OD

```

lemma *Propagate-Black*:

```

   $\vdash$  Propagate-Black
  { ' Proper  $\wedge$  Roots $\subseteq$ Blacks ' M  $\wedge$  ' obc $\subseteq$ Blacks ' M  $\wedge$  ' bc $\subseteq$ Blacks ' M
     $\wedge$  ( ' obc < Blacks ' M  $\vee$  ' Safe ) }
  apply (unfold Propagate-Black-def PBIInv-def Auxk-def collector-defs)
  apply annhoare
  apply (simp-all add: Graph6 Graph7 Graph8 Graph12)
    apply force
    apply force
    apply force
  — 5 subgoals left
  apply clarify
  apply (simp add: BtoW-def Proper-Edges-def)
  — 4 subgoals left
  apply clarify
  apply (simp add: Proper-Edges-def Graph6 Graph7 Graph8 Graph12)
  apply (erule disjE)
  apply (rule disjI1)
  apply (erule psubset-subset-trans)
  apply (erule Graph9)
  apply (case-tac M x!k x=Black)
  apply (case-tac M x ! snd (E x ! ind x)=Black)
  apply (simp add: Graph10 BtoW-def)
  apply (rule disjI2)
  apply clarify
  apply (erule less-SucE)
  apply (erule-tac x=i in allE , erule (1) notE impE)

```

```

apply simp
apply clarify
apply (drule-tac  $y = r$  in le-imp-less-or-eq)
apply (erule disjE)
  apply (subgoal-tac Suc ( $\text{ind } x \leq r$ ))
    apply fast
    apply arith
  apply fast
apply fast
apply (simp add: Graph10 BtoW-def)
apply (erule disjE)
apply (erule disjI1)
apply clarify
apply (erule less-SucE)
  apply force
apply simp
apply (subgoal-tac Suc  $R \leq r$ )
  apply fast
apply arith
apply(rule disjI1)
apply(erule subset-psubset-trans)
apply(erule Graph11)
apply fast
— 2 subgoals left
apply clarify
apply(simp add: Proper-Edges-def Graph6 Graph7 Graph8 Graph12)
apply (erule disjE)
  apply fast
apply clarify
apply (erule less-SucE)
apply (erule-tac  $x=i$  in allE , erule (1) notE impE)
apply simp
apply clarify
apply (drule-tac  $y = r$  in le-imp-less-or-eq)
apply (erule disjE)
  apply (subgoal-tac Suc ( $\text{ind } x \leq r$ ))
    apply fast
    apply arith
  apply (simp add: BtoW-def)
apply (simp add: BtoW-def)
— last
apply clarify
apply simp
apply(subgoal-tac  $\text{ind } x = \text{length } (E x)$ )
apply (simp)
apply(drule Graph1)
  apply simp
apply clarify
apply(erule allE, erule impE, assumption)

```

apply force
apply force
apply arith
done

Counting black nodes

definition *CountInv* :: *gar-coll-state* \Rightarrow *nat* \Rightarrow *bool* **where**
CountInv \equiv « $\lambda ind. \{i. i < ind \wedge 'Ma!i=Black\} \subseteq 'bc$ »

definition *Count* :: *gar-coll-state* *ann-com* **where**

Count \equiv
 $\{ \{ 'Proper \wedge Roots \subseteq Blacks \ 'M$
 $\wedge 'obc \subseteq Blacks \ 'Ma \wedge Blacks \ 'Ma \subseteq Blacks \ 'M \wedge 'bc \subseteq Blacks \ 'M$
 $\wedge length \ 'Ma = length \ 'M \wedge ('obc < Blacks \ 'Ma \vee 'Safe) \wedge 'bc = \{ \} \}$
 $'ind := 0;;$
 $\{ \{ 'Proper \wedge Roots \subseteq Blacks \ 'M$
 $\wedge 'obc \subseteq Blacks \ 'Ma \wedge Blacks \ 'Ma \subseteq Blacks \ 'M \wedge 'bc \subseteq Blacks \ 'M$
 $\wedge length \ 'Ma = length \ 'M \wedge ('obc < Blacks \ 'Ma \vee 'Safe) \wedge 'bc = \{ \}$
 $\wedge 'ind = 0 \}$
WHILE $'ind < length \ 'M$
INV $\{ \{ 'Proper \wedge Roots \subseteq Blacks \ 'M$
 $\wedge 'obc \subseteq Blacks \ 'Ma \wedge Blacks \ 'Ma \subseteq Blacks \ 'M \wedge 'bc \subseteq Blacks \ 'M$
 $\wedge length \ 'Ma = length \ 'M \wedge 'CountInv \ 'ind$
 $\wedge ('obc < Blacks \ 'Ma \vee 'Safe) \wedge 'ind \leq length \ 'M \}$
DO $\{ \{ 'Proper \wedge Roots \subseteq Blacks \ 'M$
 $\wedge 'obc \subseteq Blacks \ 'Ma \wedge Blacks \ 'Ma \subseteq Blacks \ 'M \wedge 'bc \subseteq Blacks \ 'M$
 $\wedge length \ 'Ma = length \ 'M \wedge 'CountInv \ 'ind$
 $\wedge ('obc < Blacks \ 'Ma \vee 'Safe) \wedge 'ind < length \ 'M \}$
IF $'M! 'ind = Black$
THEN $\{ \{ 'Proper \wedge Roots \subseteq Blacks \ 'M$
 $\wedge 'obc \subseteq Blacks \ 'Ma \wedge Blacks \ 'Ma \subseteq Blacks \ 'M \wedge 'bc \subseteq Blacks \ 'M$
 $\wedge length \ 'Ma = length \ 'M \wedge 'CountInv \ 'ind$
 $\wedge ('obc < Blacks \ 'Ma \vee 'Safe) \wedge 'ind < length \ 'M \wedge 'M! 'ind = Black \}$
 $'bc := insert \ 'ind \ 'bc$
FI;;
 $\{ \{ 'Proper \wedge Roots \subseteq Blacks \ 'M$
 $\wedge 'obc \subseteq Blacks \ 'Ma \wedge Blacks \ 'Ma \subseteq Blacks \ 'M \wedge 'bc \subseteq Blacks \ 'M$
 $\wedge length \ 'Ma = length \ 'M \wedge 'CountInv \ ('ind + 1)$
 $\wedge ('obc < Blacks \ 'Ma \vee 'Safe) \wedge 'ind < length \ 'M \}$
 $'ind := 'ind + 1$
OD

lemma *Count*:

\vdash *Count*
 $\{ \{ 'Proper \wedge Roots \subseteq Blacks \ 'M$
 $\wedge 'obc \subseteq Blacks \ 'Ma \wedge Blacks \ 'Ma \subseteq 'bc \wedge 'bc \subseteq Blacks \ 'M \wedge length \ 'Ma = length$
 $\ 'M$
 $\wedge ('obc < Blacks \ 'Ma \vee 'Safe) \}$

```

apply(unfold Count-def)
apply annhoare
apply(simp-all add:CountInv-def Graph6 Graph7 Graph8 Graph12 Blacks-def col-
lector-defs)
  apply force
  apply force
  apply force
  apply clarify
  apply simp
  apply(fast elim:less-SucE)
  apply clarify
  apply simp
  apply(fast elim:less-SucE)
apply force
apply force
done

```

Appending garbage nodes to the free list

axiomatization *Append-to-free* :: $\text{nat} \times \text{edges} \Rightarrow \text{edges}$

where

Append-to-free0: $\text{length} (\text{Append-to-free} (i, e)) = \text{length } e$ **and**

Append-to-free1: *Proper-Edges* (m, e)

$\implies \text{Proper-Edges} (m, \text{Append-to-free}(i, e))$ **and**

Append-to-free2: $i \notin \text{Reach } e$

$\implies n \in \text{Reach} (\text{Append-to-free}(i, e)) = (n = i \vee n \in \text{Reach } e)$

definition *AppendInv* :: $\text{gar-coll-state} \Rightarrow \text{nat} \Rightarrow \text{bool}$ **where**

AppendInv $\equiv \llbracket \lambda \text{ind}. \forall i < \text{length } 'M. \text{ind} \leq i \longrightarrow i \in \text{Reach } 'E \longrightarrow 'M!i = \text{Black} \rrbracket$

definition *Append* :: $\text{gar-coll-state} \text{ ann-com}$ **where**

Append \equiv

$\{ \{ 'Proper \wedge \text{Roots} \subseteq \text{Blacks } 'M \wedge 'Safe \} \}$

$'ind := 0;$

$\{ \{ 'Proper \wedge \text{Roots} \subseteq \text{Blacks } 'M \wedge 'Safe \wedge 'ind = 0 \} \}$

WHILE $'ind < \text{length } 'M$

INV $\{ \{ 'Proper \wedge 'AppendInv 'ind \wedge 'ind \leq \text{length } 'M \} \}$

DO $\{ \{ 'Proper \wedge 'AppendInv 'ind \wedge 'ind < \text{length } 'M \} \}$

IF $'M!'ind = \text{Black}$ *THEN*

$\{ \{ 'Proper \wedge 'AppendInv 'ind \wedge 'ind < \text{length } 'M \wedge 'M!'ind = \text{Black} \} \}$

$'M := 'M['ind := \text{White}]$

ELSE $\{ \{ 'Proper \wedge 'AppendInv 'ind \wedge 'ind < \text{length } 'M \wedge 'ind \notin \text{Reach } 'E \} \}$

$'E := \text{Append-to-free}('ind, 'E)$

FI;;

$\{ \{ 'Proper \wedge 'AppendInv ('ind + 1) \wedge 'ind < \text{length } 'M \} \}$

$'ind := 'ind + 1$

OD

lemma *Append*:

```

  ⊢ Append { ' Proper }
apply(unfold Append-def AppendInv-def)
apply annhoare
apply(simp-all add:collector-defs Graph6 Graph7 Graph8 Append-to-free0 Append-to-free1
Graph12)
  apply(force simp:Blacks-def nth-list-update)
  apply force
  apply force
  apply(force simp add:Graph-defs)
  apply force
apply clarify
apply simp
apply(rule conjI)
  apply (erule Append-to-free1)
apply clarify
apply (drule-tac n = i in Append-to-free2)
apply force
apply force
apply force
done

```

Correctness of the Collector

definition Collector :: gar-coll-state ann-com **where**

```

  Collector ≡
  { ' Proper }
  WHILE True INV { ' Proper }
  DO
  Blacken-Roots;;
  { ' Proper ∧ Roots ⊆ Blacks ' M }
  ' obc := { };
  { ' Proper ∧ Roots ⊆ Blacks ' M ∧ ' obc = { } }
  ' bc := Roots;;
  { ' Proper ∧ Roots ⊆ Blacks ' M ∧ ' obc = { } ∧ ' bc = Roots }
  ' Ma := M-init;;
  { ' Proper ∧ Roots ⊆ Blacks ' M ∧ ' obc = { } ∧ ' bc = Roots ∧ ' Ma = M-init }
  WHILE ' obc ≠ ' bc
  INV { ' Proper ∧ Roots ⊆ Blacks ' M
      ∧ ' obc ⊆ Blacks ' Ma ∧ Blacks ' Ma ⊆ ' bc ∧ ' bc ⊆ Blacks ' M
      ∧ length ' Ma = length ' M ∧ (' obc < Blacks ' Ma ∨ ' Safe) }
  DO { ' Proper ∧ Roots ⊆ Blacks ' M ∧ ' bc ⊆ Blacks ' M }
  ' obc := ' bc;;
  Propagate-Black;;
  { ' Proper ∧ Roots ⊆ Blacks ' M ∧ ' obc ⊆ Blacks ' M ∧ ' bc ⊆ Blacks ' M
    ∧ (' obc < Blacks ' M ∨ ' Safe) }
  ' Ma := ' M;;
  { ' Proper ∧ Roots ⊆ Blacks ' M ∧ ' obc ⊆ Blacks ' Ma
    ∧ Blacks ' Ma ⊆ Blacks ' M ∧ ' bc ⊆ Blacks ' M ∧ length ' Ma = length ' M
    ∧ (' obc < Blacks ' Ma ∨ ' Safe) }

```



```

      'bc:={};;
      Count
    OD;;
  Append
OD

```

lemma *Collector:*

```

  ⊢ Collector {False}
apply(unfold Collector-def)
apply annhoare
apply(simp-all add: Blacken-Roots Propagate-Black Count Append)
apply(simp-all add:Blacken-Roots-def Propagate-Black-def Count-def Append-def
collector-defs)
  apply (force simp add: Proper-Roots-def)
  apply force
  apply force
apply clarify
apply (erule disjE)
apply(simp add:psubsetI)
  apply(force dest:subset-antisym)
done

```

2.2.3 Interference Freedom

```

lemmas modules = Redirect-Edge-def Color-Target-def Blacken-Roots-def
Propagate-Black-def Count-def Append-def
lemmas Invariants = PBIInv-def Auxk-def CountInv-def AppendInv-def
lemmas abbrev = collector-defs mutator-defs Invariants

```

lemma *interfree-Blacken-Roots-Redirect-Edge:*

```

  interfree-aux (Some Blacken-Roots, {}, Some Redirect-Edge)
apply (unfold modules)
apply interfree-aux
apply safe
apply (simp-all add:Graph6 Graph12 abbrev)
done

```

lemma *interfree-Redirect-Edge-Blacken-Roots:*

```

  interfree-aux (Some Redirect-Edge, {}, Some Blacken-Roots)
apply (unfold modules)
apply interfree-aux
apply safe
apply(simp add:abbrev)+
done

```

lemma *interfree-Blacken-Roots-Color-Target:*

```

  interfree-aux (Some Blacken-Roots, {}, Some Color-Target)
apply (unfold modules)
apply interfree-aux

```

apply *safe*
apply(*simp-all add:Graph7 Graph8 nth-list-update abbrev*)
done

lemma *interfree-Color-Target-Blacken-Roots*:
interfree-aux (Some Color-Target, {}, Some Blacken-Roots)
apply (*unfold modules*)
apply *interfree-aux*
apply *safe*
apply(*simp add:abbrev*)
done

lemma *interfree-Propagate-Black-Redirect-Edge*:
interfree-aux (Some Propagate-Black, {}, Some Redirect-Edge)
apply (*unfold modules*)
apply *interfree-aux*
— 11 subgoals left
apply(*clarify, simp add:abbrev Graph6 Graph12*)
apply(*clarify, simp add:abbrev Graph6 Graph12*)
apply(*clarify, simp add:abbrev Graph6 Graph12*)
apply(*clarify, simp add:abbrev Graph6 Graph12*)
apply(*erule conjE*)
apply(*erule disjE, erule disjI1, rule disjI2, rule allI, (rule impI)+, case-tac R=i, rule conjI, erule sym*)
apply(*erule Graph4*)
apply(*simp*)
apply (*simp add:BtoW-def*)
apply (*simp add:BtoW-def*)
apply(*rule conjI*)
apply (*force simp add:BtoW-def*)
apply(*erule Graph4*)
apply *simp*
— 7 subgoals left
apply(*clarify, simp add:abbrev Graph6 Graph12*)
apply(*erule conjE*)
apply(*erule disjE, erule disjI1, rule disjI2, rule allI, (rule impI)+, case-tac R=i, rule conjI, erule sym*)
apply(*erule Graph4*)
apply(*simp*)
apply (*simp add:BtoW-def*)
apply (*simp add:BtoW-def*)
apply(*rule conjI*)
apply (*force simp add:BtoW-def*)
apply(*erule Graph4*)
apply *simp*
— 6 subgoals left
apply(*clarify, simp add:abbrev Graph6 Graph12*)
apply(*erule conjE*)
apply(*rule conjI*)

```

apply(erule disjE, erule disjI1, rule disjI2, rule allI, (rule impI)+, case-tac R=i,
rule conjI, erule sym)
  apply(erule Graph4)
    apply(simp)+
      apply (simp add:BtoW-def)
      apply (simp add:BtoW-def)
apply(rule conjI)
  apply (force simp add:BtoW-def)
apply(erule Graph4)
  apply simp+
apply(simp add:BtoW-def nth-list-update)
apply force
— 5 subgoals left
apply(clarify, simp add:abbrev Graph6 Graph12)
— 4 subgoals left
apply(clarify, simp add:abbrev Graph6 Graph12)
apply(rule conjI)
apply(erule disjE, erule disjI1, rule disjI2, rule allI, (rule impI)+, case-tac R=i,
rule conjI, erule sym)
  apply(erule Graph4)
    apply(simp)+
      apply (simp add:BtoW-def)
      apply (simp add:BtoW-def)
apply(rule conjI)
  apply (force simp add:BtoW-def)
apply(erule Graph4)
  apply simp+
apply(rule conjI)
  apply(simp add:nth-list-update)
apply force
apply(rule impI, rule impI, erule disjE, erule disjI1, case-tac R = (ind x) ,case-tac
M x ! T = Black)
  apply(force simp add:BtoW-def)
  apply(case-tac M x !snd (E x ! ind x)=Black)
  apply(rule disjI2)
  apply simp
  apply (erule Graph5)
  apply simp+
  apply(force simp add:BtoW-def)
apply(force simp add:BtoW-def)
— 3 subgoals left
apply(clarify, simp add:abbrev Graph6 Graph12)
— 2 subgoals left
apply(clarify, simp add:abbrev Graph6 Graph12)
apply(erule disjE, erule disjI1, rule disjI2, rule allI, (rule impI)+, case-tac R=i,
rule conjI, erule sym)
  apply clarify
apply(erule Graph4)
  apply(simp)+

```

```

apply (simp add:BtoW-def)
apply (simp add:BtoW-def)
apply(rule conjI)
apply (force simp add:BtoW-def)
apply(erule Graph4)
apply simp+
done

```

```

lemma interfree-Redirect-Edge-Propagate-Black:
  interfree-aux (Some Redirect-Edge, {}, Some Propagate-Black)
apply (unfold modules )
apply interfree-aux
apply(clarify, simp add:abbrev)+
done

```

```

lemma interfree-Propagate-Black-Color-Target:
  interfree-aux (Some Propagate-Black, {}, Some Color-Target)
apply (unfold modules )
apply interfree-aux
— 11 subgoals left
apply(clarify, simp add:abbrev Graph7 Graph8 Graph12)+
apply(erule conjE)+
apply(erule disjE,rule disjI1,erule psubset-subset-trans,erule Graph9,
  case-tac M x!T=Black, rule disjI2,rotate-tac -1, simp add: Graph10, clarify,
  erule allE, erule impE, assumption, erule impE, assumption,
  simp add:BtoW-def, rule disjI1, erule subset-psubset-trans, erule Graph11,
force)
— 7 subgoals left
apply(clarify, simp add:abbrev Graph7 Graph8 Graph12)
apply(erule conjE)+
apply(erule disjE,rule disjI1,erule psubset-subset-trans,erule Graph9,
  case-tac M x!T=Black, rule disjI2,rotate-tac -1, simp add: Graph10, clarify,
  erule allE, erule impE, assumption, erule impE, assumption,
  simp add:BtoW-def, rule disjI1, erule subset-psubset-trans, erule Graph11,
force)
— 6 subgoals left
apply(clarify, simp add:abbrev Graph7 Graph8 Graph12)
apply clarify
apply (rule conjI)
apply(erule disjE,rule disjI1,erule psubset-subset-trans,erule Graph9,
  case-tac M x!T=Black, rule disjI2,rotate-tac -1, simp add: Graph10, clarify,
  erule allE, erule impE, assumption, erule impE, assumption,
  simp add:BtoW-def, rule disjI1, erule subset-psubset-trans, erule Graph11,
force)
apply(simp add:nth-list-update)
— 5 subgoals left
apply(clarify, simp add:abbrev Graph7 Graph8 Graph12)
— 4 subgoals left
apply(clarify, simp add:abbrev Graph7 Graph8 Graph12)

```

```

apply (rule conjI)
  apply(erule disjE,rule disjI1,erule psubset-subset-trans,erule Graph9,
    case-tac M x!T=Black, rule disjI2,rotate-tac -1, simp add: Graph10, clarify,
    erule allE, erule impE, assumption, erule impE, assumption,
    simp add:BtoW-def, rule disjI1, erule subset-psubset-trans, erule Graph11,
    force)
apply(rule conjI)
apply(simp add:nth-list-update)
apply(rule impI,rule impI, case-tac M x!T=Black,rotate-tac -1, force simp add:
  BtoW-def Graph10,
  erule subset-psubset-trans, erule Graph11, force)
— 3 subgoals left
apply(clarify, simp add:abbrev Graph7 Graph8 Graph12)
— 2 subgoals left
apply(clarify, simp add:abbrev Graph7 Graph8 Graph12)
apply(erule disjE,rule disjI1,erule psubset-subset-trans,erule Graph9,
  case-tac M x!T=Black, rule disjI2,rotate-tac -1, simp add: Graph10, clarify,
  erule allE, erule impE, assumption, erule impE, assumption,
  simp add:BtoW-def, rule disjI1, erule subset-psubset-trans, erule Graph11,
  force)
— 3 subgoals left
apply(simp add:abbrev)
done

```

```

lemma interfree-Color-Target-Propagate-Black:
  interfree-aux (Some Color-Target, {}, Some Propagate-Black)
apply (unfold modules )
apply interfree-aux
apply(clarify, simp add:abbrev)+
done

```

```

lemma interfree-Count-Redirect-Edge:
  interfree-aux (Some Count, {}, Some Redirect-Edge)
apply (unfold modules)
apply interfree-aux
— 9 subgoals left
apply(simp-all add:abbrev Graph6 Graph12)
— 6 subgoals left
apply(clarify, simp add:abbrev Graph6 Graph12,
  erule disjE,erule disjI1,rule disjI2,rule subset-trans, erule Graph3,force,force)+
done

```

```

lemma interfree-Redirect-Edge-Count:
  interfree-aux (Some Redirect-Edge, {}, Some Count)
apply (unfold modules )
apply interfree-aux
apply(clarify,simp add:abbrev)+
apply(simp add:abbrev)
done

```

lemma *interfree-Count-Color-Target:*
interfree-aux (Some Count, {}, Some Color-Target)
apply (*unfold modules*)
apply *interfree-aux*
— 9 subgoals left
apply(*simp-all add:abbrev Graph7 Graph8 Graph12*)
— 6 subgoals left
apply(*clarify, simp add:abbrev Graph7 Graph8 Graph12,*
erule disjE, erule disjI1, rule disjI2, erule subset-trans, erule Graph9)
— 2 subgoals left
apply(*clarify, simp add:abbrev Graph7 Graph8 Graph12*)
apply(*rule conjI*)
apply(*erule disjE, erule disjI1, rule disjI2, erule subset-trans, erule Graph9*)
apply(*simp add:nth-list-update*)
— 1 subgoal left
apply(*clarify, simp add:abbrev Graph7 Graph8 Graph12,*
erule disjE, erule disjI1, rule disjI2, erule subset-trans, erule Graph9)
done

lemma *interfree-Color-Target-Count:*
interfree-aux (Some Color-Target, {}, Some Count)
apply (*unfold modules*)
apply *interfree-aux*
apply(*clarify, simp add:abbrev*)
apply(*simp add:abbrev*)
done

lemma *interfree-Append-Redirect-Edge:*
interfree-aux (Some Append, {}, Some Redirect-Edge)
apply (*unfold modules*)
apply *interfree-aux*
apply(*simp-all add:abbrev Graph6 Append-to-free0 Append-to-free1 Graph12*)
apply(*clarify, simp add:abbrev Graph6 Append-to-free0 Append-to-free1 Graph12,*
force dest:Graph3)
done

lemma *interfree-Redirect-Edge-Append:*
interfree-aux (Some Redirect-Edge, {}, Some Append)
apply (*unfold modules*)
apply *interfree-aux*
apply(*clarify, simp add:abbrev Append-to-free0*)
apply (*force simp add: Append-to-free2*)
apply(*clarify, simp add:abbrev Append-to-free0*)
done

lemma *interfree-Append-Color-Target:*
interfree-aux (Some Append, {}, Some Color-Target)
apply (*unfold modules*)

```

apply interfree-aux
apply(clarify, simp add:abbrev Graph7 Graph8 Append-to-free0 Append-to-free1
Graph12 nth-list-update)+
apply(simp add:abbrev Graph7 Graph8 Append-to-free0 Append-to-free1 Graph12
nth-list-update)
done

```

```

lemma interfree-Color-Target-Append:
  interfree-aux (Some Color-Target, {}, Some Append)
apply (unfold modules )
apply interfree-aux
apply(clarify, simp add:abbrev Append-to-free0)+
apply (force simp add: Append-to-free2)
apply(clarify,simp add:abbrev Append-to-free0)+
done

```

```

lemmas collector-mutator-interfree =
interfree-Blacken-Roots-Redirect-Edge interfree-Blacken-Roots-Color-Target
interfree-Propagate-Black-Redirect-Edge interfree-Propagate-Black-Color-Target
interfree-Count-Redirect-Edge interfree-Count-Color-Target
interfree-Append-Redirect-Edge interfree-Append-Color-Target
interfree-Redirect-Edge-Blacken-Roots interfree-Color-Target-Blacken-Roots
interfree-Redirect-Edge-Propagate-Black interfree-Color-Target-Propagate-Black
interfree-Redirect-Edge-Count interfree-Color-Target-Count
interfree-Redirect-Edge-Append interfree-Color-Target-Append

```

Interference freedom Collector-Mutator

```

lemma interfree-Collector-Mutator:
  interfree-aux (Some Collector, {}, Some Mutator)
apply(unfold Collector-def Mutator-def)
apply interfree-aux
apply(simp-all add:collector-mutator-interfree)
apply(unfold modules collector-defs Mut-init-def)
apply(tactic <TRYALL (interfree-aux-tac context)>)
— 32 subgoals left
apply(simp-all add:Graph6 Graph7 Graph8 Append-to-free0 Append-to-free1 Graph12)
— 20 subgoals left
apply(tactic<TRYALL (clarify-tac context)>)
apply(simp-all add:Graph6 Graph7 Graph8 Append-to-free0 Append-to-free1 Graph12)
apply(tactic <TRYALL (eresolve-tac context [disjE])>)
apply simp-all
apply(tactic <TRYALL(EVERY'[resolve-tac context [disjI2],
  resolve-tac context [subset-trans],
  eresolve-tac context @{thms Graph3},
  force-tac context,
  assume-tac context])>)
apply(tactic <TRYALL(EVERY'[resolve-tac context [disjI2],
  eresolve-tac context [subset-trans],

```

```

    resolve-tac context @{\thms Graph9},
    force-tac context])›)
apply(tactic ‹TRYALL(EVERY[resolve-tac context [disjI1],
    eresolve-tac context @{\thms psubset-subset-trans},
    resolve-tac context @{\thms Graph9},
    force-tac context])›)
done

```

Interference freedom Mutator-Collector

```

lemma interfree-Mutator-Collector:
  interfree-aux (Some Mutator, {}, Some Collector)
apply(unfold Collector-def Mutator-def)
apply interfree-aux
apply(simp-all add:collector-mutator-interfree)
apply(unfold modules collector-defs Mut-init-def)
apply(tactic ‹TRYALL (interfree-aux-tac context)›)
— 64 subgoals left
apply(simp-all add:nth-list-update Invariants Append-to-free0)+
apply(tactic‹TRYALL (clarify-tac context)›)
— 4 subgoals left
apply force
apply(simp add:Append-to-free2)
apply force
apply(simp add:Append-to-free2)
done

```

The Garbage Collection algorithm

In total there are 289 verification conditions.

```

lemma Gar-Coll:
  ||— {′Proper ∧ ′Mut-init ∧ ′z}
  COBEGIN
    Collector
    {False}
  ||
    Mutator
    {False}
  COEND
  {False}
apply oghoare
apply(force simp add: Mutator-def Collector-def modules)
apply(rule Collector)
apply(rule Mutator)
apply(simp add:interfree-Collector-Mutator)
apply(simp add:interfree-Mutator-Collector)
apply force
done

```


end

2.3 The Multi-Mutator Case

theory *Mul-Gar-Coll* **imports** *Graph OG-Syntax* **begin**

The full theory takes aprox. 18 minutes.

```
record mut =  
  Z :: bool  
  R :: nat  
  T :: nat
```

Declaration of variables:

```
record mul-gar-coll-state =  
  M :: nodes  
  E :: edges  
  bc :: nat set  
  obc :: nat set  
  Ma :: nodes  
  ind :: nat  
  k :: nat  
  q :: nat  
  l :: nat  
  Muts :: mut list
```

2.3.1 The Mutators

definition *Mul-mut-init* :: *mul-gar-coll-state* \Rightarrow *nat* \Rightarrow *bool* **where**
$$\begin{aligned} \textit{Mul-mut-init} \equiv \ll \lambda n. n = \textit{length } 'Muts \wedge (\forall i < n. R ('Muts!i) < \textit{length } 'E \\ \wedge T ('Muts!i) < \textit{length } 'M) \gg \end{aligned}$$

definition *Mul-Redirect-Edge* :: *nat* \Rightarrow *nat* \Rightarrow *mul-gar-coll-state* *ann-com* **where**
$$\begin{aligned} \textit{Mul-Redirect-Edge } j \ n \equiv \\ \{ \textit{Mul-mut-init } n \wedge Z ('Muts!j) \} \\ \langle \textit{IF } T ('Muts!j) \in \textit{Reach } 'E \ \textit{THEN} \\ 'E := 'E[R ('Muts!j) := (\textit{fst } ('E!R ('Muts!j))), T ('Muts!j))] \ \textit{FI}, \\ 'Muts := 'Muts[j := ('Muts!j)] \ (Z := \textit{False}) \rangle \end{aligned}$$

definition *Mul-Color-Target* :: *nat* \Rightarrow *nat* \Rightarrow *mul-gar-coll-state* *ann-com* **where**
$$\begin{aligned} \textit{Mul-Color-Target } j \ n \equiv \\ \{ \textit{Mul-mut-init } n \wedge \neg Z ('Muts!j) \} \\ \langle 'M := 'M[T ('Muts!j) := \textit{Black}], 'Muts := 'Muts[j := ('Muts!j)] \ (Z := \textit{True}) \rangle \end{aligned}$$

definition *Mul-Mutator* :: *nat* \Rightarrow *nat* \Rightarrow *mul-gar-coll-state* *ann-com* **where**
$$\begin{aligned} \textit{Mul-Mutator } j \ n \equiv \\ \{ \textit{Mul-mut-init } n \wedge Z ('Muts!j) \} \\ \textit{WHILE } \textit{True} \\ \ \ \ \ \textit{INV } \{ \textit{Mul-mut-init } n \wedge Z ('Muts!j) \} \end{aligned}$$

DO Mul-Redirect-Edge j n ;;
Mul-Color-Target j n
OD

lemmas *mul-mutator-defs = Mul-mut-init-def Mul-Redirect-Edge-def Mul-Color-Target-def*

Correctness of the proof outline of one mutator

lemma *Mul-Redirect-Edge: $0 \leq j \wedge j < n \implies$*
 \vdash *Mul-Redirect-Edge j n*
pre(Mul-Color-Target j n)
apply (*unfold mul-mutator-defs*)
apply *annhoare*
apply (*simp-all*)
apply *clarify*
apply (*simp add: nth-list-update*)
done

lemma *Mul-Color-Target: $0 \leq j \wedge j < n \implies$*
 \vdash *Mul-Color-Target j n*
 $\{\{ 'Mul-mut-init n \wedge Z ('Muts!j) \}\}$
apply (*unfold mul-mutator-defs*)
apply *annhoare*
apply (*simp-all*)
apply *clarify*
apply (*simp add: nth-list-update*)
done

lemma *Mul-Mutator: $0 \leq j \wedge j < n \implies$*
 \vdash *Mul-Mutator j n $\{\{ False \}\}$*
apply (*unfold Mul-Mutator-def*)
apply *annhoare*
apply (*simp-all add: Mul-Redirect-Edge Mul-Color-Target*)
apply (*simp add: mul-mutator-defs Mul-Redirect-Edge-def*)
done

Interference freedom between mutators

lemma *Mul-interfree-Redirect-Edge-Redirect-Edge:*
 $\llbracket 0 \leq i; i < n; 0 \leq j; j < n; i \neq j \rrbracket \implies$
interfree-awx (Some (Mul-Redirect-Edge i n), {}, Some (Mul-Redirect-Edge j n))
apply (*unfold mul-mutator-defs*)
apply *interfree-awx*
apply *safe*
apply (*simp-all add: nth-list-update*)
done

lemma *Mul-interfree-Redirect-Edge-Color-Target:*
 $\llbracket 0 \leq i; i < n; 0 \leq j; j < n; i \neq j \rrbracket \implies$
interfree-awx (Some (Mul-Redirect-Edge i n), {}, Some (Mul-Color-Target j n))

```

apply (unfold mul-mutator-defs)
apply interfree-aux
apply safe
apply(simp-all add: nth-list-update)
done

```

```

lemma Mul-interfree-Color-Target-Redirect-Edge:
   $[[0 \leq i; i < n; 0 \leq j; j < n; i \neq j]] \implies$ 
  interfree-aux (Some(Mul-Color-Target i n), {}, Some(Mul-Redirect-Edge j n))
apply (unfold mul-mutator-defs)
apply interfree-aux
apply safe
apply(simp-all add: nth-list-update)
done

```

```

lemma Mul-interfree-Color-Target-Color-Target:
   $[[0 \leq i; i < n; 0 \leq j; j < n; i \neq j]] \implies$ 
  interfree-aux (Some(Mul-Color-Target i n), {}, Some(Mul-Color-Target j n))
apply (unfold mul-mutator-defs)
apply interfree-aux
apply safe
apply(simp-all add: nth-list-update)
done

```

```

lemmas mul-mutator-interfree =
  Mul-interfree-Redirect-Edge-Redirect-Edge Mul-interfree-Redirect-Edge-Color-Target
  Mul-interfree-Color-Target-Redirect-Edge Mul-interfree-Color-Target-Color-Target

```

```

lemma Mul-interfree-Mutator-Mutator:  $[[i < n; j < n; i \neq j]] \implies$ 
  interfree-aux (Some (Mul-Mutator i n), {}, Some (Mul-Mutator j n))
apply(unfold Mul-Mutator-def)
apply(interfree-aux)
apply(simp-all add: mul-mutator-interfree)
apply(simp-all add: mul-mutator-defs)
apply(tactic <TRYALL (interfree-aux-tac context)>)
apply(tactic <ALLGOALS (clarify-tac context)>)
apply (simp-all add: nth-list-update)
done

```

Modular Parameterized Mutators

```

lemma Mul-Parameterized-Mutators:  $0 < n \implies$ 
  ||- { Mul-mut-init n  $\wedge$  ( $\forall i < n. Z ('Muts!i)$ ) }
  COBEGIN
  SCHEME [0 ≤ j < n]
  Mul-Mutator j n
  {False}
  COEND
  {False}

```

```

apply oghoare
apply(force simp add:Mul-Mutator-def mul-mutator-defs nth-list-update)
apply(erule Mul-Mutator)
apply(simp add:Mul-interfree-Mutator-Mutator)
apply(force simp add:Mul-Mutator-def mul-mutator-defs nth-list-update)
done

```

2.3.2 The Collector

definition *Queue* :: *mul-gar-coll-state* \Rightarrow *nat* **where**
Queue \equiv « *length* (*filter* ($\lambda i. \neg Z\ i \wedge 'M!(T\ i) \neq Black$) '*Muts*) »

consts *M-init* :: *nodes*

definition *Proper-M-init* :: *mul-gar-coll-state* \Rightarrow *bool* **where**
Proper-M-init \equiv « *Blacks* *M-init=Roots* \wedge *length* *M-init=length* '*M* »

definition *Mul-Proper* :: *mul-gar-coll-state* \Rightarrow *nat* \Rightarrow *bool* **where**
Mul-Proper \equiv « $\lambda n. Proper\ Roots\ 'M \wedge Proper\ Edges\ ('M, 'E) \wedge 'Proper\ M\ init$
 $\wedge n=length\ 'Muts$ »

definition *Safe* :: *mul-gar-coll-state* \Rightarrow *bool* **where**
Safe \equiv « *Reach* '*E* \subseteq *Blacks* '*M* »

lemmas *mul-collector-defs* = *Proper-M-init-def* *Mul-Proper-def* *Safe-def*

Blackening Roots

definition *Mul-Blacken-Roots* :: *nat* \Rightarrow *mul-gar-coll-state* *ann-com* **where**
Mul-Blacken-Roots *n* \equiv
 $\{ 'Mul\ Proper\ n \}$
 $'ind:=0;;$
 $\{ 'Mul\ Proper\ n \wedge 'ind=0 \}$
 $WHILE\ 'ind < length\ 'M$
 $INV\ \{ 'Mul\ Proper\ n \wedge (\forall i < 'ind. i \in Roots \longrightarrow 'M!i=Black) \wedge 'ind \leq length$
 $'M \}$
 $DO\ \{ 'Mul\ Proper\ n \wedge (\forall i < 'ind. i \in Roots \longrightarrow 'M!i=Black) \wedge 'ind < length\ 'M \}$
 $IF\ 'ind \in Roots\ THEN$
 $\{ 'Mul\ Proper\ n \wedge (\forall i < 'ind. i \in Roots \longrightarrow 'M!i=Black) \wedge 'ind < length\ 'M \wedge$
 $'ind \in Roots \}$
 $'M := 'M['ind := Black]\ FI;;$
 $\{ 'Mul\ Proper\ n \wedge (\forall i < 'ind+1. i \in Roots \longrightarrow 'M!i=Black) \wedge 'ind < length$
 $'M \}$
 $'ind := 'ind+1$
 OD

lemma *Mul-Blacken-Roots*:
 $\vdash Mul\ Blacken\ Roots\ n$
 $\{ 'Mul\ Proper\ n \wedge Roots \subseteq Blacks\ 'M \}$
apply (*unfold* *Mul-Blacken-Roots-def*)

```

apply annhoare
apply(simp-all add:mul-collector-defs Graph-defs)
apply safe
apply(simp-all add:nth-list-update)
  apply (erule less-SucE)
  apply simp+
  apply force
apply force
done

```

Propagating Black

definition *Mul-PBInv* :: *mul-gar-coll-state* \Rightarrow *bool* **where**

$$\begin{aligned} \textit{Mul-PBInv} \equiv & \langle \langle \textit{Safe} \vee \textit{obc} \subseteq \textit{Blacks} \textit{'M} \vee \textit{l} < \textit{Queue} \\ & \vee (\forall i < \textit{ind}. \neg \textit{BtoW}(\textit{E!}i, \textit{'M})) \wedge \textit{l} \leq \textit{Queue} \rangle \rangle \end{aligned}$$

definition *Mul-Auxk* :: *mul-gar-coll-state* \Rightarrow *bool* **where**

$$\textit{Mul-Auxk} \equiv \langle \langle \textit{l} < \textit{Queue} \vee \textit{M!}k \neq \textit{Black} \vee \neg \textit{BtoW}(\textit{E!} \textit{ind}, \textit{'M}) \vee \textit{obc} \subseteq \textit{Blacks} \textit{'M} \rangle \rangle$$

definition *Mul-Propagate-Black* :: *nat* \Rightarrow *mul-gar-coll-state ann-com* **where**

$$\textit{Mul-Propagate-Black} \textit{n} \equiv$$

$$\{ \{ \textit{Mul-Prop} \textit{n} \wedge \textit{Roots} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{obc} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{bc} \subseteq \textit{Blacks} \textit{'M} \\ \wedge (\textit{Safe} \vee \textit{l} \leq \textit{Queue} \vee \textit{obc} \subseteq \textit{Blacks} \textit{'M}) \} \}$$

$$\textit{ind} := 0;;$$

$$\{ \{ \textit{Mul-Prop} \textit{n} \wedge \textit{Roots} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge \textit{obc} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{Blacks} \textit{'M} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{bc} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge (\textit{Safe} \vee \textit{l} \leq \textit{Queue} \vee \textit{obc} \subseteq \textit{Blacks} \textit{'M}) \wedge \textit{ind} = 0 \}$$

WHILE $\textit{ind} < \textit{length} \textit{'E}$

INV $\{ \{ \textit{Mul-Prop} \textit{n} \wedge \textit{Roots} \subseteq \textit{Blacks} \textit{'M}$

$$\wedge \textit{obc} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{bc} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge \textit{Mul-PBInv} \wedge \textit{ind} \leq \textit{length} \textit{'E} \}$$

DO $\{ \{ \textit{Mul-Prop} \textit{n} \wedge \textit{Roots} \subseteq \textit{Blacks} \textit{'M}$

$$\wedge \textit{obc} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{bc} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge \textit{Mul-PBInv} \wedge \textit{ind} < \textit{length} \textit{'E} \}$$

IF $\textit{M!}(\textit{fst}(\textit{E!} \textit{ind})) = \textit{Black}$ **THEN**

$$\{ \{ \textit{Mul-Prop} \textit{n} \wedge \textit{Roots} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge \textit{obc} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{bc} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge \textit{Mul-PBInv} \wedge (\textit{M!} \textit{fst}(\textit{E!} \textit{ind})) = \textit{Black} \wedge \textit{ind} < \textit{length} \textit{'E} \}$$

$$\textit{k} := \textit{snd}(\textit{E!} \textit{ind});;$$

$$\{ \{ \textit{Mul-Prop} \textit{n} \wedge \textit{Roots} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge \textit{obc} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{bc} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge (\textit{Safe} \vee \textit{obc} \subseteq \textit{Blacks} \textit{'M} \vee \textit{l} < \textit{Queue} \vee (\forall i < \textit{ind}. \neg \textit{BtoW}(\textit{E!}i, \textit{'M}))$$

$$\wedge \textit{l} \leq \textit{Queue} \wedge \textit{Mul-Auxk}) \wedge \textit{k} < \textit{length} \textit{'M} \wedge \textit{M!} \textit{fst}(\textit{E!} \textit{ind}) = \textit{Black}$$

$$\wedge \textit{ind} < \textit{length} \textit{'E} \}$$

$$\langle \textit{M} := \textit{M}[\textit{k} := \textit{Black}], \textit{ind} := \textit{ind} + 1 \rangle$$

ELSE $\{ \{ \textit{Mul-Prop} \textit{n} \wedge \textit{Roots} \subseteq \textit{Blacks} \textit{'M}$

$$\wedge \textit{obc} \subseteq \textit{Blacks} \textit{'M} \wedge \textit{bc} \subseteq \textit{Blacks} \textit{'M}$$

$$\wedge \textit{Mul-PBInv} \wedge \textit{ind} < \textit{length} \textit{'E} \}$$

$\langle IF \ 'M!(fst \ ('E! \ 'ind)) \neq Black \ THEN \ 'ind := \ 'ind + 1 \ FI \rangle \ FI$
OD

lemma *Mul-Propagate-Black:*

$\vdash \text{Mul-Propagate-Black } n$
 $\{ \{ \text{Mul-Propagate-Black } n \wedge \text{Roots} \subseteq \text{Blacks } \ 'M \wedge \ 'obc \subseteq \text{Blacks } \ 'M \wedge \ 'bc \subseteq \text{Blacks } \ 'M$
 $\wedge \ ('Safe \vee \ 'obc \subseteq \text{Blacks } \ 'M \vee \ 'l < \text{Queue} \wedge \ ('l \leq \text{Queue} \vee \ 'obc \subseteq \text{Blacks}$
 $\ 'M)) \}$
apply (*unfold Mul-Propagate-Black-def*)
apply *annhoare*
apply (*simp-all add: Mul-PBInv-def mul-collector-defs Mul-Auxk-def Graph6 Graph7*
Graph8 Graph12 mul-collector-defs Queue-def)
— 8 subgoals left
apply *force*
apply *force*
apply *force*
apply (*force simp add: BtoW-def Graph-defs*)
— 4 subgoals left
apply *clarify*
apply (*simp add: mul-collector-defs Graph12 Graph6 Graph7 Graph8*)
apply (*disjE-tac*)
apply (*simp-all add: Graph12 Graph13*)
apply (*case-tac M x! k x=Black*)
apply (*simp add: Graph10*)
apply (*rule disjI2, rule disjI1, erule subset-psubset-trans, erule Graph11, force*)
apply (*case-tac M x! k x=Black*)
apply (*simp add: Graph10 BtoW-def*)
apply (*rule disjI2, clarify, erule less-SucE, force*)
apply (*case-tac M x! snd(E x! ind x)=Black*)
apply (*force*)
apply (*force*)
apply (*rule disjI2, rule disjI1, erule subset-psubset-trans, erule Graph11, force*)
— 2 subgoals left
apply *clarify*
apply (*conjI-tac*)
apply (*disjE-tac*)
apply (*simp-all*)
apply *clarify*
apply (*erule less-SucE*)
apply *force*
apply (*simp add: BtoW-def*)
— 1 subgoal left
apply *clarify*
apply *simp*
apply (*disjE-tac*)
apply (*simp-all*)
apply (*rule disjI1, rule Graph1*)
apply *simp-all*
done

Counting Black Nodes

definition *Mul-CountInv* :: *mul-gar-coll-state* \Rightarrow *nat* \Rightarrow *bool* **where**
Mul-CountInv \equiv « $\lambda ind. \{i. i < ind \wedge 'Ma!i=Black\} \subseteq 'bc$ »

definition *Mul-Count* :: *nat* \Rightarrow *mul-gar-coll-state ann-com* **where**

Mul-Count *n* \equiv
 $\{ \{ 'Mul-Prop\text{-}n \wedge \text{Roots} \subseteq \text{Blacks } 'M$
 $\wedge 'obc \subseteq \text{Blacks } 'Ma \wedge \text{Blacks } 'Ma \subseteq \text{Blacks } 'M \wedge 'bc \subseteq \text{Blacks } 'M$
 $\wedge \text{length } 'Ma = \text{length } 'M$
 $\wedge ('Safe \vee 'obc \subseteq \text{Blacks } 'Ma \vee 'l < 'q \wedge ('q \leq 'Queue \vee 'obc \subseteq \text{Blacks } 'M))$
 $\wedge 'q < n+1 \wedge 'bc = \{\}\}$
 $'ind := 0;$
 $\{ \{ 'Mul-Prop\text{-}n \wedge \text{Roots} \subseteq \text{Blacks } 'M$
 $\wedge 'obc \subseteq \text{Blacks } 'Ma \wedge \text{Blacks } 'Ma \subseteq \text{Blacks } 'M \wedge 'bc \subseteq \text{Blacks } 'M$
 $\wedge \text{length } 'Ma = \text{length } 'M$
 $\wedge ('Safe \vee 'obc \subseteq \text{Blacks } 'Ma \vee 'l < 'q \wedge ('q \leq 'Queue \vee 'obc \subseteq \text{Blacks } 'M))$
 $\wedge 'q < n+1 \wedge 'bc = \{\} \wedge 'ind = 0\}$
WHILE $'ind < \text{length } 'M$
INV $\{ \{ 'Mul-Prop\text{-}n \wedge \text{Roots} \subseteq \text{Blacks } 'M$
 $\wedge 'obc \subseteq \text{Blacks } 'Ma \wedge \text{Blacks } 'Ma \subseteq \text{Blacks } 'M \wedge 'bc \subseteq \text{Blacks } 'M$
 $\wedge \text{length } 'Ma = \text{length } 'M \wedge 'Mul-CountInv 'ind$
 $\wedge ('Safe \vee 'obc \subseteq \text{Blacks } 'Ma \vee 'l < 'q \wedge ('q \leq 'Queue \vee 'obc \subseteq \text{Blacks } 'M))$
 $\wedge 'q < n+1 \wedge 'ind \leq \text{length } 'M\}$
DO $\{ \{ 'Mul-Prop\text{-}n \wedge \text{Roots} \subseteq \text{Blacks } 'M$
 $\wedge 'obc \subseteq \text{Blacks } 'Ma \wedge \text{Blacks } 'Ma \subseteq \text{Blacks } 'M \wedge 'bc \subseteq \text{Blacks } 'M$
 $\wedge \text{length } 'Ma = \text{length } 'M \wedge 'Mul-CountInv 'ind$
 $\wedge ('Safe \vee 'obc \subseteq \text{Blacks } 'Ma \vee 'l < 'q \wedge ('q \leq 'Queue \vee 'obc \subseteq \text{Blacks } 'M))$
 $\wedge 'q < n+1 \wedge 'ind < \text{length } 'M\}$
IF $'M!'ind = Black$
THEN $\{ \{ 'Mul-Prop\text{-}n \wedge \text{Roots} \subseteq \text{Blacks } 'M$
 $\wedge 'obc \subseteq \text{Blacks } 'Ma \wedge \text{Blacks } 'Ma \subseteq \text{Blacks } 'M \wedge 'bc \subseteq \text{Blacks } 'M$
 $\wedge \text{length } 'Ma = \text{length } 'M \wedge 'Mul-CountInv 'ind$
 $\wedge ('Safe \vee 'obc \subseteq \text{Blacks } 'Ma \vee 'l < 'q \wedge ('q \leq 'Queue \vee 'obc \subseteq \text{Blacks } 'M))$
 $'M))$
 $\wedge 'q < n+1 \wedge 'ind < \text{length } 'M \wedge 'M!'ind = Black\}$
 $'bc := \text{insert } 'ind 'bc$
FI;;
 $\{ \{ 'Mul-Prop\text{-}n \wedge \text{Roots} \subseteq \text{Blacks } 'M$
 $\wedge 'obc \subseteq \text{Blacks } 'Ma \wedge \text{Blacks } 'Ma \subseteq \text{Blacks } 'M \wedge 'bc \subseteq \text{Blacks } 'M$
 $\wedge \text{length } 'Ma = \text{length } 'M \wedge 'Mul-CountInv ('ind+1)$
 $\wedge ('Safe \vee 'obc \subseteq \text{Blacks } 'Ma \vee 'l < 'q \wedge ('q \leq 'Queue \vee 'obc \subseteq \text{Blacks } 'M))$
 $\wedge 'q < n+1 \wedge 'ind < \text{length } 'M\}$
 $'ind := 'ind+1$
OD

lemma *Mul-Count*:

$\vdash 'Mul-Count 'n$

$\{ \{ 'Mul-Prop\text{-}n \wedge \text{Roots} \subseteq \text{Blacks } 'M$
 $\wedge 'obc \subseteq \text{Blacks } 'Ma \wedge \text{Blacks } 'Ma \subseteq \text{Blacks } 'M \wedge 'bc \subseteq \text{Blacks } 'M$

$\wedge \text{length } 'Ma = \text{length } 'M \wedge \text{Blacks } 'Ma \subseteq 'bc$
 $\wedge ('Safe \vee 'obc \subseteq \text{Blacks } 'Ma \vee 'l < 'q \wedge ('q \leq 'Queue \vee 'obc \subseteq \text{Blacks } 'M))$
 $\wedge 'q < n + 1 \}$
apply (*unfold Mul-Count-def*)
apply *annhoare*
apply(*simp-all add:Mul-CountInv-def mul-collector-defs Mul-Auxk-def Graph6 Graph7 Graph8 Graph12 mul-collector-defs Queue-def*)
— 7 subgoals left
apply *force*
apply *force*
apply *force*
— 4 subgoals left
apply *clarify*
apply(*conjI-tac*)
apply(*disjE-tac*)
apply *simp-all*
apply(*simp add:Blacks-def*)
apply *clarify*
apply(*erule less-SucE*)
back
apply *force*
apply *force*
— 3 subgoals left
apply *clarify*
apply(*conjI-tac*)
apply(*disjE-tac*)
apply *simp-all*
apply *clarify*
apply(*erule less-SucE*)
back
apply *force*
apply *simp*
apply(*rotate-tac -1*)
apply (*force simp add:Blacks-def*)
— 2 subgoals left
apply *force*
— 1 subgoal left
apply *clarify*
apply(*drule-tac x = ind x in le-imp-less-or-eq*)
apply (*simp-all add:Blacks-def*)
done

Appending garbage nodes to the free list

axiomatization *Append-to-free* :: $\text{nat} \times \text{edges} \Rightarrow \text{edges}$

where

Append-to-free0: $\text{length } (\text{Append-to-free } (i, e)) = \text{length } e$ **and**

Append-to-free1: $\text{Proper-Edges } (m, e)$

$\implies \text{Proper-Edges } (m, \text{Append-to-free}(i, e))$ **and**

Append-to-free2: $i \notin \text{Reach } e$
 $\implies n \in \text{Reach } (\text{Append-to-free}(i, e)) = (n = i \vee n \in \text{Reach } e)$

definition *Mul-AppendInv* :: *mul-gar-coll-state* \Rightarrow *nat* \Rightarrow *bool* **where**
Mul-AppendInv \equiv « $\lambda \text{ind. } (\forall i. \text{ind} \leq i \longrightarrow i < \text{length } 'M \longrightarrow i \in \text{Reach } 'E \longrightarrow 'M!i = \text{Black})$ »

definition *Mul-Append* :: *nat* \Rightarrow *mul-gar-coll-state* *ann-com* **where**
Mul-Append $n \equiv$
 $\{ \text{'Mul-Prop} \text{er } n \wedge \text{Roots} \subseteq \text{Blacks } 'M \wedge \text{'Safe} \}$
 $\text{'ind} := 0;$
 $\{ \text{'Mul-Prop} \text{er } n \wedge \text{Roots} \subseteq \text{Blacks } 'M \wedge \text{'Safe} \wedge \text{'ind} = 0 \}$
WHILE $\text{'ind} < \text{length } 'M$
 INV $\{ \text{'Mul-Prop} \text{er } n \wedge \text{'Mul-AppendInv } \text{'ind} \wedge \text{'ind} \leq \text{length } 'M \}$
 DO $\{ \text{'Mul-Prop} \text{er } n \wedge \text{'Mul-AppendInv } \text{'ind} \wedge \text{'ind} < \text{length } 'M \}$
 IF $\text{'M}! \text{'ind} = \text{Black}$ **THEN**
 $\{ \text{'Mul-Prop} \text{er } n \wedge \text{'Mul-AppendInv } \text{'ind} \wedge \text{'ind} < \text{length } 'M \wedge \text{'M}! \text{'ind} = \text{Black} \}$
 $\text{'M} := \text{'M}[\text{'ind} := \text{White}]$
 ELSE
 $\{ \text{'Mul-Prop} \text{er } n \wedge \text{'Mul-AppendInv } \text{'ind} \wedge \text{'ind} < \text{length } 'M \wedge \text{'ind} \notin \text{Reach } 'E \}$
 $\text{'E} := \text{Append-to-free}(\text{'ind}, 'E)$
 FI;
 $\{ \text{'Mul-Prop} \text{er } n \wedge \text{'Mul-AppendInv } (\text{'ind} + 1) \wedge \text{'ind} < \text{length } 'M \}$
 $\text{'ind} := \text{'ind} + 1$
 OD

lemma *Mul-Append*:
 $\vdash \text{Mul-Append } n$
 $\{ \text{'Mul-Prop} \text{er } n \}$
apply (*unfold Mul-Append-def*)
apply *annhoare*
apply (*simp-all add: mul-collector-defs Mul-AppendInv-def*
 Graph6 Graph7 Graph8 Append-to-free0 Append-to-free1 Graph12)
apply (*force simp add: Blacks-def*)
apply (*force simp add: Blacks-def*)
apply (*force simp add: Blacks-def*)
apply (*force simp add: Graph-defs*)
apply *force*
apply (*force simp add: Append-to-free1 Append-to-free2*)
apply *force*
apply *force*
done

Collector

definition *Mul-Collector* :: *nat* \Rightarrow *mul-gar-coll-state* *ann-com* **where**
Mul-Collector $n \equiv$
 $\{ \text{'Mul-Prop} \text{er } n \}$

```

WHILE True INV { Mul-Propser n }
DO
  Mul-Blacken-Roots n ;;
  { Mul-Propser n ∧ Roots ⊆ Blacks 'M }
  'obc := { } ;;
  { Mul-Propser n ∧ Roots ⊆ Blacks 'M ∧ 'obc = { } }
  'bc := Roots ;;
  { Mul-Propser n ∧ Roots ⊆ Blacks 'M ∧ 'obc = { } ∧ 'bc = Roots }
  'l := 0 ;;
  { Mul-Propser n ∧ Roots ⊆ Blacks 'M ∧ 'obc = { } ∧ 'bc = Roots ∧ 'l = 0 }
  WHILE 'l < n + 1
    INV { Mul-Propser n ∧ Roots ⊆ Blacks 'M ∧ 'bc ⊆ Blacks 'M ∧
          ( 'Safe ∨ ( 'l ≤ 'Queue ∨ 'bc ⊆ Blacks 'M ) ) ∧ 'l < n + 1 }
    DO { Mul-Propser n ∧ Roots ⊆ Blacks 'M ∧ 'bc ⊆ Blacks 'M
          ∧ ( 'Safe ∨ 'l ≤ 'Queue ∨ 'bc ⊆ Blacks 'M ) }
        'obc := 'bc ;;
        Mul-Propagate-Black n ;;
        { Mul-Propser n ∧ Roots ⊆ Blacks 'M
          ∧ 'obc ⊆ Blacks 'M ∧ 'bc ⊆ Blacks 'M
          ∧ ( 'Safe ∨ 'obc ⊆ Blacks 'M ∨ 'l < 'Queue
            ∧ ( 'l ≤ 'Queue ∨ 'obc ⊆ Blacks 'M ) ) }
        'bc := { } ;;
        { Mul-Propser n ∧ Roots ⊆ Blacks 'M
          ∧ 'obc ⊆ Blacks 'M ∧ 'bc ⊆ Blacks 'M
          ∧ ( 'Safe ∨ 'obc ⊆ Blacks 'M ∨ 'l < 'Queue
            ∧ ( 'l ≤ 'Queue ∨ 'obc ⊆ Blacks 'M ) ) ∧ 'bc = { } }
        { 'Ma := 'M,, 'q := 'Queue } ;;
        Mul-Count n ;;
        { Mul-Propser n ∧ Roots ⊆ Blacks 'M
          ∧ 'obc ⊆ Blacks 'Ma ∧ Blacks 'Ma ⊆ Blacks 'M ∧ 'bc ⊆ Blacks 'M
          ∧ length 'Ma = length 'M ∧ Blacks 'Ma ⊆ 'bc
          ∧ ( 'Safe ∨ 'obc ⊆ Blacks 'Ma ∨ 'l < 'q ∧ ( 'q ≤ 'Queue ∨ 'obc ⊆ Blacks 'M ) )
          ∧ 'q < n + 1 }
        IF 'obc = 'bc THEN
          { Mul-Propser n ∧ Roots ⊆ Blacks 'M
            ∧ 'obc ⊆ Blacks 'Ma ∧ Blacks 'Ma ⊆ Blacks 'M ∧ 'bc ⊆ Blacks 'M
            ∧ length 'Ma = length 'M ∧ Blacks 'Ma ⊆ 'bc
            ∧ ( 'Safe ∨ 'obc ⊆ Blacks 'Ma ∨ 'l < 'q ∧ ( 'q ≤ 'Queue ∨ 'obc ⊆ Blacks 'M ) )
            ∧ 'q < n + 1 ∧ 'obc = 'bc }
          'l := 'l + 1
        ELSE { Mul-Propser n ∧ Roots ⊆ Blacks 'M
              ∧ 'obc ⊆ Blacks 'Ma ∧ Blacks 'Ma ⊆ Blacks 'M ∧ 'bc ⊆ Blacks 'M
              ∧ length 'Ma = length 'M ∧ Blacks 'Ma ⊆ 'bc
              ∧ ( 'Safe ∨ 'obc ⊆ Blacks 'Ma ∨ 'l < 'q ∧ ( 'q ≤ 'Queue ∨ 'obc ⊆ Blacks 'M ) )
              ∧ 'q < n + 1 ∧ 'obc ≠ 'bc }
            'l := 0 FI
    OD ;;
  Mul-Append n
OD

```

lemmas *mul-modules* = *Mul-Redirect-Edge-def* *Mul-Color-Target-def*
Mul-Blacken-Roots-def *Mul-Propagate-Black-def*
Mul-Count-def *Mul-Append-def*

lemma *Mul-Collector*:
 \vdash *Mul-Collector* *n*
 $\{\{False\}\}$
apply (*unfold* *Mul-Collector-def*)
apply *annhoare*
apply (*simp-all* *only:pre.simps* *Mul-Blacken-Roots*
Mul-Propagate-Black *Mul-Count* *Mul-Append*)
apply (*simp-all* *add:mul-modules*)
apply (*simp-all* *add:mul-collector-defs* *Queue-def*)
apply *force*
apply *force*
apply *force*
apply (*force* *simp* *add:less-Suc-eq-le*)
apply *force*
apply (*force* *dest:subset-antisym*)
apply *force*
apply *force*
apply *force*
done

2.3.3 Interference Freedom

lemma *le-length-filter-update* [*rule-format*]:
 $\forall i. (\neg P (list!i) \vee P j) \wedge i < \text{length } list$
 $\longrightarrow \text{length}(\text{filter } P \text{ } list) \leq \text{length}(\text{filter } P (list[i:=j]))$
apply (*induct-tac* *list*)
apply (*simp*)
apply (*clarify*)
apply (*case-tac* *i*)
apply (*simp*)
apply (*simp*)
done

lemma *less-length-filter-update* [*rule-format*]:
 $\forall i. P j \wedge \neg(P (list!i)) \wedge i < \text{length } list$
 $\longrightarrow \text{length}(\text{filter } P \text{ } list) < \text{length}(\text{filter } P (list[i:=j]))$
apply (*induct-tac* *list*)
apply (*simp*)
apply (*clarify*)
apply (*case-tac* *i*)
apply (*simp*)
apply (*simp*)
done

lemma *Mul-interfree-Blacken-Roots-Redirect-Edge*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
interfree-aux (Some(*Mul-Blacken-Roots* n), {}, Some(*Mul-Redirect-Edge* j n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply *safe*
apply(*simp-all add: Graph6 Graph9 Graph12 nth-list-update mul-mutator-defs mul-collector-defs*)
done

lemma *Mul-interfree-Redirect-Edge-Blacken-Roots*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
interfree-aux (Some(*Mul-Redirect-Edge* j n), {}, Some (*Mul-Blacken-Roots* n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply *safe*
apply(*simp-all add:mul-mutator-defs nth-list-update*)
done

lemma *Mul-interfree-Blacken-Roots-Color-Target*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
interfree-aux (Some(*Mul-Blacken-Roots* n), {}, Some (*Mul-Color-Target* j n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply *safe*
apply(*simp-all add:mul-mutator-defs mul-collector-defs nth-list-update Graph7 Graph8 Graph9 Graph12*)
done

lemma *Mul-interfree-Color-Target-Blacken-Roots*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
interfree-aux (Some(*Mul-Color-Target* j n), {}, Some (*Mul-Blacken-Roots* n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply *safe*
apply(*simp-all add:mul-mutator-defs nth-list-update*)
done

lemma *Mul-interfree-Propagate-Black-Redirect-Edge*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
interfree-aux (Some(*Mul-Propagate-Black* n), {}, Some (*Mul-Redirect-Edge* j n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply(*simp-all add:mul-mutator-defs mul-collector-defs Mul-PBInv-def nth-list-update Graph6*)
— 7 subgoals left
apply *clarify*
apply(*disjE-tac*)
apply(*simp-all add:Graph6*)
apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
apply(*rule conjI*)
apply(*rule impI,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*rule impI,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)

— 6 subgoals left
apply *clarify*
apply(*disjE-tac*)
 apply(*simp-all add:Graph6*)
 apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
apply(*rule conjI*)
 apply(*rule impI,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*rule impI,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
— 5 subgoals left
apply *clarify*
apply(*disjE-tac*)
 apply(*simp-all add:Graph6*)
 apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
apply(*rule conjI*)
 apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule less-le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule less-le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*erule conjE*)
apply(*case-tac M x!(T (Muts x!j))=Black*)
apply(*rule conjI*)
 apply(*rule impI,(rule disjI2)+,rule conjI*)
 apply *clarify*
 apply(*case-tac R (Muts x! j)=i*)
 apply (*force simp add: nth-list-update BtoW-def*)
 apply (*force simp add: nth-list-update*)
 apply(*erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
 apply(*rule impI,(rule disjI2)+, erule le-trans*)
 apply(*force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*rule conjI*)
 apply(*rule impI,rule disjI2,rule disjI2,rule disjI1, erule le-less-trans*)
 apply(*force simp add:Queue-def less-Suc-eq-le less-length-filter-update*)
 apply(*rule impI,rule disjI2,rule disjI2,rule disjI1, erule le-less-trans*)
 apply(*force simp add:Queue-def less-Suc-eq-le less-length-filter-update*)
— 4 subgoals left
apply *clarify*
apply(*disjE-tac*)
 apply(*simp-all add:Graph6*)
 apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
apply(*rule conjI*)
 apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule less-le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule less-le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*erule conjE*)
apply(*case-tac M x!(T (Muts x!j))=Black*)
 apply(*rule conjI*)

```

apply(rule impI,(rule disjI2)+,rule conjI)
apply clarify
apply(case-tac R (Muts x! j)=i)
  apply (force simp add: nth-list-update BtoW-def)
  apply (force simp add: nth-list-update)
apply(erule le-trans,force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(rule impI,(rule disjI2)+, erule le-trans)
apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(rule conjI)
apply(rule impI,rule disjI2,rule disjI2,rule disjI1, erule le-less-trans)
apply(force simp add: Queue-def less-Suc-eq-le less-length-filter-update)
apply(rule impI,rule disjI2,rule disjI2,rule disjI1, erule le-less-trans)
apply(force simp add: Queue-def less-Suc-eq-le less-length-filter-update)
— 3 subgoals left
apply clarify
apply(disjE-tac)
  apply(simp-all add: Graph6)
  apply (rule impI)
  apply(rule conjI)
    apply(rule disjI1,rule subset-trans,erule Graph3,simp,simp)
    apply(case-tac R (Muts x ! j)= ind x)
      apply(simp add: nth-list-update)
      apply(simp add: nth-list-update)
    apply(case-tac R (Muts x ! j)= ind x)
      apply(simp add: nth-list-update)
      apply(simp add: nth-list-update)
apply(case-tac M x!(T (Muts x!j))=Black)
apply(rule conjI)
apply(rule impI)
apply(rule conjI)
  apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
  apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(case-tac R (Muts x ! j)= ind x)
  apply(simp add: nth-list-update)
  apply(simp add: nth-list-update)
apply(rule impI)
apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(rule conjI)
apply(rule impI)
  apply(rule conjI)
    apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
    apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
  apply(case-tac R (Muts x ! j)= ind x)
    apply(simp add: nth-list-update)
    apply(simp add: nth-list-update)
apply(rule impI)
apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)

```

```

apply(erule conjE)
apply(rule conjI)
apply(case-tac M x!(T (Muts x!j))=Black)
apply(rule impI,rule conjI,(rule disjI2)+,rule conjI)
apply clarify
apply(case-tac R (Muts x! j)=i)
apply (force simp add: nth-list-update BtoW-def)
apply (force simp add: nth-list-update)
apply(erule le-trans,force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(case-tac R (Muts x ! j)= ind x)
apply(simp add:nth-list-update)
apply(simp add:nth-list-update)
apply(rule impI,rule conjI)
apply(rule disjI2,rule disjI2,rule disjI1, erule le-less-trans)
apply(force simp add: Queue-def less-Suc-eq-le less-length-filter-update)
apply(case-tac R (Muts x! j)=ind x)
apply (force simp add: nth-list-update)
apply (force simp add: nth-list-update)
apply(rule impI, (rule disjI2)+, erule le-trans)
apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
— 2 subgoals left
apply clarify
apply(rule conjI)
apply(disjE-tac)
apply(simp-all add:Mul-Auxk-def Graph6)
apply (rule impI)
apply(rule conjI)
apply(rule disjI1,rule subset-trans,erule Graph3,simp,simp)
apply(case-tac R (Muts x ! j)= ind x)
apply(simp add:nth-list-update)
apply(simp add:nth-list-update)
apply(case-tac R (Muts x ! j)= ind x)
apply(simp add:nth-list-update)
apply(simp add:nth-list-update)
apply(case-tac M x!(T (Muts x!j))=Black)
apply(rule impI)
apply(rule conjI)
apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(case-tac R (Muts x ! j)= ind x)
apply(simp add:nth-list-update)
apply(simp add:nth-list-update)
apply(rule impI)
apply(rule conjI)
apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(case-tac R (Muts x ! j)= ind x)
apply(simp add:nth-list-update)
apply(simp add:nth-list-update)

```

```

apply(rule impI)
apply(rule conjI)
apply(erule conjE)+
apply(case-tac M x!(T (Muts x!j))=Black)
apply((rule disjI2)+,rule conjI)
apply clarify
apply(case-tac R (Muts x! j)=i)
  apply (force simp add: nth-list-update BtoW-def)
apply (force simp add: nth-list-update)
apply(rule conjI)
apply(erule le-trans,force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(rule impI)
apply(case-tac R (Muts x ! j)= ind x)
  apply(simp add:nth-list-update BtoW-def)
apply (simp add:nth-list-update)
apply(rule impI)
apply simp
apply(disjE-tac)
  apply(rule disjI1, erule less-le-trans)
  apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply force
apply(rule disjI2,rule disjI2,rule disjI1, erule le-less-trans)
apply(force simp add: Queue-def less-Suc-eq-le less-length-filter-update)
apply(case-tac R (Muts x ! j)= ind x)
  apply(simp add:nth-list-update)
apply(simp add:nth-list-update)
apply(disjE-tac)
apply simp-all
apply(conjI-tac)
  apply(rule impI)
  apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
  apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(erule conjE)+
apply(rule impI,(rule disjI2)+,rule conjI)
  apply(erule le-trans,force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply(rule impI)+
apply simp
apply(disjE-tac)
  apply(rule disjI1, erule less-le-trans)
  apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update)
apply force
— 1 subgoal left
apply clarify
apply(disjE-tac)
  apply(simp-all add: Graph6)
  apply(rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp)
apply(rule conjI)
  apply(rule impI,rule disjI2,rule disjI2,rule disjI1,erule less-le-trans,force simp
add: Queue-def less-Suc-eq-le le-length-filter-update)

```



```

apply(rule impI,rule disjI2,rule disjI2,rule disjI1,erule less-le-trans,force simp add:Queue-def
less-Suc-eq-le le-length-filter-update)
apply(erule conjE)
apply(case-tac M x!(T (Muts x!j))=Black)
apply(rule conjI)
apply(rule impI,(rule disjI2)+,rule conjI)
apply clarify
apply(case-tac R (Muts x! j)=i)
apply (force simp add: nth-list-update BtoW-def)
apply (force simp add: nth-list-update)
apply(erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update)
apply(rule impI,(rule disjI2)+, erule le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update)
apply(rule conjI)
apply(rule impI,rule disjI2,rule disjI2,rule disjI1, erule le-less-trans)
apply(force simp add:Queue-def less-Suc-eq-le less-length-filter-update)
apply(rule impI,rule disjI2,rule disjI2,rule disjI1, erule le-less-trans)
apply(force simp add:Queue-def less-Suc-eq-le less-length-filter-update)
done

```

```

lemma Mul-interfree-Redirect-Edge-Propagate-Black:  $\llbracket 0 \leq j; j < n \rrbracket \implies$ 
interfree-aux (Some (Mul-Redirect-Edge j n ), {}, Some (Mul-Propagate-Black n))
apply (unfold mul-modules)
apply interfree-aux
apply safe
apply(simp-all add:mul-mutator-defs nth-list-update)
done

```

```

lemma Mul-interfree-Propagate-Black-Color-Target:  $\llbracket 0 \leq j; j < n \rrbracket \implies$ 
interfree-aux (Some (Mul-Propagate-Black n), {}, Some (Mul-Color-Target j n ))
apply (unfold mul-modules)
apply interfree-aux
apply(simp-all add: mul-collector-defs mul-mutator-defs)
— 7 subgoals left
apply clarify
apply (simp add:Graph7 Graph8 Graph12)
apply(disjE-tac)
apply(simp add:Graph7 Graph8 Graph12)
apply(case-tac M x!(T (Muts x!j))=Black)
apply(rule disjI2,rule disjI1, erule le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply((rule disjI2)+,erule subset-psubset-trans, erule Graph11, simp)
apply((rule disjI2)+,erule psubset-subset-trans, simp add: Graph9)
— 6 subgoals left
apply clarify
apply (simp add:Graph7 Graph8 Graph12)
apply(disjE-tac)
apply(simp add:Graph7 Graph8 Graph12)
apply(case-tac M x!(T (Muts x!j))=Black)

```

```

apply(rule disjI2,rule disjI1, erule le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply((rule disjI2)+,erule subset-psubset-trans, erule Graph11, simp)
apply((rule disjI2)+,erule psubset-subset-trans, simp add: Graph9)
— 5 subgoals left
apply clarify
apply (simp add:mul-collector-defs Mul-PBInv-def Graph7 Graph8 Graph12)
apply(disjE-tac)
  apply(simp add:Graph7 Graph8 Graph12)
  apply(rule disjI2,rule disjI1, erule psubset-subset-trans,simp add:Graph9)
  apply(case-tac M x!(T (Muts x!j))=Black)
  apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
  apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
  apply(rule disjI2,rule disjI1,erule subset-psubset-trans, erule Graph11, simp)
apply(erule conjE)
apply(case-tac M x!(T (Muts x!j))=Black)
apply((rule disjI2)+)
apply (rule conjI)
  apply(simp add:Graph10)
apply(erule le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply(rule disjI2,rule disjI1,erule subset-psubset-trans, erule Graph11, simp)
— 4 subgoals left
apply clarify
apply (simp add:mul-collector-defs Mul-PBInv-def Graph7 Graph8 Graph12)
apply(disjE-tac)
  apply(simp add:Graph7 Graph8 Graph12)
  apply(rule disjI2,rule disjI1, erule psubset-subset-trans,simp add:Graph9)
  apply(case-tac M x!(T (Muts x!j))=Black)
  apply(rule disjI2,rule disjI2,rule disjI1, erule less-le-trans)
  apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
  apply(rule disjI2,rule disjI1,erule subset-psubset-trans, erule Graph11, simp)
apply(erule conjE)
apply(case-tac M x!(T (Muts x!j))=Black)
apply((rule disjI2)+)
apply (rule conjI)
  apply(simp add:Graph10)
apply(erule le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply(rule disjI2,rule disjI1,erule subset-psubset-trans, erule Graph11, simp)
— 3 subgoals left
apply clarify
apply (simp add:mul-collector-defs Mul-PBInv-def Graph7 Graph8 Graph12)
apply(case-tac M x!(T (Muts x!j))=Black)
apply(simp add:Graph10)
apply(disjE-tac)
apply simp-all
apply(rule disjI2, rule disjI2, rule disjI1,erule less-le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)

```

```

apply(erule conjE)
apply((rule disjI2)+,erule le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply(rule conjI)
apply(rule disjI2,rule disjI1, erule subset-psubset-trans,simp add:Graph11)
apply (force simp add:nth-list-update)
— 2 subgoals left
apply clarify
apply(simp add:Mul-Auxk-def Graph7 Graph8 Graph12)
apply(case-tac M x!(T (Muts x!j))=Black)
apply(simp add:Graph10)
apply(disjE-tac)
apply simp-all
apply(rule disjI2, rule disjI2, rule disjI1,erule less-le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply(erule conjE)+
apply((rule disjI2)+,rule conjI, erule le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply((rule impI)+)
apply simp
apply(erule disjE)
apply(rule disjI1, erule less-le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply force
apply(rule conjI)
apply(rule disjI2,rule disjI1, erule subset-psubset-trans,simp add:Graph11)
apply (force simp add:nth-list-update)
— 1 subgoal left
apply clarify
apply (simp add:mul-collector-defs Mul-PBInv-def Graph7 Graph8 Graph12)
apply(case-tac M x!(T (Muts x!j))=Black)
apply(simp add:Graph10)
apply(disjE-tac)
apply simp-all
apply(rule disjI2, rule disjI2, rule disjI1,erule less-le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply(erule conjE)
apply((rule disjI2)+,erule le-trans)
apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
apply(rule disjI2,rule disjI1, erule subset-psubset-trans,simp add:Graph11)
done

```

```

lemma Mul-interfree-Color-Target-Propagate-Black:  $\llbracket 0 \leq j; j < n \rrbracket \implies$ 
  interfree-aux (Some(Mul-Color-Target j n),{\},Some(Mul-Propagate-Black n ))
apply (unfold mul-modules)
apply interfree-aux
apply safe
apply(simp-all add:mul-mutator-defs nth-list-update)
done

```

lemma *Mul-interfree-Count-Redirect-Edge*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
interfree-aux (Some(Mul-Count n), {}, Some(Mul-Redirect-Edge j n))
apply (*unfold mul-modules*)
apply *interfree-aux*
— 9 subgoals left
apply(*simp add:mul-mutator-defs mul-collector-defs Mul-CountInv-def Graph6*)
apply *clarify*
apply *disjE-tac*
 apply(*simp add:Graph6*)
 apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
 apply(*simp add:Graph6*)
apply *clarify*
apply *disjE-tac*
 apply(*simp add:Graph6*)
 apply(*rule conjI*)
 apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
 apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*simp add:Graph6*)
— 8 subgoals left
apply(*simp add:mul-mutator-defs nth-list-update*)
— 7 subgoals left
apply(*simp add:mul-mutator-defs mul-collector-defs*)
apply *clarify*
apply *disjE-tac*
 apply(*simp add:Graph6*)
 apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
 apply(*simp add:Graph6*)
apply *clarify*
apply *disjE-tac*
 apply(*simp add:Graph6*)
 apply(*rule conjI*)
 apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
 apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*simp add:Graph6*)
— 6 subgoals left
apply(*simp add:mul-mutator-defs mul-collector-defs Mul-CountInv-def*)
apply *clarify*
apply *disjE-tac*
 apply(*simp add:Graph6 Queue-def*)
 apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
 apply(*simp add:Graph6*)
apply *clarify*
apply *disjE-tac*
 apply(*simp add:Graph6*)

apply(*rule conjI*)
apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*simp add:Graph6*)
— 5 subgoals left
apply(*simp add:mul-mutator-defs mul-collector-defs Mul-CountInv-def*)
apply clarify
apply disjE-tac
apply(*simp add:Graph6*)
apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
apply(*simp add:Graph6*)
apply clarify
apply disjE-tac
apply(*simp add:Graph6*)
apply(*rule conjI*)
apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*simp add:Graph6*)
— 4 subgoals left
apply(*simp add:mul-mutator-defs mul-collector-defs Mul-CountInv-def*)
apply clarify
apply disjE-tac
apply(*simp add:Graph6*)
apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
apply(*simp add:Graph6*)
apply clarify
apply disjE-tac
apply(*simp add:Graph6*)
apply(*rule conjI*)
apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*rule impI,rule disjI2,rule disjI2,rule disjI1,erule le-trans,force simp add:Queue-def less-Suc-eq-le le-length-filter-update*)
apply(*simp add:Graph6*)
— 3 subgoals left
apply(*simp add:mul-mutator-defs nth-list-update*)
— 2 subgoals left
apply(*simp add:mul-mutator-defs mul-collector-defs Mul-CountInv-def*)
apply clarify
apply disjE-tac
apply(*simp add:Graph6*)
apply(*rule impI,rule disjI1,rule subset-trans,erule Graph3,simp,simp*)
apply(*simp add:Graph6*)
apply clarify
apply disjE-tac

```

apply(simp add: Graph6)
apply(rule conjI)
apply(rule impI, rule disjI2, rule disjI2, rule disjI1, erule le-trans, force simp add: Queue-def
less-Suc-eq-le le-length-filter-update)
apply(rule impI, rule disjI2, rule disjI2, rule disjI1, erule le-trans, force simp add: Queue-def
less-Suc-eq-le le-length-filter-update)
apply(simp add: Graph6)
— 1 subgoal left
apply(simp add: mul-mutator-defs nth-list-update)
done

```

```

lemma Mul-interfree-Redirect-Edge-Count:  $\llbracket 0 \leq j; j < n \rrbracket \implies$ 
  interfree-aux (Some(Mul-Redirect-Edge j n), {}, Some(Mul-Count n ))
apply (unfold mul-modules)
apply interfree-aux
apply safe
apply(simp-all add: mul-mutator-defs nth-list-update)
done

```

```

lemma Mul-interfree-Count-Color-Target:  $\llbracket 0 \leq j; j < n \rrbracket \implies$ 
  interfree-aux (Some(Mul-Count n ), {}, Some(Mul-Color-Target j n ))
apply (unfold mul-modules)
apply interfree-aux
apply(simp-all add: mul-collector-defs mul-mutator-defs Mul-CountInv-def)
— 6 subgoals left
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply (simp add: Graph7 Graph8 Graph12)
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply(case-tac M x!(T (Muts x!j))=Black)
  apply(rule disjI2, rule disjI2, rule disjI1, erule le-trans)
  apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update Graph10)
  apply((rule disjI2)+, (erule subset-psubset-trans)+, simp add: Graph11)
apply (simp add: Graph7 Graph8 Graph12)
apply((rule disjI2)+, erule psubset-subset-trans, simp add: Graph9)
— 5 subgoals left
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply (simp add: Graph7 Graph8 Graph12)
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply(case-tac M x!(T (Muts x!j))=Black)
  apply(rule disjI2, rule disjI2, rule disjI1, erule le-trans)
  apply(force simp add: Queue-def less-Suc-eq-le le-length-filter-update Graph10)

```

```

  apply((rule disjI2)+,(erule subset-psubset-trans)+, simp add: Graph11)
apply (simp add: Graph7 Graph8 Graph12)
apply((rule disjI2)+,erule psubset-subset-trans, simp add: Graph9)
— 4 subgoals left
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply (simp add: Graph7 Graph8 Graph12)
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply(case-tac M x!(T (Muts x!j))=Black)
  apply(rule disjI2,rule disjI2, rule disjI1, erule le-trans)
  apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
  apply((rule disjI2)+,(erule subset-psubset-trans)+, simp add: Graph11)
apply (simp add: Graph7 Graph8 Graph12)
apply((rule disjI2)+,erule psubset-subset-trans, simp add: Graph9)
— 3 subgoals left
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply (simp add: Graph7 Graph8 Graph12)
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply(case-tac M x!(T (Muts x!j))=Black)
  apply(rule disjI2,rule disjI2, rule disjI1, erule le-trans)
  apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
  apply((rule disjI2)+,(erule subset-psubset-trans)+, simp add: Graph11)
apply (simp add: Graph7 Graph8 Graph12)
apply((rule disjI2)+,erule psubset-subset-trans, simp add: Graph9)
— 2 subgoals left
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12 nth-list-update)
  apply (simp add: Graph7 Graph8 Graph12 nth-list-update)
apply clarify
apply disjE-tac
  apply (simp add: Graph7 Graph8 Graph12)
  apply(rule conjI)
  apply(case-tac M x!(T (Muts x!j))=Black)
  apply(rule disjI2,rule disjI2, rule disjI1, erule le-trans)
  apply(force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10)
  apply((rule disjI2)+,(erule subset-psubset-trans)+, simp add: Graph11)
  apply (simp add: nth-list-update)
apply (simp add: Graph7 Graph8 Graph12)
apply(rule conjI)
  apply((rule disjI2)+,erule psubset-subset-trans, simp add: Graph9)
apply (simp add: nth-list-update)

```

— 1 subgoal left
apply *clarify*
apply *disjE-tac*
 apply (*simp add: Graph7 Graph8 Graph12*)
 apply (*simp add: Graph7 Graph8 Graph12*)
apply *clarify*
apply *disjE-tac*
 apply (*simp add: Graph7 Graph8 Graph12*)
 apply(*case-tac M x!(T (Muts x!j))=Black*)
 apply(*rule disjI2,rule disjI2, rule disjI1, erule le-trans*)
 apply(*force simp add:Queue-def less-Suc-eq-le le-length-filter-update Graph10*)
 apply((*rule disjI2*)+,(*erule subset-psubset-trans*)+, *simp add: Graph11*)
apply (*simp add: Graph7 Graph8 Graph12*)
apply((*rule disjI2*)+,*erule psubset-subset-trans, simp add: Graph9*)
done

lemma *Mul-interfree-Color-Target-Count*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
 interfree-aux (Some(Mul-Color-Target j n),{\}, Some(Mul-Count n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply *safe*
apply(*simp-all add:mul-mutator-defs nth-list-update*)
done

lemma *Mul-interfree-Append-Redirect-Edge*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
 interfree-aux (Some(Mul-Append n),{\}, Some(Mul-Redirect-Edge j n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply(*tactic <ALLGOALS (clarify-tac context)>*)
apply(*simp-all add:Graph6 Append-to-free0 Append-to-free1 mul-collector-defs mul-mutator-defs*
 Mul-AppendInv-def)
apply(*erule-tac x=j in alle, force dest:Graph3*)+
done

lemma *Mul-interfree-Redirect-Edge-Append*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
 interfree-aux (Some(Mul-Redirect-Edge j n),{\},Some(Mul-Append n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply(*tactic <ALLGOALS (clarify-tac context)>*)
apply(*simp-all add:mul-collector-defs Append-to-free0 Mul-AppendInv-def mul-mutator-defs*
 nth-list-update)
done

lemma *Mul-interfree-Append-Color-Target*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
 interfree-aux (Some(Mul-Append n),{\}, Some(Mul-Color-Target j n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply(*tactic <ALLGOALS (clarify-tac context)>*)
apply(*simp-all add:mul-mutator-defs mul-collector-defs Mul-AppendInv-def Graph7*)

Graph8 Append-to-free0 Append-to-free1
Graph12 nth-list-update)

done

lemma *Mul-interfree-Color-Target-Append*: $\llbracket 0 \leq j; j < n \rrbracket \implies$
interfree-aux (Some (Mul-Color-Target j n), {}, Some (Mul-Append n))
apply (*unfold mul-modules*)
apply *interfree-aux*
apply (*tactic <ALLGOALS (clarify-tac context)>*)
apply (*simp-all add: mul-mutator-defs nth-list-update*)
apply (*simp add: Mul-AppendInv-def Append-to-free0*)
done

Interference freedom Collector-Mutator

lemmas *mul-collector-mutator-interfree =*
Mul-interfree-Blacken-Roots-Redirect-Edge Mul-interfree-Blacken-Roots-Color-Target
Mul-interfree-Propagate-Black-Redirect-Edge Mul-interfree-Propagate-Black-Color-Target
Mul-interfree-Count-Redirect-Edge Mul-interfree-Count-Color-Target
Mul-interfree-Append-Redirect-Edge Mul-interfree-Append-Color-Target
Mul-interfree-Redirect-Edge-Blacken-Roots Mul-interfree-Color-Target-Blacken-Roots
Mul-interfree-Redirect-Edge-Propagate-Black Mul-interfree-Color-Target-Propagate-Black
Mul-interfree-Redirect-Edge-Count Mul-interfree-Color-Target-Count
Mul-interfree-Redirect-Edge-Append Mul-interfree-Color-Target-Append

lemma *Mul-interfree-Collector-Mutator*: $j < n \implies$
interfree-aux (Some (Mul-Collector n), {}, Some (Mul-Mutator j n))
apply (*unfold Mul-Collector-def Mul-Mutator-def*)
apply *interfree-aux*
apply (*simp-all add: mul-collector-mutator-interfree*)
apply (*unfold mul-modules mul-collector-defs mul-mutator-defs*)
apply (*tactic <TRYALL (interfree-aux-tac context)>*)
— 42 subgoals left
apply (*clarify, simp add: Graph6 Graph7 Graph8 Append-to-free0 Append-to-free1*
Graph12)+)
— 24 subgoals left
apply (*simp-all add: Graph6 Graph7 Graph8 Append-to-free0 Append-to-free1 Graph12*)
— 14 subgoals left
apply (*tactic <TRYALL (clarify-tac context)>*)
apply (*simp-all add: Graph6 Graph7 Graph8 Append-to-free0 Append-to-free1 Graph12*)
apply (*tactic <TRYALL (resolve-tac context [conjI])>*)
apply (*tactic <TRYALL (resolve-tac context [impI])>*)
apply (*tactic <TRYALL (eresolve-tac context [disjE])>*)
apply (*tactic <TRYALL (eresolve-tac context [conjE])>*)
apply (*tactic <TRYALL (eresolve-tac context [disjE])>*)
apply (*tactic <TRYALL (eresolve-tac context [disjE])>*)
— 72 subgoals left
apply (*simp-all add: Graph6 Graph7 Graph8 Append-to-free0 Append-to-free1 Graph12*)
— 35 subgoals left

apply(*tactic* ‹*TRYALL*(*EVERY*'[*resolve-tac context* [*disjI1*],
resolve-tac context [*subset-trans*],
eresolve-tac context @{*thms Graph3*},
force-tac context,
assume-tac context])›)

— 28 subgoals left

apply(*tactic* ‹*TRYALL* (*eresolve-tac context* [*conjE*])›)

apply(*tactic* ‹*TRYALL* (*eresolve-tac context* [*disjE*])›)

— 34 subgoals left

apply(*rule disjI2*,*rule disjI1*,*erule le-trans*,*force simp add:Queue-def less-Suc-eq-le*
le-length-filter-update)

apply(*rule disjI2*,*rule disjI1*,*erule le-trans*,*force simp add:Queue-def less-Suc-eq-le*
le-length-filter-update)

apply(*case-tac* [!] *M x!(T (Muts x ! j))=Black*)

apply(*simp-all add:Graph10*)

— 47 subgoals left

apply(*tactic* ‹*TRYALL*(*EVERY*'[*REPEAT o resolve-tac context* [*disjI2*],
eresolve-tac context @{*thms subset-psubset-trans*},
eresolve-tac context @{*thms Graph11*},
force-tac context])›)

— 41 subgoals left

apply(*tactic* ‹*TRYALL*(*EVERY*'[*resolve-tac context* [*disjI2*],
resolve-tac context [*disjI1*],
eresolve-tac context @{*thms le-trans*},
force-tac (context addsimps @{thms Queue-def less-Suc-eq-le le-length-filter-update})])›)

— 35 subgoals left

apply(*tactic* ‹*TRYALL*(*EVERY*'[*resolve-tac context* [*disjI2*],
resolve-tac context [*disjI1*],
eresolve-tac context @{*thms psubset-subset-trans*},
resolve-tac context @{*thms Graph9*},
force-tac context])›)

— 31 subgoals left

apply(*tactic* ‹*TRYALL*(*EVERY*'[*resolve-tac context* [*disjI2*],
resolve-tac context [*disjI1*],
eresolve-tac context @{*thms subset-psubset-trans*},
eresolve-tac context @{*thms Graph11*},
force-tac context])›)

— 29 subgoals left

apply(*tactic* ‹*TRYALL*(*EVERY*'[*REPEAT o resolve-tac context* [*disjI2*],
eresolve-tac context @{*thms subset-psubset-trans*},
eresolve-tac context @{*thms subset-psubset-trans*},
eresolve-tac context @{*thms Graph11*},
force-tac context])›)

— 25 subgoals left

apply(*tactic* ‹*TRYALL*(*EVERY*'[*resolve-tac context* [*disjI2*],
resolve-tac context [*disjI2*],
resolve-tac context [*disjI1*],
eresolve-tac context @{*thms le-trans*},
force-tac (context addsimps @{thms Queue-def less-Suc-eq-le le-length-filter-update})])›)

— 10 subgoals left
apply(rule *disjI2*,rule *disjI2*,rule *conjI*,erule *less-le-trans*,force *simp add:Queue-def less-Suc-eq-le le-length-filter-update*, rule *disjI1*, rule *less-imp-le*, erule *less-le-trans*, force *simp add:Queue-def less-Suc-eq-le le-length-filter-update*)+
done

Interference freedom Mutator-Collector

lemma *Mul-interfree-Mutator-Collector*: $j < n \implies$
interfree-aux (Some (Mul-Mutator j n), {}, Some (Mul-Collector n))
apply(unfold *Mul-Collector-def Mul-Mutator-def*)
apply *interfree-aux*
apply(*simp-all add:mul-collector-mutator-interfree*)
apply(unfold *mul-modules mul-collector-defs mul-mutator-defs*)
apply(tactic ‹TRYALL (*interfree-aux-tac context*)›)
— 76 subgoals left
apply (*clarsimp simp add: nth-list-update*)+
— 56 subgoals left
apply (*clarsimp simp add: Mul-AppendInv-def Append-to-free0 nth-list-update*)+
done

The Multi-Mutator Garbage Collection Algorithm

The total number of verification conditions is 328

lemma *Mul-Gar-Coll*:
 $\| - \{ \text{Mul-Proper } n \wedge \text{Mul-mut-init } n \wedge (\forall i < n. Z (\text{Muts!}i)) \}$
COBEGIN
Mul-Collector n
 $\{ \text{False} \}$
 $\|$
SCHEME $[0 \leq j < n]$
Mul-Mutator j n
 $\{ \text{False} \}$
COEND
 $\{ \text{False} \}$
apply *oghoare*
— Strengthening the precondition
apply(rule *Int-greatest*)
apply (*case-tac n*)
apply(force *simp add: Mul-Collector-def mul-mutator-defs mul-collector-defs nth-append*)
apply(*simp add: Mul-Mutator-def mul-collector-defs mul-mutator-defs nth-append*)
apply *force*
apply *clarify*
apply(*case-tac i*)
apply(*simp add: Mul-Collector-def mul-mutator-defs mul-collector-defs nth-append*)
apply(*simp add: Mul-Mutator-def mul-mutator-defs mul-collector-defs nth-append nth-map-upt*)
— Collector
apply(rule *Mul-Collector*)

```
— Mutator
apply(erule Mul-Mutator)
— Interference freedom
apply(simp add:Mul-interfree-Collector-Mutator)
apply(simp add:Mul-interfree-Mutator-Collector)
apply(simp add:Mul-interfree-Mutator-Mutator)
— Weakening of the postcondition
apply(case-tac n)
  apply(simp add:Mul-Collector-def mul-mutator-defs mul-collector-defs nth-append)
apply(simp add:Mul-Mutator-def mul-mutator-defs mul-collector-defs nth-append)
done

end
```

Chapter 3

The Rely-Guarantee Method

3.1 Abstract Syntax

```
theory RG-Com imports Main begin
```

Semantics of assertions and boolean expressions (*bexp*) as sets of states.
Syntax of commands *com* and parallel commands *par-com*.

```
type-synonym 'a bexp = 'a set
```

```
datatype 'a com =  
  Basic 'a ⇒ 'a  
  | Seq 'a com 'a com  
  | Cond 'a bexp 'a com 'a com  
  | While 'a bexp 'a com  
  | Await 'a bexp 'a com
```

```
type-synonym 'a par-com = 'a com option list
```

```
end
```

3.2 Operational Semantics

```
theory RG-Tran  
imports RG-Com  
begin
```

3.2.1 Semantics of Component Programs

Environment transitions

```
type-synonym 'a conf = (('a com) option) × 'a
```

inductive-set

```
etran :: ('a conf × 'a conf) set  
and etran' :: 'a conf ⇒ 'a conf ⇒ bool (⟦-⟧ ⇒ [81,81] 80)
```

where

$P -e\rightarrow Q \equiv (P, Q) \in etran$
| $Env: (P, s) -e\rightarrow (P, t)$

lemma $etranE: c -e\rightarrow c' \implies (\bigwedge P s t. c = (P, s) \implies c' = (P, t) \implies Q) \implies Q$
by ($induct\ c, induct\ c', erule\ etran.cases, blast$)

Component transitions

inductive-set

$ctran :: ('a\ conf \times 'a\ conf)\ set$
and $ctran' :: 'a\ conf \Rightarrow 'a\ conf \Rightarrow bool$ ($\langle - -c\rightarrow - \rangle [81,81]\ 80$)
and $ctrans :: 'a\ conf \Rightarrow 'a\ conf \Rightarrow bool$ ($\langle - -c*\rightarrow - \rangle [81,81]\ 80$)

where

$P -c\rightarrow Q \equiv (P, Q) \in ctran$
| $P -c*\rightarrow Q \equiv (P, Q) \in ctran^*$

| $Basic: (Some(Basic\ f), s) -c\rightarrow (None, f\ s)$

| $Seq1: (Some\ P0, s) -c\rightarrow (None, t) \implies (Some(Seq\ P0\ P1), s) -c\rightarrow (Some\ P1, t)$

| $Seq2: (Some\ P0, s) -c\rightarrow (Some\ P2, t) \implies (Some(Seq\ P0\ P1), s) -c\rightarrow (Some(Seq\ P2\ P1), t)$

| $CondT: s \in b \implies (Some(Cond\ b\ P1\ P2), s) -c\rightarrow (Some\ P1, s)$

| $CondF: s \notin b \implies (Some(Cond\ b\ P1\ P2), s) -c\rightarrow (Some\ P2, s)$

| $WhileF: s \notin b \implies (Some(While\ b\ P), s) -c\rightarrow (None, s)$

| $WhileT: s \in b \implies (Some(While\ b\ P), s) -c\rightarrow (Some(Seq\ P\ (While\ b\ P)), s)$

| $Await: \llbracket s \in b; (Some\ P, s) -c*\rightarrow (None, t) \rrbracket \implies (Some(Await\ b\ P), s) -c\rightarrow (None, t)$

monos $rtrancl\ mono$

3.2.2 Semantics of Parallel Programs

type-synonym $'a\ par\ conf = ('a\ par\ com) \times 'a$

inductive-set

$par\ etran :: ('a\ par\ conf \times 'a\ par\ conf)\ set$
and $par\ etran' :: ['a\ par\ conf, 'a\ par\ conf] \Rightarrow bool$ ($\langle - -pe\rightarrow - \rangle [81,81]\ 80$)

where

$P -pe\rightarrow Q \equiv (P, Q) \in par\ etran$
| $ParEnv: (Ps, s) -pe\rightarrow (Ps, t)$

inductive-set

$par\ ctran :: ('a\ par\ conf \times 'a\ par\ conf)\ set$
and $par\ ctran' :: ['a\ par\ conf, 'a\ par\ conf] \Rightarrow bool$ ($\langle - -pc\rightarrow - \rangle [81,81]\ 80$)

where

$P -pc \rightarrow Q \equiv (P, Q) \in \text{par-ctran}$
 $| \text{ParComp}: \llbracket i < \text{length } Ps; (Ps!i, s) -c \rightarrow (r, t) \rrbracket \Longrightarrow (Ps, s) -pc \rightarrow (Ps[i:=r], t)$

lemma *par-ctranE*: $c -pc \rightarrow c' \Longrightarrow$

$(\bigwedge i \text{ Ps } s \ r \ t. c = (Ps, s) \Longrightarrow c' = (Ps[i := r], t) \Longrightarrow i < \text{length } Ps \Longrightarrow$
 $(Ps ! i, s) -c \rightarrow (r, t) \Longrightarrow P) \Longrightarrow P$

by (*induct c, induct c', erule par-ctran.cases, blast*)

3.2.3 Computations

Sequential computations

type-synonym $'a \text{ confs} = 'a \text{ conf list}$

inductive-set $\text{cptn} :: 'a \text{ confs set}$

where

CptnOne: $\llbracket (P, s) \rrbracket \in \text{cptn}$
 $| \text{CptnEnv}$: $(P, t) \# xs \in \text{cptn} \Longrightarrow (P, s) \# (P, t) \# xs \in \text{cptn}$
 $| \text{CptnComp}$: $\llbracket (P, s) -c \rightarrow (Q, t); (Q, t) \# xs \in \text{cptn} \rrbracket \Longrightarrow (P, s) \# (Q, t) \# xs \in \text{cptn}$

definition $\text{cp} :: ('a \text{ com}) \text{ option} \Rightarrow 'a \Rightarrow ('a \text{ confs}) \text{ set}$ **where**

$\text{cp } P \ s \equiv \{l. !!0=(P, s) \wedge l \in \text{cptn}\}$

Parallel computations

type-synonym $'a \text{ par-conf} = 'a \text{ par-conf list}$

inductive-set $\text{par-cptn} :: 'a \text{ par-conf} \text{ set}$

where

ParCptnOne: $\llbracket (P, s) \rrbracket \in \text{par-cptn}$
 $| \text{ParCptnEnv}$: $(P, t) \# xs \in \text{par-cptn} \Longrightarrow (P, s) \# (P, t) \# xs \in \text{par-cptn}$
 $| \text{ParCptnComp}$: $\llbracket (P, s) -pc \rightarrow (Q, t); (Q, t) \# xs \in \text{par-cptn} \rrbracket \Longrightarrow (P, s) \# (Q, t) \# xs \in \text{par-cptn}$

definition $\text{par-cp} :: 'a \text{ par-com} \Rightarrow 'a \Rightarrow ('a \text{ par-conf}) \text{ set}$ **where**

$\text{par-cp } P \ s \equiv \{l. !!0=(P, s) \wedge l \in \text{par-cptn}\}$

3.2.4 Modular Definition of Computation

definition $\text{lift} :: 'a \text{ com} \Rightarrow 'a \text{ conf} \Rightarrow 'a \text{ conf}$ **where**

$\text{lift } Q \equiv \lambda(P, s). (\text{if } P = \text{None} \text{ then } (\text{Some } Q, s) \text{ else } (\text{Some}(\text{Seq}(\text{the } P) \ Q), s))$

inductive-set $\text{cptn-mod} :: ('a \text{ confs}) \text{ set}$

where

CptnModOne: $\llbracket (P, s) \rrbracket \in \text{cptn-mod}$
 $| \text{CptnModEnv}$: $(P, t) \# xs \in \text{cptn-mod} \Longrightarrow (P, s) \# (P, t) \# xs \in \text{cptn-mod}$
 $| \text{CptnModNone}$: $\llbracket (\text{Some } P, s) -c \rightarrow (\text{None}, t); (\text{None}, t) \# xs \in \text{cptn-mod} \rrbracket \Longrightarrow$
 $(\text{Some } P, s) \# (\text{None}, t) \# xs \in \text{cptn-mod}$

| *CptnModCondT*: $\llbracket (Some\ P0, s)\#ys \in\ cptn\text{-}mod; s \in\ b \rrbracket \implies (Some(Cond\ b\ P0\ P1), s)\#(Some\ P0, s)\#ys \in\ cptn\text{-}mod$
| *CptnModCondF*: $\llbracket (Some\ P1, s)\#ys \in\ cptn\text{-}mod; s \notin\ b \rrbracket \implies (Some(Cond\ b\ P0\ P1), s)\#(Some\ P1, s)\#ys \in\ cptn\text{-}mod$
| *CptnModSeq1*: $\llbracket (Some\ P0, s)\#xs \in\ cptn\text{-}mod; zs = map\ (lift\ P1)\ xs \rrbracket \implies (Some(Seq\ P0\ P1), s)\#zs \in\ cptn\text{-}mod$
| *CptnModSeq2*:
 $\llbracket (Some\ P0, s)\#xs \in\ cptn\text{-}mod; fst(last\ ((Some\ P0, s)\#xs)) = None;$
 $(Some\ P1, snd(last\ ((Some\ P0, s)\#xs)))\#ys \in\ cptn\text{-}mod;$
 $zs = (map\ (lift\ P1)\ xs)\@ys \rrbracket \implies (Some(Seq\ P0\ P1), s)\#zs \in\ cptn\text{-}mod$

| *CptnModWhile1*:
 $\llbracket (Some\ P, s)\#xs \in\ cptn\text{-}mod; s \in\ b; zs = map\ (lift\ (While\ b\ P))\ xs \rrbracket \implies (Some(While\ b\ P), s)\#(Some(Seq\ P\ (While\ b\ P)), s)\#zs \in\ cptn\text{-}mod$
| *CptnModWhile2*:
 $\llbracket (Some\ P, s)\#xs \in\ cptn\text{-}mod; fst(last\ ((Some\ P, s)\#xs)) = None; s \in\ b;$
 $zs = (map\ (lift\ (While\ b\ P))\ xs)\@ys;$
 $(Some(While\ b\ P), snd(last\ ((Some\ P, s)\#xs)))\#ys \in\ cptn\text{-}mod \rrbracket \implies (Some(While\ b\ P), s)\#(Some(Seq\ P\ (While\ b\ P)), s)\#zs \in\ cptn\text{-}mod$

3.2.5 Equivalence of Both Definitions.

lemma *last-length*: $((a\#xs)!(length\ xs)) = last\ (a\#xs)$
by (*induct xs*) *auto*

lemma *div-seq* [*rule-format*]: $list \in\ cptn\text{-}mod \implies$
 $(\forall s\ P\ Q\ zs.\ list = (Some\ (Seq\ P\ Q), s)\#zs \longrightarrow$
 $(\exists xs.\ (Some\ P, s)\#xs \in\ cptn\text{-}mod \wedge (zs = (map\ (lift\ Q)\ xs) \vee$
 $(fst(((Some\ P, s)\#xs)!\length\ xs) = None \wedge$
 $(\exists ys.\ (Some\ Q, snd(((Some\ P, s)\#xs)!\length\ xs))\#ys \in\ cptn\text{-}mod$
 $\wedge zs = (map\ (lift\ (Q))\ xs)\@ys))))$

apply (*erule cptn-mod.induct*)
apply *simp-all*
apply *clarify*
apply (*force intro:CptnModOne*)
apply *clarify*
apply (*erule-tac x=Pa in allE*)
apply (*erule-tac x=Q in allE*)
apply *simp*
apply *clarify*
apply (*erule disjE*)
apply (*rule-tac x=(Some Pa,t)#xsa in exI*)
apply (*rule conjI*)
apply *clarify*
apply (*erule CptnModEnv*)
apply (*rule disjI1*)
apply (*simp add:lift-def*)
apply *clarify*
apply (*rule-tac x=(Some Pa,t)#xsa in exI*)


```

apply(rule conjI)
  apply(erule CptnModEnv)
apply(rule disjI2)
apply(rule conjI)
  apply(case-tac xsa,simp,simp)
apply(rule-tac x=ys in exI)
apply(rule conjI)
  apply simp
apply(simp add:lift-def)
apply clarify
apply(erule ctran.cases,simp-all)
apply clarify
apply(rule-tac x=xs in exI)
apply simp
apply clarify
apply(rule-tac x=xs in exI)
apply(simp add: last-length)
done

```

```

lemma cptn-onlyif-cptn-mod-aux [rule-format]:
   $\forall s Q t xs. ((Some a, s), Q, t) \in ctran \longrightarrow (Q, t) \# xs \in cptn-mod$ 
   $\longrightarrow (Some a, s) \# (Q, t) \# xs \in cptn-mod$ 
  supply [[simproc del: defined-all]]
apply(induct a)
apply simp-all
— basic
apply clarify
apply(erule ctran.cases,simp-all)
apply(rule CptnModNone,rule Basic,simp)
apply clarify
apply(erule ctran.cases,simp-all)
— Seq1
apply(rule-tac xs=[[None,ta]] in CptnModSeq2)
  apply(erule CptnModNone)
  apply(rule CptnModOne)
  apply simp
apply simp
apply(simp add:lift-def)
— Seq2
apply(erule-tac x=sa in allE)
apply(erule-tac x=Some P2 in allE)
apply(erule allE,erule impE, assumption)
apply(drule div-seq,simp)
apply clarify
apply(erule disjE)
  apply clarify
apply(erule allE,erule impE, assumption)
apply(erule-tac CptnModSeq1)
apply(simp add:lift-def)

```

```

apply clarify
apply(erule allE,erule impE, assumption)
apply(erule-tac CptnModSeq2)
  apply (simp add:last-length)
  apply (simp add:last-length)
apply(simp add:lift-def)
— Cond
apply clarify
apply(erule ctran.cases,simp-all)
apply(force elim: CptnModCondT)
apply(force elim: CptnModCondF)
— While
apply clarify
apply(erule ctran.cases,simp-all)
apply(rule CptnModNone,erule WhileF,simp)
apply(drule div-seq,force)
apply clarify
apply (erule disjE)
  apply(force elim:CptnModWhile1)
apply clarify
apply(force simp add:last-length elim:CptnModWhile2)
— await
apply clarify
apply(erule ctran.cases,simp-all)
apply(rule CptnModNone,erule Await,simp+)
done

lemma cptn-onlyif-cptn-mod [rule-format]:  $c \in \text{cptn} \implies c \in \text{cptn-mod}$ 
apply(erule cptn.induct)
  apply(rule CptnModOne)
  apply(erule CptnModEnv)
apply(case-tac P)
  apply simp
  apply(erule ctran.cases,simp-all)
apply(force elim:cptn-onlyif-cptn-mod-aux)
done

lemma lift-is-cptn:  $c \in \text{cptn} \implies \text{map } (\text{lift } P) c \in \text{cptn}$ 
apply(erule cptn.induct)
  apply(force simp add:lift-def CptnOne)
  apply(force intro:CptnEnv simp add:lift-def)
apply(force simp add:lift-def intro:CptnComp Seq2 Seq1 elim:ctran.cases)
done

lemma cptn-append-is-cptn [rule-format]:
 $\forall b a. b\#c1 \in \text{cptn} \longrightarrow a\#c2 \in \text{cptn} \longrightarrow (b\#c1)!\text{length } c1 = a \longrightarrow b\#c1@c2 \in \text{cptn}$ 
apply(induct c1)
  apply simp
apply clarify

```

apply(*erule cptn.cases,simp-all*)
apply(*force intro:CptnEnv*)
apply(*force elim:CptnComp*)
done

lemma last-lift: $\llbracket xs \neq [] \rrbracket; \text{fst}(xs!(\text{length } xs - (\text{Suc } 0))) = \text{None}$
 $\implies \text{fst}((\text{map } (\text{lift } P) \text{ xs})!(\text{length } (\text{map } (\text{lift } P) \text{ xs}) - (\text{Suc } 0))) = (\text{Some } P)$
by (*cases (xs ! (length xs - (Suc 0))) (simp add:lift-def)*)

lemma last-fst [*rule-format*]: $P((a \# x)!\text{length } x) \longrightarrow \neg P a \longrightarrow P (x!(\text{length } x - (\text{Suc } 0)))$
by (*induct x simp-all*)

lemma last-fst-esp:
 $\text{fst}(((\text{Some } a,s) \# xs)!(\text{length } xs)) = \text{None} \implies \text{fst}(xs!(\text{length } xs - (\text{Suc } 0))) = \text{None}$
apply(*erule last-fst*)
apply *simp*
done

lemma last-snd: $xs \neq [] \implies$
 $\text{snd}((\text{map } (\text{lift } P) \text{ xs})!(\text{length } (\text{map } (\text{lift } P) \text{ xs}) - (\text{Suc } 0))) = \text{snd}(xs!(\text{length } xs - (\text{Suc } 0)))$
by (*cases (xs ! (length xs - (Suc 0))) (simp-all add:lift-def)*)

lemma Cons-lift: $(\text{Some } (\text{Seq } P \ Q), s) \# (\text{map } (\text{lift } Q) \text{ xs}) = \text{map } (\text{lift } Q) ((\text{Some } P, s) \# xs)$
by (*simp add:lift-def*)

lemma Cons-lift-append:
 $(\text{Some } (\text{Seq } P \ Q), s) \# (\text{map } (\text{lift } Q) \text{ xs}) @ ys = \text{map } (\text{lift } Q) ((\text{Some } P, s) \# xs) @ ys$
by (*simp add:lift-def*)

lemma lift-nth: $i < \text{length } xs \implies \text{map } (\text{lift } Q) \text{ xs } ! i = \text{lift } Q (xs ! i)$
by (*simp add:lift-def*)

lemma snd-lift: $i < \text{length } xs \implies \text{snd}(\text{lift } Q (xs ! i)) = \text{snd } (xs ! i)$
by (*cases xs!i (simp add:lift-def)*)

lemma cptn-if-cptn-mod: $c \in \text{cptn-mod} \implies c \in \text{cptn}$
apply(*erule cptn-mod.induct*)
apply(*rule CptnOne*)
apply(*erule CptnEnv*)
apply(*erule CptnComp,simp*)
apply(*rule CptnComp*)
apply(*erule CondT,simp*)
apply(*rule CptnComp*)
apply(*erule CondF,simp*)
— Seq1

```

apply(erule cptn.cases,simp-all)
  apply(rule CptnOne)
  apply clarify
  apply(drule-tac  $P=P1$  in lift-is-cptn)
  apply(simp add:lift-def)
  apply(rule CptnEnv,simp)
apply clarify
apply(simp add:lift-def)
apply(rule conjI)
  apply clarify
  apply(rule CptnComp)
  apply(rule Seq1,simp)
  apply(drule-tac  $P=P1$  in lift-is-cptn)
  apply(simp add:lift-def)
apply clarify
apply(rule CptnComp)
  apply(rule Seq2,simp)
apply(drule-tac  $P=P1$  in lift-is-cptn)
apply(simp add:lift-def)
— Seq2
apply(rule cptn-append-is-cptn)
  apply(drule-tac  $P=P1$  in lift-is-cptn)
  apply(simp add:lift-def)
  apply simp
apply(simp split: if-split-asm)
apply(frule-tac  $P=P1$  in last-lift)
  apply(rule last-fst-esp)
  apply (simp add:last-length)
apply(simp add:Cons-lift lift-def split-def last-conv-nth)
— While1
apply(rule CptnComp)
  apply(rule WhileT,simp)
apply(drule-tac  $P=While\ b\ P$  in lift-is-cptn)
apply(simp add:lift-def)
— While2
apply(rule CptnComp)
  apply(rule WhileT,simp)
apply(rule cptn-append-is-cptn)
  apply(drule-tac  $P=While\ b\ P$  in lift-is-cptn)
  apply(simp add:lift-def)
  apply simp
apply(simp split: if-split-asm)
apply(frule-tac  $P=While\ b\ P$  in last-lift)
  apply(rule last-fst-esp,simp add:last-length)
apply(simp add:Cons-lift lift-def split-def last-conv-nth)
done

theorem cptn-iff-cptn-mod:  $(c \in cptn) = (c \in cptn-mod)$ 
apply(rule iffI)

```

apply(*erule cptn-onlyif-cptn-mod*)
apply(*erule cptn-if-cptn-mod*)
done

3.3 Validity of Correctness Formulas

3.3.1 Validity for Component Programs.

type-synonym *'a rgformula* =
'a com × *'a set* × (*'a* × *'a*) *set* × (*'a* × *'a*) *set* × *'a set*

definition *assum* :: (*'a set* × (*'a* × *'a*) *set*) ⇒ (*'a confs*) *set* **where**
assum ≡ λ(*pre*, *rely*). {*c*. *snd*(*c!*0) ∈ *pre* ∧ (∀ *i*. *Suc* *i* < *length* *c* →
*c!**i* - *e* → *c!*(*Suc* *i*) → (*snd*(*c!**i*), *snd*(*c!**Suc* *i*)) ∈ *rely*)}

definition *comm* :: ((*'a* × *'a*) *set* × *'a set*) ⇒ (*'a confs*) *set* **where**
comm ≡ λ(*guar*, *post*). {*c*. (∀ *i*. *Suc* *i* < *length* *c* →
*c!**i* - *c* → *c!*(*Suc* *i*) → (*snd*(*c!**i*), *snd*(*c!**Suc* *i*)) ∈ *guar*) ∧
(*fst* (*last* *c*) = *None* → *snd* (*last* *c*) ∈ *post*)}

definition *com-validity* :: *'a com* ⇒ *'a set* ⇒ (*'a* × *'a*) *set* ⇒ (*'a* × *'a*) *set* ⇒ *'a set* ⇒ *bool*

(⊨ - *sat* [-, -, -, -] › [60,0,0,0,0] 45) **where**
⊨ *P sat* [*pre*, *rely*, *guar*, *post*] ≡
∀ *s*. *cp* (*Some* *P*) *s* ∩ *assum*(*pre*, *rely*) ⊆ *comm*(*guar*, *post*)

3.3.2 Validity for Parallel Programs.

definition *All-None* :: ((*'a com*) *option*) *list* ⇒ *bool* **where**
All-None *xs* ≡ ∀ *c* ∈ *set* *xs*. *c* = *None*

definition *par-assum* :: (*'a set* × (*'a* × *'a*) *set*) ⇒ (*'a par-confs*) *set* **where**
par-assum ≡ λ(*pre*, *rely*). {*c*. *snd*(*c!*0) ∈ *pre* ∧ (∀ *i*. *Suc* *i* < *length* *c* →
*c!**i* - *pe* → *c!**Suc* *i* → (*snd*(*c!**i*), *snd*(*c!**Suc* *i*)) ∈ *rely*)}

definition *par-comm* :: ((*'a* × *'a*) *set* × *'a set*) ⇒ (*'a par-confs*) *set* **where**
par-comm ≡ λ(*guar*, *post*). {*c*. (∀ *i*. *Suc* *i* < *length* *c* →
*c!**i* - *pc* → *c!**Suc* *i* → (*snd*(*c!**i*), *snd*(*c!**Suc* *i*)) ∈ *guar*) ∧
(*All-None* (*fst* (*last* *c*)) → *snd* (*last* *c*) ∈ *post*)}

definition *par-com-validity* :: *'a par-com* ⇒ *'a set* ⇒ (*'a* × *'a*) *set* ⇒ (*'a* × *'a*) *set* ⇒ *bool*

(⊨ - *SAT* [-, -, -, -] › [60,0,0,0,0] 45) **where**
⊨ *Ps SAT* [*pre*, *rely*, *guar*, *post*] ≡
∀ *s*. *par-cp* *Ps* *s* ∩ *par-assum*(*pre*, *rely*) ⊆ *par-comm*(*guar*, *post*)

3.3.3 Compositionality of the Semantics

Definition of the conjoin operator

definition *same-length* :: 'a par-confs \Rightarrow ('a confs) list \Rightarrow bool **where**
same-length c clist \equiv ($\forall i < \text{length clist}. \text{length}(clist!i) = \text{length } c$)

definition *same-state* :: 'a par-confs \Rightarrow ('a confs) list \Rightarrow bool **where**
same-state c clist \equiv ($\forall i < \text{length clist}. \forall j < \text{length } c. \text{snd}(c!j) = \text{snd}((clist!i)!j)$)

definition *same-program* :: 'a par-confs \Rightarrow ('a confs) list \Rightarrow bool **where**
same-program c clist \equiv ($\forall j < \text{length } c. \text{fst}(c!j) = \text{map } (\lambda x. \text{fst}(\text{nth } x \ j)) \text{ clist}$)

definition *compat-label* :: 'a par-confs \Rightarrow ('a confs) list \Rightarrow bool **where**
compat-label c clist \equiv ($\forall j. \text{Suc } j < \text{length } c \longrightarrow$
 $(c!j - pc \rightarrow c!\text{Suc } j \wedge (\exists i < \text{length clist}. (clist!i)!j - c \rightarrow (clist!i)! \text{Suc } j \wedge$
 $(\forall l < \text{length clist}. l \neq i \longrightarrow (clist!l)!j - e \rightarrow (clist!l)! \text{Suc } j))) \vee$
 $(c!j - pe \rightarrow c!\text{Suc } j \wedge (\forall i < \text{length clist}. (clist!i)!j - e \rightarrow (clist!i)! \text{Suc } j))$)

definition *conjoin* :: 'a par-confs \Rightarrow ('a confs) list \Rightarrow bool ($\leftarrow \propto \rightarrow$ [65,65] 64)
where
c \propto *clist* \equiv (*same-length* c *clist*) \wedge (*same-state* c *clist*) \wedge (*same-program* c *clist*)
 \wedge (*compat-label* c *clist*)

Some previous lemmas

lemma *list-eq-if* [rule-format]:
 $\forall ys. xs = ys \longrightarrow (\text{length } xs = \text{length } ys) \longrightarrow (\forall i < \text{length } xs. xs!i = ys!i)$
by (*induct* xs) *auto*

lemma *list-eq*: $(\text{length } xs = \text{length } ys \wedge (\forall i < \text{length } xs. xs!i = ys!i)) = (xs = ys)$
apply (*rule iffI*)
apply *clarify*
apply (*erule nth-equalityI*)
apply *simp+*
done

lemma *nth-tl*: $\llbracket ys!0 = a; ys \neq [] \rrbracket \Longrightarrow ys = (a \# (\text{tl } ys))$
by (*cases* ys) *simp-all*

lemma *nth-tl-if* [rule-format]: $ys \neq [] \longrightarrow ys!0 = a \longrightarrow P \text{ } ys \longrightarrow P (a \# (\text{tl } ys))$
by (*induct* ys) *simp-all*

lemma *nth-tl-onlyif* [rule-format]: $ys \neq [] \longrightarrow ys!0 = a \longrightarrow P (a \# (\text{tl } ys)) \longrightarrow P \text{ } ys$
by (*induct* ys) *simp-all*

lemma *seq-not-eq1*: $\text{Seq } c1 \ c2 \neq c1$
by (*induct* c1) *auto*

lemma *seq-not-eq2*: $\text{Seq } c1 \ c2 \neq c2$

by (*induct c2*) *auto*

lemma *if-not-eq1*: *Cond b c1 c2 ≠ c1*
by (*induct c1*) *auto*

lemma *if-not-eq2*: *Cond b c1 c2 ≠ c2*
by (*induct c2*) *auto*

lemmas *seq-and-if-not-eq* [*simp*] = *seq-not-eq1 seq-not-eq2*
seq-not-eq1 [*THEN not-sym*] *seq-not-eq2* [*THEN not-sym*]
if-not-eq1 if-not-eq2 if-not-eq1 [*THEN not-sym*] *if-not-eq2* [*THEN not-sym*]

lemma *prog-not-eq-in-ctran-aux*:
assumes *c*: $(P,s) -c \rightarrow (Q,t)$
shows $P \neq Q$ **using** *c*
by (*induct x1 ≡ (P,s) x2 ≡ (Q,t) arbitrary: P s Q t*) *auto*

lemma *prog-not-eq-in-ctran* [*simp*]: $\neg (P,s) -c \rightarrow (P,t)$
apply *clarify*
apply (*drule prog-not-eq-in-ctran-aux*)
apply *simp*
done

lemma *prog-not-eq-in-par-ctran-aux* [*rule-format*]: $(P,s) -pc \rightarrow (Q,t) \implies (P \neq Q)$
apply (*erule par-ctran.induct*)
apply (*drule prog-not-eq-in-ctran-aux*)
apply *clarify*
apply (*drule list-eq-if*)
apply *simp-all*
apply *force*
done

lemma *prog-not-eq-in-par-ctran* [*simp*]: $\neg (P,s) -pc \rightarrow (P,t)$
apply *clarify*
apply (*drule prog-not-eq-in-par-ctran-aux*)
apply *simp*
done

lemma *tl-in-cptn*: $\llbracket a \# xs \in \text{cptn}; xs \neq [] \rrbracket \implies xs \in \text{cptn}$
by (*force elim: cptn.cases*)

lemma *tl-zero* [*rule-format*]:
 $P (ys! \text{Suc } j) \longrightarrow \text{Suc } j < \text{length } ys \longrightarrow ys \neq [] \longrightarrow P (\text{tl}(ys)!j)$
by (*induct ys*) *simp-all*

3.3.4 The Semantics is Compositional

lemma *aux-if* [*rule-format*]:
 $\forall xs \ s \ \text{clist}. (\text{length } \text{clist} = \text{length } xs \wedge (\forall i < \text{length } xs. (xs!i,s) \# \text{clist}!i \in \text{cptn}))$

```

     $\wedge ((xs, s)\#ys \propto \text{map } (\lambda i. (\text{fst } i, s)\#\text{snd } i) (\text{zip } xs \text{ clist}))$ 
     $\longrightarrow (xs, s)\#ys \in \text{par-cptn}$ 
apply(induct ys)
apply(clarify)
apply(rule ParCptnOne)
apply(clarify)
apply(simp add: conjoin-def compat-label-def)
apply clarify
apply(erule-tac x=0 and P= $\lambda j. H j \longrightarrow (P j \vee Q j)$  for H P Q in all-dupE, simp)
apply(erule disjE)
— first step is a Component step
apply clarify
apply simp
apply(subgoal-tac a=(xs[i:=fst(clist!i!0)]))
apply(subgoal-tac b=snd(clist!i!0), simp)
prefer 2
apply(simp add: same-state-def)
apply(erule-tac x=i in allE, erule impE, assumption,
    erule-tac x=1 and P= $\lambda j. (H j) \longrightarrow (\text{snd } (d j))=(\text{snd } (e j))$  for H d e in
allE, simp)
prefer 2
apply(simp add: same-program-def)
apply(erule-tac x=1 and P= $\lambda j. H j \longrightarrow (\text{fst } (s j))=(t j)$  for H s t in allE, simp)
apply(rule nth-equalityI, simp)
apply clarify
apply(case-tac i=ia, simp, simp)
apply(erule-tac x=ia and P= $\lambda j. H j \longrightarrow I j \longrightarrow J j$  for H I J in allE)
apply(drule-tac t=i in not-sym, simp)
apply(erule etranE, simp)
apply(rule ParCptnComp)
apply(erule ParComp, simp)
— applying the induction hypothesis
apply(erule-tac x=xs[i := fst (clist ! i ! 0)] in allE)
apply(erule-tac x=snd (clist ! i ! 0) in allE)
apply(erule mp)
apply(rule-tac x=map tl clist in exI, simp)
apply(rule conjI, clarify)
apply(case-tac i=ia, simp)
apply(rule nth-tl-if)
apply(force simp add: same-length-def length-Suc-conv)
apply simp
apply(erule allE, erule impE, assumption, erule tl-in-cptn)
apply(force simp add: same-length-def length-Suc-conv)
apply(rule nth-tl-if)
apply(force simp add: same-length-def length-Suc-conv)
apply(simp add: same-state-def)
apply(erule-tac x=ia in allE, erule impE, assumption,
    erule-tac x=1 and P= $\lambda j. H j \longrightarrow (\text{snd } (d j))=(\text{snd } (e j))$  for H d e in allE)

```



```

apply(erule-tac x=ia and P= $\lambda j. H j \longrightarrow I j \longrightarrow J j$  for H I J in allE)
apply(drule-tac t=i in not-sym,simp)
apply(erule etranE,simp)
apply(erule allE,erule impE,assumption,erule tl-in-cptn)
apply(force simp add:same-length-def length-Suc-conv)
apply(simp add:same-length-def same-state-def)
apply(rule conjI)
apply clarify
apply(case-tac j,simp,simp)
apply(erule-tac x=ia in allE, erule impE, assumption,
      erule-tac x=Suc(Suc nat) and P= $\lambda j. H j \longrightarrow (snd (d j))=(snd (e j))$  for
H d e in allE,simp)
apply(force simp add:same-length-def length-Suc-conv)
apply(rule conjI)
apply(simp add:same-program-def)
apply clarify
apply(case-tac j,simp)
apply(rule nth-equalityI,simp)
apply clarify
apply(case-tac i=ia,simp,simp)
apply(erule-tac x=Suc(Suc nat) and P= $\lambda j. H j \longrightarrow (fst (s j))=(t j)$  for H s t
in allE,simp)
apply(rule nth-equalityI,simp,simp)
apply(force simp add:length-Suc-conv)
apply(rule allI,rule impI)
apply(erule-tac x=Suc j and P= $\lambda j. H j \longrightarrow (I j \vee J j)$  for H I J in allE,simp)
apply(erule disjE)
apply clarify
apply(rule-tac x=ia in exI,simp)
apply(case-tac i=ia,simp)
apply(rule conjI)
apply(force simp add: length-Suc-conv)
apply clarify
apply(erule-tac x=l and P= $\lambda j. H j \longrightarrow I j \longrightarrow J j$  for H I J in allE,erule
impE,assumption)
apply(erule-tac x=l and P= $\lambda j. H j \longrightarrow I j \longrightarrow J j$  for H I J in allE,erule
impE,assumption)
apply simp
apply(case-tac j,simp)
apply(rule tl-zero)
apply(erule-tac x=l in allE, erule impE, assumption,
      erule-tac x=1 and P= $\lambda j. (H j) \longrightarrow (snd (d j))=(snd (e j))$  for H d e
in allE,simp)
apply(force elim:etranE intro:Env)
apply force
apply force
apply simp
apply(rule tl-zero)
apply(erule tl-zero)

```

apply force
apply force
apply force
apply force
apply(rule conjI,simp)
apply(rule nth-tl-if)
apply force
apply(erule-tac x=ia in allE, erule impE, assumption,
erule-tac x=1 and P= λ j. H j \longrightarrow (snd (d j))=(snd (e j)) for H d e in
allE)
apply(erule-tac x=ia and P= λ j. H j \longrightarrow I j \longrightarrow J j for H I J in allE)
apply(drule-tac t=i in not-sym,simp)
apply(erule etranE,simp)
apply(erule tl-zero)
apply force
apply force
apply clarify
apply(case-tac i=l,simp)
apply(rule nth-tl-if)
apply(erule-tac x=l and P= λ j. H j \longrightarrow (length (s j) = t) for H s t in
allE,force)
apply simp
apply(erule-tac P= λ j. H j \longrightarrow I j \longrightarrow J j for H I J in allE,erule impE,assumption,erule
impE,assumption)
apply(erule tl-zero,force)
apply(erule-tac x=l and P= λ j. H j \longrightarrow (length (s j) = t) for H s t in allE,force)
apply(rule nth-tl-if)
apply(erule-tac x=l and P= λ j. H j \longrightarrow (length (s j) = t) for H s t in
allE,force)
apply(erule-tac x=l in allE, erule impE, assumption,
erule-tac x=1 and P= λ j. H j \longrightarrow (snd (d j))=(snd (e j)) for H d e in
allE)
apply(erule-tac x=l and P= λ j. H j \longrightarrow I j \longrightarrow J j for H I J in allE,erule
impE, assumption,simp)
apply(erule etranE,simp)
apply(rule tl-zero)
apply force
apply force
apply(erule-tac x=l and P= λ j. H j \longrightarrow (length (s j) = t) for H s t in allE,force)
apply(rule disjI2)
apply(case-tac j,simp)
apply clarify
apply(rule tl-zero)
apply(erule-tac x=ia and P= λ j. H j \longrightarrow I j \in etran for H I in allE,erule impE,
assumption)
apply(case-tac i=ia,simp,simp)
apply(erule-tac x=ia in allE, erule impE, assumption,
erule-tac x=1 and P= λ j. H j \longrightarrow (snd (d j))=(snd (e j)) for H d e in allE)
apply(erule-tac x=ia and P= λ j. H j \longrightarrow I j \longrightarrow J j for H I J in allE,erule

```

impE, assumption, simp)
  apply(force elim: etranE intro: Env)
  apply force
  apply(erule-tac x=ia and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in
allE, force)
  apply simp
  apply clarify
  apply(rule tl-zero)
  apply(rule tl-zero, force)
  apply force
  apply(erule-tac x=ia and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in
allE, force)
  apply force
  apply(erule-tac x=ia and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in allE, force)
— first step is an environmental step
  apply clarify
  apply(erule par-etran.cases)
  apply simp
  apply(rule ParCptnEnv)
  apply(erule-tac x=Ps in allE)
  apply(erule-tac x=t in allE)
  apply(erule mp)
  apply(rule-tac x=map tl clist in exI, simp)
  apply(rule conjI)
  apply clarify
  apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow I j \in cptn$  for H I in allE, simp)
  apply(erule cptn.cases)
  apply(simp add:same-length-def)
  apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in allE, force)
  apply(simp add:same-state-def)
  apply(erule-tac x=i in allE, erule impE, assumption,
erule-tac x=1 and P= $\lambda j. H j \longrightarrow (snd (d j))=(snd (e j))$  for H d e in allE, simp)
  apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow J j \in etran$  for H J in allE, simp)
  apply(erule etranE, simp)
  apply(simp add:same-state-def same-length-def)
  apply(rule conjI, clarify)
  apply(case-tac j, simp, simp)
  apply(erule-tac x=i in allE, erule impE, assumption,
erule-tac x=Suc(Suc nat) and P= $\lambda j. H j \longrightarrow (snd (d j))=(snd (e j))$  for H
d e in allE, simp)
  apply(rule tl-zero)
  apply(simp)
  apply force
  apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in allE, force)
  apply(rule conjI)
  apply(simp add:same-program-def)
  apply clarify
  apply(case-tac j, simp)
  apply(rule nth-equalityI, simp)

```

```

apply clarify
apply simp
apply(erule-tac x=Suc(Suc nat) and P= $\lambda j. H j \longrightarrow (fst (s j))=(t j)$  for H s t
in allE,simp)
apply(rule nth-equalityI,simp,simp)
apply(force simp add:length-Suc-conv)
apply(rule allI,rule impI)
apply(erule-tac x=Suc j and P= $\lambda j. H j \longrightarrow (I j \vee J j)$  for H I J in allE,simp)
apply(erule disjE)
apply clarify
apply(rule-tac x=i in exI,simp)
apply(rule conjI)
apply(erule-tac x=i and P= $\lambda i. H i \longrightarrow J i \in etran$  for H J in allE, erule impE,
assumption)
apply(erule etranE,simp)
apply(erule-tac x=i in allE, erule impE, assumption,
erule-tac x=1 and P= $\lambda j. H j \longrightarrow (snd (d j))=(snd (e j))$  for H d e in
allE,simp)
apply(rule nth-tl-if)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in allE,force)
apply simp
apply(erule tl-zero,force)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in allE,force)
apply clarify
apply(erule-tac x=l and P= $\lambda i. H i \longrightarrow J i \in etran$  for H J in allE, erule impE,
assumption)
apply(erule etranE,simp)
apply(erule-tac x=l in allE, erule impE, assumption,
erule-tac x=1 and P= $\lambda j. H j \longrightarrow (snd (d j))=(snd (e j))$  for H d e in
allE,simp)
apply(rule nth-tl-if)
apply(erule-tac x=l and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in allE,force)
apply simp
apply(rule tl-zero,force)
apply force
apply(erule-tac x=l and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in allE,force)
apply(rule disjI2)
apply simp
apply clarify
apply(case-tac j,simp)
apply(rule tl-zero)
apply(erule-tac x=i and P= $\lambda i. H i \longrightarrow J i \in etran$  for H J in allE, erule
impE, assumption)
apply(erule-tac x=i and P= $\lambda i. H i \longrightarrow J i \in etran$  for H J in allE, erule
impE, assumption)
apply(force elim:etranE intro:Env)
apply force
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (length (s j) = t)$  for H s t in allE,force)
apply simp

```

```

apply(rule tl-zero)
  apply(rule tl-zero,force)
  apply force
  apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE,force)
  apply force
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE,force)
done

```

```

lemma aux-onlyif [rule-format]:  $\forall xs s. (xs, s)\#ys \in \text{par-cptn} \longrightarrow$ 
  ( $\exists \text{clist. } (\text{length } \text{clist} = \text{length } xs) \wedge$ 
   $(xs, s)\#ys \propto \text{map } (\lambda i. (\text{fst } i, s)\#(\text{snd } i)) (\text{zip } xs \text{ clist}) \wedge$ 
   $(\forall i < \text{length } xs. (xs!i, s)\#(\text{clist}!i) \in \text{cptn})$ )
  supply [[simproc del: defined-all]]
apply(induct ys)
apply(clarify)
apply(rule-tac x=map ( $\lambda i. []$ ) [0.. $\text{length } xs$ ] in exI)
apply(simp add: conjoin-def same-length-def same-state-def same-program-def compat-label-def)
apply(rule conjI)
  apply(rule nth-equalityI, simp, simp)
apply(force intro: cptn.intros)
apply(clarify)
apply(erule par-cptn.cases, simp)
  apply simp
  apply(erule-tac x=xs in allE)
  apply(erule-tac x=t in allE, simp)
  apply clarify
  apply(rule-tac x=(map ( $\lambda j. (P!j, t)\#(\text{clist}!j)$ ) [0.. $\text{length } P$ ]) in exI, simp)
  apply(rule conjI)
  prefer 2
  apply clarify
  apply(rule CptnEnv, simp)
apply(simp add: conjoin-def same-length-def same-state-def)
apply (rule conjI)
  apply clarify
  apply(case-tac j, simp, simp)
apply(rule conjI)
  apply(simp add: same-program-def)
  apply clarify
  apply(case-tac j, simp)
  apply(rule nth-equalityI, simp, simp)
apply simp
  apply(rule nth-equalityI, simp, simp)
apply(simp add: compat-label-def)
apply clarify
apply(case-tac j, simp)
apply(simp add: ParEnv)
apply clarify
apply(simp add: Env)

```

```

apply simp
apply(erule-tac  $x=\text{nat}$  in allE,erule impE, assumption)
apply(erule disjE,simp)
  apply clarify
    apply(rule-tac  $x=i$  in exI,simp)
  apply force
apply(erule par-ctran.cases,simp)
apply(erule-tac  $x=Ps[i:=r]$  in allE)
apply(erule-tac  $x=ta$  in allE,simp)
apply clarify
apply(rule-tac  $x=(\text{map } (\lambda j. (Ps!j, ta)\#(clist!j)) [0..\text{length } Ps]) [i:=((r, ta)\#(clist!i))]$ 
in exI,simp)
apply(rule conjI)
  prefer 2
  apply clarify
    apply(case-tac  $i=ia$ ,simp)
    apply(erule CptnComp)
    apply(erule-tac  $x=ia$  and  $P=\lambda j. H j \longrightarrow (I j \in \text{cptn})$  for  $H I$  in allE,simp)
  apply simp
apply(erule-tac  $x=ia$  in allE)
apply(rule CptnEnv,simp)
apply(simp add:conjoin-def)
apply (rule conjI)
  apply(simp add:same-length-def)
  apply clarify
    apply(case-tac  $i=ia$ ,simp,simp)
apply(rule conjI)
  apply(simp add:same-state-def)
  apply clarify
    apply(case-tac  $j, simp, simp$  (no-asm-simp))
    apply(case-tac  $i=ia$ ,simp,simp)
apply(rule conjI)
  apply(simp add:same-program-def)
  apply clarify
    apply(case-tac  $j$ ,simp)
    apply(rule nth-equalityI,simp,simp)
  apply simp
apply(rule nth-equalityI,simp,simp)
apply(erule-tac  $x=\text{nat}$  and  $P=\lambda j. H j \longrightarrow (\text{fst } (a j))=((b j))$  for  $H a b$  in allE)
apply(case-tac nat)
  apply clarify
    apply(case-tac  $i=ia$ ,simp,simp)
  apply clarify
    apply(case-tac  $i=ia$ ,simp,simp)
apply(simp add:compat-label-def)
apply clarify
apply(case-tac  $j$ )
  apply(rule conjI,simp)
  apply(erule ParComp,assumption)

```

```

apply clarify
apply(rule-tac x=i in exI,simp)
apply clarify
apply(rule Env)
apply simp
apply(erule-tac x=nat and P=λj. H j → (P j ∨ Q j) for H P Q in allE,simp)
apply(erule disjE)
apply clarify
apply(rule-tac x=ia in exI,simp)
apply(rule conjI)
apply(case-tac i=ia,simp,simp)
apply clarify
apply(case-tac i=l,simp)
apply(case-tac l=ia,simp,simp)
apply(erule-tac x=l in allE,erule impE,assumption,erule impE, assumption,simp)
apply simp
apply(erule-tac x=l in allE,erule impE,assumption,erule impE, assumption,simp)
apply clarify
apply(erule-tac x=ia and P=λj. H j → (P j)∈etran for H P in allE, erule
impE, assumption)
apply(case-tac i=ia,simp,simp)
done

```

```

lemma one-iff-aux: xs≠[] ⇒ (∀ ys. ((xs, s)#ys ∈ par-cptn) =
(∃ clist. length clist= length xs ∧
((xs, s)#ys ∝ map (λi. (fst i,s)#(snd i)) (zip xs clist)) ∧
(∀ i<length xs. (xs!i,s)#(clist!i) ∈ cptn))) =
(par-cp (xs) s = {c. ∃ clist. (length clist)=(length xs) ∧
(∀ i<length clist. (clist!i) ∈ cp(xs!i) s) ∧ c ∝ clist})
apply (rule iffI)
apply(rule subset-antisym)
apply(rule subsetI)
apply(clarify)
apply(simp add:par-cp-def cp-def)
apply(case-tac x)
apply(force elim:par-cptn.cases)
apply simp
apply(rename-tac a list)
apply(erule-tac x=list in allE)
apply clarify
apply simp
apply(rule-tac x=map (λi. (fst i, s) # snd i) (zip xs clist) in exI,simp)
apply(rule subsetI)
apply(clarify)
apply(case-tac x)
apply(erule-tac x=0 in allE)
apply(simp add:cp-def conjoin-def same-length-def same-program-def same-state-def
compat-label-def)
apply clarify

```

```

apply(erule cptn.cases.force,force,force)
apply(simp add:par-cp-def conjoin-def same-length-def same-program-def same-state-def
compat-label-def)
apply clarify
apply(erule-tac x=0 and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in all-dupE)
apply(subgoal-tac a = xs)
apply(subgoal-tac b = s,simp)
prefer 3
apply(erule-tac x=0 and P= $\lambda j. H j \longrightarrow (\text{fst } (s j))=((t j))$  for H s t in allE)
apply (simp add:cp-def)
apply(rule nth-equalityI,simp,simp)
prefer 2
apply(erule-tac x=0 in allE)
apply (simp add:cp-def)
apply(erule-tac x=0 and P= $\lambda j. H j \longrightarrow (\forall i. T i \longrightarrow (\text{snd } (d j i))=(\text{snd } (e j i)))$  for H T d e in allE,simp)
apply(erule-tac x=0 and P= $\lambda j. H j \longrightarrow (\text{snd } (d j))=(\text{snd } (e j))$  for H d e in
allE,simp)
apply(erule-tac x=list in allE)
apply(rule-tac x=map tl clist in exI,simp)
apply(rule conjI)
apply clarify
apply(case-tac j,simp)
apply(erule-tac x=i in allE, erule impE, assumption,
erule-tac x=0 and P= $\lambda j. H j \longrightarrow (\text{snd } (d j))=(\text{snd } (e j))$  for H d e in
allE,simp)
apply(erule-tac x=i in allE, erule impE, assumption,
erule-tac x=Suc nat and P= $\lambda j. H j \longrightarrow (\text{snd } (d j))=(\text{snd } (e j))$  for H d e
in allE)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply(rule conjI)
apply clarify
apply(rule nth-equalityI,simp,simp)
apply(case-tac j)
apply clarify
apply(erule-tac x=i in allE)
apply(simp add:cp-def)
apply clarify
apply simp
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply(thin-tac H = ( $\exists i. J i$ ) for H J)
apply(rule conjI)
apply clarify
apply(erule-tac x=j in allE,erule impE, assumption,erule disjE)
apply clarify
apply(rule-tac x=i in exI,simp)
apply(case-tac j,simp)

```



```

apply(rule conjI)
apply(erule-tac x=i in allE)
apply(simp add:cp-def)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply clarify
apply(erule-tac x=l in allE)
apply(erule-tac x=l and P= $\lambda j. H j \longrightarrow I j \longrightarrow J j$  for H I J in allE)
apply clarify
apply(simp add:cp-def)
apply(erule-tac x=l and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE)
apply(case-tac clist!l,simp,simp)
apply simp
apply(rule conjI)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply clarify
apply(erule-tac x=l and P= $\lambda j. H j \longrightarrow I j \longrightarrow J j$  for H I J in allE)
apply(erule-tac x=l and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE)
apply(case-tac clist!l,simp,simp)
apply clarify
apply(erule-tac x=i in allE)
apply(simp add:cp-def)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE)
apply(case-tac clist!i,simp)
apply(rule nth-tl-if,simp,simp)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (P j) \in \text{etran}$  for H P in allE, erule
impE, assumption,simp)
apply(simp add:cp-def)
apply clarify
apply(rule nth-tl-if)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{length } (s j) = t)$  for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply force
apply force
apply clarify
apply(rule iffI)
apply(simp add:par-cp-def)
apply(erule-tac c=(xs, s) # ys in equalityCE)
apply simp
apply clarify
apply(rule-tac x=map tl clist in exI)
apply simp
apply (rule conjI)
apply(simp add:conjoin-def cp-def)
apply(rule conjI)
apply clarify
apply(unfold same-length-def)
apply clarify

```

```

apply(erule-tac x=i and P= $\lambda$ j. H j  $\longrightarrow$  (length (s j) = t) for H s t in
allE,simp)
apply(rule conjI)
apply(simp add:same-state-def)
apply clarify
apply(erule-tac x=i in allE, erule impE, assumption,
erule-tac x=j and P= $\lambda$ j. H j  $\longrightarrow$  (snd (d j))=(snd (e j)) for H d e in allE)
apply(case-tac j,simp)
apply(erule-tac x=i and P= $\lambda$ j. H j  $\longrightarrow$  (length (s j) = t) for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply(rule conjI)
apply(simp add:same-program-def)
apply clarify
apply(rule nth-equalityI,simp,simp)
apply(case-tac j,simp)
apply clarify
apply(erule-tac x=i and P= $\lambda$ j. H j  $\longrightarrow$  (length (s j) = t) for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply clarify
apply(simp add:compat-label-def)
apply(rule allI,rule impI)
apply(erule-tac x=j in allE,erule impE, assumption)
apply(erule disjE)
apply clarify
apply(rule-tac x=i in exI,simp)
apply(rule conjI)
apply(erule-tac x=i in allE)
apply(case-tac j,simp)
apply(erule-tac x=i and P= $\lambda$ j. H j  $\longrightarrow$  (length (s j) = t) for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply(erule-tac x=i and P= $\lambda$ j. H j  $\longrightarrow$  (length (s j) = t) for H s t in allE)
apply(case-tac clist!i,simp,simp)
apply clarify
apply(erule-tac x=l and P= $\lambda$ j. H j  $\longrightarrow$  I j  $\longrightarrow$  J j for H I J in allE)
apply(erule-tac x=l and P= $\lambda$ j. H j  $\longrightarrow$  (length (s j) = t) for H s t in allE)
apply(case-tac clist!l,simp,simp)
apply(erule-tac x=l in allE,simp)
apply(rule disjI2)
apply clarify
apply(rule tl-zero)
apply(case-tac j,simp,simp)
apply(rule tl-zero,force)
apply force
apply(erule-tac x=i and P= $\lambda$ j. H j  $\longrightarrow$  (length (s j) = t) for H s t in
allE,force)
apply force
apply(erule-tac x=i and P= $\lambda$ j. H j  $\longrightarrow$  (length (s j) = t) for H s t in allE,force)
apply clarify
apply(erule-tac x=i in allE)

```

```

apply(simp add:cp-def)
apply(rule nth-tl-if)
  apply(simp add:conjoin-def)
  apply clarify
  apply(simp add:same-length-def)
  apply(erule-tac x=i in allE,simp)
  apply simp
  apply simp
  apply simp
apply clarify
apply(erule-tac c=(xs, s) # ys in equalityCE)
  apply(simp add:par-cp-def)
apply simp
apply(erule-tac x=map (λi. (fst i, s) # snd i) (zip xs clist) in allE)
apply simp
apply clarify
apply(simp add:cp-def)
done

```

```

theorem one: xs≠[] ⇒
  par-cp xs s = {c. ∃ clist. (length clist)=(length xs) ∧
    (∀ i<length clist. (clist!i) ∈ cp(xs!i) s) ∧ c ∝ clist}

```

```

apply(frule one-iff-aux)
apply(drule sym)
apply(erule iffD2)
apply clarify
apply(rule iffI)
  apply(erule aux-onlyif)
apply clarify
apply(force intro:aux-if)
done

```

end

3.4 The Proof System

theory *RG-Hoare* **imports** *RG-Tran* **begin**

3.4.1 Proof System for Component Programs

declare *Un-subset-iff* [*simp del*] *sup.bounded-iff* [*simp del*]

definition *stable* :: 'a set ⇒ ('a × 'a) set ⇒ bool **where**
stable ≡ λf g. (∀ x y. x ∈ f → (x, y) ∈ g → y ∈ f)

inductive

rghoare :: ['a com, 'a set, ('a × 'a) set, ('a × 'a) set, 'a set] ⇒ bool
 (⟦ - sat [-, -, -, -] [60,0,0,0,0] 45)

where

Basic: $\llbracket pre \subseteq \{s. f s \in post\}; \{(s,t). s \in pre \wedge (t=f s \vee t=s)\} \subseteq guar;$
 $stable\ pre\ rely; stable\ post\ rely \rrbracket$
 $\implies \vdash Basic\ f\ sat\ [pre, rely, guar, post]$

| *Seq*: $\llbracket \vdash P\ sat\ [pre, rely, guar, mid]; \vdash Q\ sat\ [mid, rely, guar, post] \rrbracket$
 $\implies \vdash Seq\ P\ Q\ sat\ [pre, rely, guar, post]$

| *Cond*: $\llbracket stable\ pre\ rely; \vdash P1\ sat\ [pre \cap b, rely, guar, post];$
 $\vdash P2\ sat\ [pre \cap \neg b, rely, guar, post]; \forall s. (s,s) \in guar \rrbracket$
 $\implies \vdash Cond\ b\ P1\ P2\ sat\ [pre, rely, guar, post]$

| *While*: $\llbracket stable\ pre\ rely; (pre \cap \neg b) \subseteq post; stable\ post\ rely;$
 $\vdash P\ sat\ [pre \cap b, rely, guar, pre]; \forall s. (s,s) \in guar \rrbracket$
 $\implies \vdash While\ b\ P\ sat\ [pre, rely, guar, post]$

| *Await*: $\llbracket stable\ pre\ rely; stable\ post\ rely;$
 $\forall V. \vdash P\ sat\ [pre \cap b \cap \{V\}, \{(s, t). s = t\},$
 $UNIV, \{s. (V, s) \in guar\} \cap post] \rrbracket$
 $\implies \vdash Await\ b\ P\ sat\ [pre, rely, guar, post]$

| *Conseq*: $\llbracket pre \subseteq pre'; rely \subseteq rely'; guar' \subseteq guar; post' \subseteq post;$
 $\vdash P\ sat\ [pre', rely', guar', post'] \rrbracket$
 $\implies \vdash P\ sat\ [pre, rely, guar, post]$

definition *Pre* :: 'a rgformula \Rightarrow 'a set **where**
Pre $x \equiv fst(snd\ x)$

definition *Post* :: 'a rgformula \Rightarrow 'a set **where**
Post $x \equiv snd(snd(snd\ x))$

definition *Rely* :: 'a rgformula \Rightarrow ('a \times 'a) set **where**
Rely $x \equiv fst(snd(snd\ x))$

definition *Guar* :: 'a rgformula \Rightarrow ('a \times 'a) set **where**
Guar $x \equiv snd(snd(snd\ x))$

definition *Com* :: 'a rgformula \Rightarrow 'a com **where**
Com $x \equiv fst\ x$

3.4.2 Proof System for Parallel Programs

type-synonym 'a par-rgformula =
('a rgformula) list \times 'a set \times ('a \times 'a) set \times ('a \times 'a) set \times 'a set

inductive

par-rghoare :: ('a rgformula) list \Rightarrow 'a set \Rightarrow ('a \times 'a) set \Rightarrow ('a \times 'a) set \Rightarrow 'a set \Rightarrow bool
 $\langle \vdash - SAT [-, -, -, -] [60,0,0,0,0] 45 \rangle$

where

Parallel:

$$\llbracket \forall i < \text{length } xs. \text{rely} \cup (\bigcup_{j \in \{j. j < \text{length } xs \wedge j \neq i\}}. \text{Guar}(xs!j)) \subseteq \text{Rely}(xs!i);$$

$$(\bigcup_{j \in \{j. j < \text{length } xs\}}. \text{Guar}(xs!j)) \subseteq \text{guar};$$

$$\text{pre} \subseteq (\bigcap_{i \in \{i. i < \text{length } xs\}}. \text{Pre}(xs!i));$$

$$(\bigcap_{i \in \{i. i < \text{length } xs\}}. \text{Post}(xs!i)) \subseteq \text{post};$$

$$\forall i < \text{length } xs. \vdash \text{Com}(xs!i) \text{ sat } [\text{Pre}(xs!i), \text{Rely}(xs!i), \text{Guar}(xs!i), \text{Post}(xs!i)] \rrbracket$$

$$\implies \vdash xs \text{ SAT } [\text{pre}, \text{rely}, \text{guar}, \text{post}]$$

3.5 Soundness

Some previous lemmas

lemma *tl-of-assum-in-assum:*

$$(P, s) \# (P, t) \# xs \in \text{assum}(\text{pre}, \text{rely}) \implies \text{stable pre rely}$$

$$\implies (P, t) \# xs \in \text{assum}(\text{pre}, \text{rely})$$
apply(*simp add:assum-def*)
apply *clarify*
apply(*rule conjI*)
apply(*erule-tac x=0 in allE*)
apply(*simp (no-asm-use) only:stable-def*)
apply(*erule allE,erule allE,erule impE,assumption,erule mp*)
apply(*simp add:Env*)
apply *clarify*
apply(*erule-tac x=Suc i in allE*)
apply *simp*
done

lemma *etran-in-comm:*

$$(P, t) \# xs \in \text{comm}(\text{guar}, \text{post}) \implies (P, s) \# (P, t) \# xs \in \text{comm}(\text{guar}, \text{post})$$
apply(*simp add:comm-def*)
apply *clarify*
apply(*case-tac i, simp+*)
done

lemma *ctran-in-comm:*

$$\llbracket (s, s) \in \text{guar}; (Q, s) \# xs \in \text{comm}(\text{guar}, \text{post}) \rrbracket$$

$$\implies (P, s) \# (Q, s) \# xs \in \text{comm}(\text{guar}, \text{post})$$
apply(*simp add:comm-def*)
apply *clarify*
apply(*case-tac i, simp+*)
done

lemma *takecptn-is-cptn* [*rule-format, elim!*]:

$$\forall j. c \in \text{cptn} \longrightarrow \text{take}(\text{Suc } j) c \in \text{cptn}$$
apply(*induct c*)
apply(*force elim: cptn.cases*)
apply *clarify*
apply(*case-tac j*)
apply *simp*

```

  apply(rule CptnOne)
  apply simp
  apply(force intro:cptn.intros elim:cptn.cases)
  done

```

```

lemma dropcptn-is-cptn [rule-format,elim!]:
   $\forall j < \text{length } c. c \in \text{cptn} \longrightarrow \text{drop } j \ c \in \text{cptn}$ 
  apply(induct c)
  apply(force elim: cptn.cases)
  apply clarify
  apply(case-tac j,simp+)
  apply(erule cptn.cases)
  apply simp
  apply force
  apply force
  done

```

```

lemma takepar-cptn-is-par-cptn [rule-format,elim]:
   $\forall j. c \in \text{par-cptn} \longrightarrow \text{take } (\text{Suc } j) \ c \in \text{par-cptn}$ 
  apply(induct c)
  apply(force elim: cptn.cases)
  apply clarify
  apply(case-tac j,simp)
  apply(rule ParCptnOne)
  apply(force intro:par-cptn.intros elim:par-cptn.cases)
  done

```

```

lemma droppar-cptn-is-par-cptn [rule-format]:
   $\forall j < \text{length } c. c \in \text{par-cptn} \longrightarrow \text{drop } j \ c \in \text{par-cptn}$ 
  apply(induct c)
  apply(force elim: par-cptn.cases)
  apply clarify
  apply(case-tac j,simp+)
  apply(erule par-cptn.cases)
  apply simp
  apply force
  apply force
  done

```

```

lemma tl-of-cptn-is-cptn:  $\llbracket x \# xs \in \text{cptn}; xs \neq [] \rrbracket \implies xs \in \text{cptn}$ 
  apply(subgoal-tac 1 < length (x # xs))
  apply(drule dropcptn-is-cptn,simp+)
  done

```

```

lemma not-ctran-None [rule-format]:
   $\forall s. (\text{None}, s) \# xs \in \text{cptn} \longrightarrow (\forall i < \text{length } xs. ((\text{None}, s) \# xs)!i -e \rightarrow xs!i)$ 
  apply(induct xs,simp+)
  apply clarify
  apply(erule cptn.cases,simp)

```

```

apply simp
apply(case-tac i, simp)
  apply(rule Env)
apply simp
apply(force elim:ctran.cases)
done

```

```

lemma cptn-not-empty [simp]:[]  $\notin$  cptn
apply(force elim:cptn.cases)
done

```

```

lemma etran-or-ctran [rule-format]:
   $\forall m i. x \in \text{cptn} \longrightarrow m \leq \text{length } x$ 
   $\longrightarrow (\forall i. \text{Suc } i < m \longrightarrow \neg x!i -c \rightarrow x!\text{Suc } i) \longrightarrow \text{Suc } i < m$ 
   $\longrightarrow x!i -e \rightarrow x!\text{Suc } i$ 
  supply [[simproc del: defined-all]]
apply(induct x, simp)
apply clarify
apply(erule cptn.cases, simp)
  apply(case-tac i, simp)
  apply(rule Env)
  apply simp
  apply(erule-tac x=m - 1 in allE)
  apply(case-tac m, simp, simp)
  apply(subgoal-tac ( $\forall i. \text{Suc } i < \text{ nata} \longrightarrow (((P, t) \# xs) ! i, xs ! i) \notin \text{ctran}$ ))
  apply force
  apply clarify
  apply(erule-tac x=Suc ia in allE, simp)
apply(erule-tac x=0 and P= $\lambda j. H j \longrightarrow (J j) \notin \text{ctran}$  for H J in allE, simp)
done

```

```

lemma etran-or-ctran2 [rule-format]:
   $\forall i. \text{Suc } i < \text{length } x \longrightarrow x \in \text{cptn} \longrightarrow (x!i -c \rightarrow x!\text{Suc } i \longrightarrow \neg x!i -e \rightarrow x!\text{Suc } i)$ 
   $\vee (x!i -e \rightarrow x!\text{Suc } i \longrightarrow \neg x!i -c \rightarrow x!\text{Suc } i)$ 
apply(induct x)
  apply simp
apply clarify
apply(erule cptn.cases, simp)
  apply(case-tac i, simp+)
apply(case-tac i, simp)
  apply(force elim:etran.cases)
apply simp
done

```

```

lemma etran-or-ctran2-disjI1:
  [ $x \in \text{cptn}; \text{Suc } i < \text{length } x; x!i -c \rightarrow x!\text{Suc } i$ ]  $\implies \neg x!i -e \rightarrow x!\text{Suc } i$ 
by(drule etran-or-ctran2, simp-all)

```

```

lemma etran-or-ctran2-disjI2:

```

$\llbracket x \in \text{cptn}; \text{Suc } i < \text{length } x; x!i -e \rightarrow x!\text{Suc } i \rrbracket \Longrightarrow \neg x!i -c \rightarrow x!\text{Suc } i$
by(*drule etran-or-ctran2,simp-all*)

lemma *not-ctran-None2* [*rule-format*]:

$\llbracket (\text{None}, s) \# xs \in \text{cptn}; i < \text{length } xs \rrbracket \Longrightarrow \neg ((\text{None}, s) \# xs) ! i -c \rightarrow xs ! i$
apply(*frule not-ctran-None,simp*)
apply(*case-tac i,simp*)
apply(*force elim:etranE*)
apply *simp*
apply(*rule etran-or-ctran2-disjI2,simp-all*)
apply(*force intro:tl-of-cptn-is-cptn*)
done

lemma *Ex-first-occurrence* [*rule-format*]: $P (n::\text{nat}) \longrightarrow (\exists m. P m \wedge (\forall i < m. \neg P i))$

apply(*rule nat-less-induct*)
apply *clarify*
apply(*case-tac $\forall m. m < n \longrightarrow \neg P m$*)
apply *auto*
done

lemma *stability* [*rule-format*]:

$\forall j k. x \in \text{cptn} \longrightarrow \text{stable } p \text{ rely} \longrightarrow j \leq k \longrightarrow k < \text{length } x \longrightarrow \text{snd}(x!j) \in p \longrightarrow$
 $(\forall i. (\text{Suc } i) < \text{length } x \longrightarrow$
 $(x!i -e \rightarrow x!(\text{Suc } i)) \longrightarrow (\text{snd}(x!i), \text{snd}(x!(\text{Suc } i))) \in \text{rely}) \longrightarrow$
 $(\forall i. j \leq i \wedge i < k \longrightarrow x!i -e \rightarrow x!\text{Suc } i) \longrightarrow \text{snd}(x!k) \in p \wedge \text{fst}(x!j) = \text{fst}(x!k)$
supply [*simproc del: defined-all*]
apply(*induct x*)
apply *clarify*
apply(*force elim:cptn.cases*)
apply *clarify*
apply(*erule cptn.cases,simp*)
apply *simp*
apply(*case-tac k,simp,simp*)
apply(*case-tac j,simp*)
apply(*erule-tac x=0 in allE*)
apply(*erule-tac x=nat and P= $\lambda j. (0 \leq j) \longrightarrow (J j)$ for J in allE,simp*)
apply(*subgoal-tac t \in p*)
apply(*subgoal-tac $(\forall i. i < \text{length } xs \longrightarrow ((P, t) \# xs) ! i -e \rightarrow xs ! i \longrightarrow (\text{snd}((P, t) \# xs) ! i), \text{snd}(xs ! i) \in \text{rely}))$*)
apply *clarify*
apply(*erule-tac x=Suc i and P= $\lambda j. (H j) \longrightarrow (J j) \in \text{etran}$ for H J in allE,simp*)
apply *clarify*
apply(*erule-tac x=Suc i and P= $\lambda j. (H j) \longrightarrow (J j) \longrightarrow (T j) \in \text{rely}$ for H J T in allE,simp*)
apply(*erule-tac x=0 and P= $\lambda j. (H j) \longrightarrow (J j) \in \text{etran} \longrightarrow T j$ for H J T in allE,simp*)
apply(*simp(no-asm-use) only:stable-def*)
apply(*erule-tac x=s in allE*)


```

apply(erule-tac x=t in allE)
apply simp
apply(erule mp)
apply(erule mp)
apply(rule Env)
apply simp
apply(erule-tac x=nata in allE)
apply(erule-tac x=nat and P= $\lambda j. (s \leq j) \longrightarrow (J j)$  for s J in allE,simp)
apply(subgoal-tac ( $\forall i. i < \text{length } xs \longrightarrow ((P, t) \# xs) ! i -e \longrightarrow xs ! i \longrightarrow (\text{snd } ((P, t) \# xs) ! i), \text{snd } (xs ! i)) \in \text{rely}$ )
apply clarify
apply(erule-tac x=Suc i and P= $\lambda j. (H j) \longrightarrow (J j) \in \text{etran}$  for H J in allE,simp)
apply clarify
apply(erule-tac x=Suc i and P= $\lambda j. (H j) \longrightarrow (J j) \longrightarrow (T j) \in \text{rely}$  for H J T in allE,simp)
apply(case-tac k,simp,simp)
apply(case-tac j)
apply(erule-tac x=0 and P= $\lambda j. (H j) \longrightarrow (J j) \in \text{etran}$  for H J in allE,simp)
apply(erule etran.cases,simp)
apply(erule-tac x=nata in allE)
apply(erule-tac x=nat and P= $\lambda j. (s \leq j) \longrightarrow (J j)$  for s J in allE,simp)
apply(subgoal-tac ( $\forall i. i < \text{length } xs \longrightarrow ((Q, t) \# xs) ! i -e \longrightarrow xs ! i \longrightarrow (\text{snd } ((Q, t) \# xs) ! i), \text{snd } (xs ! i)) \in \text{rely}$ )
apply clarify
apply(erule-tac x=Suc i and P= $\lambda j. (H j) \longrightarrow (J j) \in \text{etran}$  for H J in allE,simp)
apply clarify
apply(erule-tac x=Suc i and P= $\lambda j. (H j) \longrightarrow (J j) \longrightarrow (T j) \in \text{rely}$  for H J T in allE,simp)
done

```

3.5.1 Soundness of the System for Component Programs

Soundness of the Basic rule

lemma *unique-ctran-Basic* [rule-format]:

$$\begin{aligned}
& \forall s i. x \in \text{cptn} \longrightarrow x ! 0 = (\text{Some } (\text{Basic } f), s) \longrightarrow \\
& \text{Suc } i < \text{length } x \longrightarrow x ! i -c \longrightarrow x ! \text{Suc } i \longrightarrow \\
& (\forall j. \text{Suc } j < \text{length } x \longrightarrow i \neq j \longrightarrow x ! j -e \longrightarrow x ! \text{Suc } j)
\end{aligned}$$

```

apply(induct x,simp)
apply simp
apply clarify
apply(erule cptn.cases,simp)
apply(case-tac i,simp+)
apply clarify
apply(case-tac j,simp)
apply(rule Env)
apply simp
apply clarify
apply simp
apply(case-tac i)

```

```

apply(case-tac j,simp,simp)
apply(erule ctran.cases,simp-all)
apply(force elim: not-ctran-None)
apply(ind-cases ((Some (Basic f), sa), Q, t) ∈ ctran for sa Q t)
apply simp
apply(drule-tac i=nat in not-ctran-None,simp)
apply(erule etranE,simp)
done

```

```

lemma exists-ctran-Basic-None [rule-format]:
   $\forall s i. x \in \text{cptn} \longrightarrow x ! 0 = (\text{Some } (\text{Basic } f), s)$ 
   $\longrightarrow i < \text{length } x \longrightarrow \text{fst}(x!i) = \text{None} \longrightarrow (\exists j < i. x!j -c \rightarrow x! \text{Suc } j)$ 
apply(induct x,simp)
apply simp
apply clarify
apply(erule cptn.cases,simp)
apply(case-tac i,simp,simp)
apply(erule-tac x=nat in allE,simp)
apply clarify
apply(rule-tac x=Suc j in exI,simp,simp)
apply clarify
apply(case-tac i,simp,simp)
apply(rule-tac x=0 in exI,simp)
done

```

```

lemma Basic-sound:
   $\llbracket \text{pre} \subseteq \{s. f s \in \text{post}\}; \{(s, t). s \in \text{pre} \wedge t = f s\} \subseteq \text{guar};$ 
  stable pre rely; stable post rely
   $\implies \models \text{Basic } f \text{ sat } [\text{pre}, \text{rely}, \text{guar}, \text{post}]$ 
  supply [simproc del: defined-all]
apply(unfold com-validity-def)
apply clarify
apply(simp add:comm-def)
apply(rule conjI)
apply clarify
apply(simp add:cp-def assum-def)
apply clarify
apply(frule-tac j=0 and k=i and p=pre in stability)
  apply simp-all
  apply(erule-tac x=ia in allE,simp)
  apply(erule-tac i=i and f=f in unique-ctran-Basic,simp-all)
apply(erule subsetD,simp)
apply(case-tac x!i)
apply clarify
apply(drule-tac s=Some (Basic f) in sym,simp)
apply(thin-tac  $\forall j. H j$  for H)
apply(force elim:ctran.cases)
apply clarify
apply(simp add:cp-def)

```

```

apply clarify
apply(frule-tac  $i = \text{length } x - 1$  and  $f = f$  in exists-ctran-Basic-None,simp+)
  apply(case-tac  $x$ ,simp+)
    apply(rule last-fst-esp,simp add:last-length)
  apply (case-tac  $x$ ,simp+)
apply(simp add:assum-def)
apply clarify
apply(frule-tac  $j = 0$  and  $k = j$  and  $p = \text{pre}$  in stability)
  apply simp-all
  apply(erule-tac  $x = i$  in allE,simp)
  apply(erule-tac  $i = j$  and  $f = f$  in unique-ctran-Basic,simp-all)
apply(case-tac  $x!j$ )
apply clarify
apply simp
apply(drule-tac  $s = \text{Some } (Basic\ f)$  in sym,simp)
apply(case-tac  $x!Suc\ j$ ,simp)
apply(rule ctran.cases,simp)
apply(simp-all)
apply(drule-tac  $c = sa$  in subsetD,simp)
apply clarify
apply(frule-tac  $j = Suc\ j$  and  $k = \text{length } x - 1$  and  $p = \text{post}$  in stability,simp-all)
  apply(case-tac  $x$ ,simp+)
  apply(erule-tac  $x = i$  in allE)
apply(erule-tac  $i = j$  and  $f = f$  in unique-ctran-Basic,simp-all)
  apply arith+
apply(case-tac  $x$ )
apply(simp add:last-length)+
done

```

Soundness of the Await rule

lemma *unique-ctran-Await* [rule-format]:

$$\forall s\ i.\ x \in \text{cptn} \longrightarrow x!0 = (\text{Some } (\text{Await } b\ c),\ s) \longrightarrow$$

$$Suc\ i < \text{length } x \longrightarrow x!i -c \rightarrow x!Suc\ i \longrightarrow$$

$$(\forall j.\ Suc\ j < \text{length } x \longrightarrow i \neq j \longrightarrow x!j -e \rightarrow x!Suc\ j)$$

```

apply(induct  $x$ ,simp+)
apply clarify
apply(erule cptn.cases,simp)
  apply(case-tac  $i$ ,simp+)
  apply clarify
  apply(case-tac  $j$ ,simp)
  apply(rule Env)
apply simp
apply clarify
apply simp
apply(case-tac  $i$ )
  apply(case-tac  $j$ ,simp,simp)
apply(erule ctran.cases,simp-all)
apply(force elim: not-ctran-None)

```

apply(*ind-cases* ((*Some* (*Await* *b c*), *sa*), *Q*, *t*) \in *ctran* **for** *sa Q t,simp*)
apply(*drule-tac* *i=nat* **in** *not-ctran-None,simp*)
apply(*erule etranE,simp*)
done

lemma *exists-ctran-Await-None* [*rule-format*]:
 $\forall s i. x \in \text{cptn} \longrightarrow x ! 0 = (\text{Some } (\text{Await } b c), s)$
 $\longrightarrow i < \text{length } x \longrightarrow \text{fst}(x!i) = \text{None} \longrightarrow (\exists j < i. x!j -c \rightarrow x! \text{Suc } j)$
apply(*induct* *x,simp+*)
apply *clarify*
apply(*erule cptn.cases,simp*)
apply(*case-tac* *i,simp+*)
apply(*erule-tac* *x=nat* **in** *allE,simp*)
apply *clarify*
apply(*rule-tac* *x=Suc j* **in** *exI,simp,simp*)
apply *clarify*
apply(*case-tac* *i,simp,simp*)
apply(*rule-tac* *x=0* **in** *exI,simp*)
done

lemma *Star-imp-cptn*:
 $(P, s) -c^* \rightarrow (R, t) \implies \exists l \in \text{cp } P \text{ s. } (\text{last } l) = (R, t)$
 $\wedge (\forall i. \text{Suc } i < \text{length } l \longrightarrow !i -c \rightarrow ! \text{Suc } i)$
apply (*erule converse-rtrancl-induct2*)
apply(*rule-tac* *x=[(R,t)]* **in** *bestI*)
apply *simp*
apply(*simp add:cp-def*)
apply(*rule CptnOne*)
apply *clarify*
apply(*rule-tac* *x=(a, b)#l* **in** *bestI*)
apply (*rule conjI*)
apply(*case-tac* *l,simp add:cp-def*)
apply(*simp add:last-length*)
apply *clarify*
apply(*case-tac* *i,simp*)
apply(*simp add:cp-def*)
apply *force*
apply(*simp add:cp-def*)
apply(*case-tac* *l*)
apply(*force elim:cptn.cases*)
apply *simp*
apply(*erule CptnComp*)
apply *clarify*
done

lemma *Await-sound*:
 $\llbracket \text{stable pre rely; stable post rely;}$
 $\forall V. \vdash P \text{ sat } [\text{pre} \cap b \cap \{s. s = V\}, \{(s, t). s = t\},$
 $\text{UNIV}, \{s. (V, s) \in \text{guar}\} \cap \text{post}] \wedge$

```

  ⊨ P sat [pre ∩ b ∩ {s. s = V}, {(s, t). s = t},
           UNIV, {s. (V, s) ∈ guar} ∩ post] ]
  ⇒ ⊨ Await b P sat [pre, rely, guar, post]
apply(unfold com-validity-def)
apply clarify
apply(simp add:comm-def)
apply(rule conjI)
apply clarify
apply(simp add:cp-def assum-def)
apply clarify
apply(frule-tac j=0 and k=i and p=pre in stability,simp-all)
  apply(erule-tac x=ia in allE,simp)
  apply(subgoal-tac x ∈ cp (Some(Await b P)) s)
  apply(erule-tac i=i in unique-ctran-Await,force,simp-all)
  apply(simp add:cp-def)
— here starts the different part.
apply(erule ctran.cases,simp-all)
apply(drule Star-imp-cptn)
apply clarify
apply(erule-tac x=sa in allE)
apply clarify
apply(erule-tac x=sa in allE)
apply(drule-tac c=l in subsetD)
  apply (simp add:cp-def)
  apply clarify
apply(erule-tac x=ia and P=λi. H i → (J i, I i) ∈ ctran for H J I in allE,simp)
  apply(erule etranE,simp)
apply simp
apply clarify
apply(simp add:cp-def)
apply clarify
apply(frule-tac i=length x - 1 in exists-ctran-Await-None,force)
  apply (case-tac x,simp+)
  apply(rule last-fst-esp,simp add:last-length)
  apply(case-tac x, simp+)
apply clarify
apply(simp add:assum-def)
apply clarify
apply(frule-tac j=0 and k=j and p=pre in stability,simp-all)
  apply(erule-tac x=i in allE,simp)
  apply(erule-tac i=j in unique-ctran-Await,force,simp-all)
apply(case-tac x!j)
apply clarify
apply simp
apply(drule-tac s=Some (Await b P) in sym,simp)
apply(case-tac x!Suc j,simp)
apply(rule ctran.cases,simp)
apply(simp-all)
apply(drule Star-imp-cptn)

```

```

apply clarify
apply(erule-tac  $x=sa$  in allE)
apply clarify
apply(erule-tac  $x=sa$  in allE)
apply(drule-tac  $c=l$  in subsetD)
  apply (simp add:cp-def)
  apply clarify
  apply(erule-tac  $x=i$  and  $P=\lambda i. H i \longrightarrow (J i, I i) \in ctran$  for  $H J I$  in allE, simp)
  apply(erule etranE, simp)
apply simp
apply clarify
apply(frule-tac  $j=Suc j$  and  $k=length\ x - 1$  and  $p=post$  in stability, simp-all)
  apply(case-tac  $x, simp+$ )
  apply(erule-tac  $x=i$  in allE)
apply(erule-tac  $i=j$  in unique-ctran-Await, force, simp-all)
  apply arith+
apply(case-tac  $x$ )
apply(simp add:last-length)+
done

```

Soundness of the Conditional rule

lemma *Cond-sound*:

```

  [| stable  $pre\ rely$ ;  $\models P1\ sat\ [pre \cap b, rely, guar, post]$ ;
  |  $\models P2\ sat\ [pre \cap \neg b, rely, guar, post]$ ;  $\forall s. (s,s) \in guar$  |]
   $\implies \models (Cond\ b\ P1\ P2)\ sat\ [pre, rely, guar, post]$ 
apply(unfold com-validity-def)
apply clarify
apply(simp add:cp-def comm-def)
apply(case-tac  $\exists i. Suc\ i < length\ x \wedge x!i - c \rightarrow x!Suc\ i$ )
  prefer 2
  apply simp
  apply clarify
apply(frule-tac  $j=0$  and  $k=length\ x - 1$  and  $p=pre$  in stability, simp+)
  apply(case-tac  $x, simp+$ )
  apply(simp add:assum-def)
  apply(simp add:assum-def)
  apply(erule-tac  $m=length\ x$  in etran-or-ctran, simp+)
apply(case-tac  $x, (simp\ add:last-length)$ +)
apply(erule exE)
apply(drule-tac  $n=i$  and  $P=\lambda i. H i \wedge (J i, I i) \in ctran$  for  $H J I$  in Ex-first-occurrence)
apply clarify
apply (simp add:assum-def)
apply(frule-tac  $j=0$  and  $k=m$  and  $p=pre$  in stability, simp+)
apply(erule-tac  $m=Suc\ m$  in etran-or-ctran, simp+)
apply(erule ctran.cases, simp-all)
apply(erule-tac  $x=sa$  in allE)
apply(drule-tac  $c=drop\ (Suc\ m)\ x$  in subsetD)
  apply simp

```

```

  apply clarify
  apply simp
  apply clarify
  apply(case-tac i≤m)
  apply(drule le-imp-less-or-eq)
  apply(erule disjE)
  apply(erule-tac x=i in allE, erule impE, assumption)
  apply simp+
  apply(erule-tac x=i - (Suc m) and P=λj. H j → J j → (I j)∈guar for H J
  I in allE)
  apply(subgoal-tac (Suc m)+(i - Suc m) ≤ length x)
  apply(subgoal-tac (Suc m)+Suc (i - Suc m) ≤ length x)
  apply(rotate-tac -2)
  apply simp
  apply arith
  apply arith
  apply(case-tac length (drop (Suc m) x),simp)
  apply(erule-tac x=sa in allE)
  back
  apply(drule-tac c=drop (Suc m) x in subsetD,simp)
  apply clarify
  apply simp
  apply clarify
  apply(case-tac i≤m)
  apply(drule le-imp-less-or-eq)
  apply(erule disjE)
  apply(erule-tac x=i in allE, erule impE, assumption)
  apply simp
  apply simp
  apply(erule-tac x=i - (Suc m) and P=λj. H j → J j → (I j)∈guar for H J I
  in allE)
  apply(subgoal-tac (Suc m)+(i - Suc m) ≤ length x)
  apply(subgoal-tac (Suc m)+Suc (i - Suc m) ≤ length x)
  apply(rotate-tac -2)
  apply simp
  apply arith
  apply arith
  done

```

Soundness of the Sequential rule

inductive-cases *Seq-cases* [*elim!*]: $(Some (Seq P Q), s) -c \rightarrow t$

lemma *last-lift-not-None*: $fst ((lift Q) ((x\#xs)!(length xs))) \neq None$

```

  apply(subgoal-tac length xs<length (x # xs))
  apply(drule-tac Q=Q in lift-nth)
  apply(erule ssubst)
  apply (simp add:lift-def)
  apply(case-tac (x # xs) ! length xs,simp)

```

apply simp
done

lemma Seq-sound1 [rule-format]:

$x \in \text{cptn-mod} \implies \forall s P. x \neq (\text{Some } (\text{Seq } P \ Q), s) \longrightarrow$
 $(\forall i < \text{length } x. \text{fst}(x!i) \neq \text{Some } Q) \longrightarrow$
 $(\exists xs \in \text{cp } (\text{Some } P) s. x = \text{map } (\text{lift } Q) \ xs)$
supply [[simproc del: defined-all]]
apply(erule cptn-mod.induct)
apply(unfold cp-def)
apply safe
apply simp-all
apply(simp add:lift-def)
apply(rule-tac x=[(Some Pa, sa)] in exI,simp add:CptnOne)
apply(subgoal-tac ($\forall i < \text{Suc } (\text{length } xs). \text{fst } (((\text{Some } (\text{Seq } Pa \ Q), t) \# xs) ! i) \neq \text{Some } Q$))
apply clarify
apply(rule-tac x=(Some Pa, sa) # (Some Pa, t) # zs in exI,simp)
apply(rule conjI,erule CptnEnv)
apply(simp (no-asm-use) add:lift-def)
apply clarify
apply(erule-tac x=Suc i in allE, simp)
apply(ind-cases ((Some (Seq Pa Q), sa), None, t) \in ctran for Pa sa t)
apply(rule-tac x=(Some P, sa) # xs in exI, simp add:cptn-iff-cptn-mod lift-def)
apply(erule-tac x=length xs in allE, simp)
apply(simp only:Cons-lift-append)
apply(subgoal-tac length xs < length ((Some P, sa) # xs))
apply(simp only :nth-append length-map last-length nth-map)
apply(case-tac last((Some P, sa) # xs))
apply(simp add:lift-def)
apply simp
done

lemma Seq-sound2 [rule-format]:

$x \in \text{cptn} \implies \forall s P i. x \neq (\text{Some } (\text{Seq } P \ Q), s) \longrightarrow i < \text{length } x$
 $\longrightarrow \text{fst}(x!i) = \text{Some } Q \longrightarrow$
 $(\forall j < i. \text{fst}(x!j) \neq (\text{Some } Q)) \longrightarrow$
 $(\exists xs \ ys. xs \in \text{cp } (\text{Some } P) s \wedge \text{length } xs = \text{Suc } i$
 $\wedge ys \in \text{cp } (\text{Some } Q) (\text{snd}(xs !i)) \wedge x = (\text{map } (\text{lift } Q) \ xs) @ \text{tl } ys)$
supply [[simproc del: defined-all]]
apply(erule cptn.induct)
apply(unfold cp-def)
apply safe
apply simp-all
apply(case-tac i,simp+)
apply(erule allE,erule impE,assumption,simp)
apply clarify
apply(subgoal-tac ($\forall j < \text{nat}. \text{fst } (((\text{Some } (\text{Seq } Pa \ Q), t) \# xs) ! j) \neq \text{Some } Q$),clarify)


```

prefer 2
apply force
apply(case-tac xsa,simp,simp)
apply(rename-tac list)
apply(rule-tac x=(Some Pa, sa) # (Some Pa, t) # list in exI,simp)
apply(rule conjI,erule CptnEnv)
apply(simp (no-asm-use) add:lift-def)
apply(rule-tac x=ys in exI,simp)
apply(ind-cases ((Some (Seq Pa Q), sa), t) ∈ ctran for Pa sa t)
apply simp
apply(rule-tac x=(Some Pa, sa)#[(None, ta)] in exI,simp)
apply(rule conjI)
  apply(drule-tac xs=[] in CptnComp,force simp add:CptnOne,simp)
apply(case-tac i, simp+)
apply(case-tac nat,simp+)
apply(rule-tac x=(Some Q,ta)#xs in exI,simp add:lift-def)
apply(case-tac nat,simp+)
apply(force)
apply(case-tac i, simp+)
apply(case-tac nat,simp+)
apply(erule-tac x=Suc nata in allE,simp)
apply clarify
apply(subgoal-tac (∀ j<Suc nata. fst (((Some (Seq P2 Q), ta) # xs) ! j) ≠ Some Q),clarify)
prefer 2
apply clarify
apply force
apply(rule-tac x=(Some Pa, sa)#(Some P2, ta)#(tl xsa) in exI,simp)
apply(rule conjI,erule CptnComp)
apply(rule nth-tl-if,force,simp+)
apply(rule-tac x=ys in exI,simp)
apply(rule conjI)
apply(rule nth-tl-if,force,simp+)
  apply(rule tl-zero,simp+)
  apply force
apply(rule conjI,simp add:lift-def)
apply(subgoal-tac lift Q (Some P2, ta) =(Some (Seq P2 Q), ta))
  apply(simp add:Cons-lift del:list.map)
  apply(rule nth-tl-if)
    apply force
    apply simp+
apply(simp add:lift-def)
done

```

```

lemma last-lift-not-None2: fst ((lift Q) (last (x#xs))) ≠ None
apply(simp only:last-length [THEN sym])
apply(subgoal-tac length xs<length (x # xs))
  apply(drule-tac Q=Q in lift-nth)

```

```

apply(erule ssubst)
apply (simp add:lift-def)
apply(case-tac (x # xs) ! length xs,simp)
apply simp
done

lemma Seq-sound:
  [[ $\models P \text{ sat } [pre, rely, guar, mid]; \models Q \text{ sat } [mid, rely, guar, post]$ ]]
   $\implies \models \text{Seq } P \ Q \text{ sat } [pre, rely, guar, post]$ 
apply(unfold com-validity-def)
apply clarify
apply(case-tac  $\exists i < \text{length } x. \text{fst}(x!i) = \text{Some } Q$ )
prefer 2
apply (simp add:cp-def cptn-iff-cptn-mod)
apply clarify
apply(erule-tac Seq-sound1,force)
apply force
apply clarify
apply(erule-tac  $x=s$  in allE,simp)
apply(erule-tac  $c=xs$  in subsetD,simp add:cp-def cptn-iff-cptn-mod)
apply(simp add:assum-def)
apply clarify
apply(erule-tac  $P=\lambda j. H \ j \longrightarrow J \ j \longrightarrow I \ j$  for H J I in allE,erule impE,
assumption)
apply(simp add:snd-lift)
apply(erule mp)
apply(force elim:etranE intro:Env simp add:lift-def)
apply(simp add:comm-def)
apply(rule conjI)
apply clarify
apply(erule-tac  $P=\lambda j. H \ j \longrightarrow J \ j \longrightarrow I \ j$  for H J I in allE,erule impE,
assumption)
apply(simp add:snd-lift)
apply(erule mp)
apply(case-tac (xs!i))
apply(case-tac (xs! Suc i))
apply(case-tac fst(xs!i))
apply(erule-tac  $x=i$  in allE, simp add:lift-def)
apply(case-tac fst(xs!Suc i))
apply(force simp add:lift-def)
apply(force simp add:lift-def)
apply clarify
apply(case-tac xs,simp add:cp-def)
apply clarify
apply (simp del:list.map)
apply (rename-tac list)
apply(subgoal-tac (map (lift Q) ((a, b) # list)) $\neq$ [])
apply(erule last-conv-nth)
apply (simp del:list.map)

```

```

  apply(simp only:last-lift-not-None)
  apply simp
  —  $\exists i < \text{length } x. \text{fst } (x ! i) = \text{Some } Q$ 
  apply(erule exE)
  apply(drule-tac n=i and P= $\lambda i. i < \text{length } x \wedge \text{fst } (x ! i) = \text{Some } Q$  in Ex-first-occurrence)
  apply clarify
  apply (simp add:cp-def)
  apply clarify
  apply(frule-tac i=m in Seq-sound2,force)
  apply simp+
  apply clarify
  apply(simp add:comm-def)
  apply(erule-tac x=s in allE)
  apply(drule-tac c=xs in subsetD,simp)
  apply(case-tac xs=[],simp)
  apply(simp add:cp-def assum-def nth-append)
  apply clarify
  apply(erule-tac x=i in allE)
  back
  apply(simp add:snd-lift)
  apply(erule mp)
  apply(force elim:etranE intro:Env simp add:lift-def)
  apply simp
  apply clarify
  apply(erule-tac x=snd(xs!m) in allE)
  apply(drule-tac c=ys in subsetD,simp add:cp-def assum-def)
  apply(case-tac xs $\neq$ [])
  apply(drule last-conv-nth,simp)
  apply(rule conjI)
  apply(erule mp)
  apply(case-tac xs!m)
  apply(case-tac fst(xs!m),simp)
  apply(simp add:lift-def nth-append)
  apply clarify
  apply(erule-tac x=m+i in allE)
  back
  back
  apply(case-tac ys,(simp add:nth-append)+)
  apply (case-tac i, (simp add:snd-lift)+)
  apply(erule mp)
  apply(case-tac xs!m)
  apply(force elim:etran.cases intro:Env simp add:lift-def)
  apply simp
  apply simp
  apply clarify
  apply(rule conjI,clarify)
  apply(case-tac i<m,simp add:nth-append)
  apply(simp add:snd-lift)
  apply(erule allE, erule impE, assumption, erule mp)

```

```

apply(case-tac (xs ! i))
apply(case-tac (xs ! Suc i))
apply(case-tac fst(xs ! i),force simp add:lift-def)
apply(case-tac fst(xs ! Suc i))
  apply (force simp add:lift-def)
  apply (force simp add:lift-def)
apply(erule-tac x=i-m in allE)
back
back
apply(subgoal-tac Suc (i - m) < length ys,simp)
prefer 2
apply arith
apply(simp add:nth-append snd-lift)
apply(rule conjI,clarify)
apply(subgoal-tac i=m)
  prefer 2
  apply arith
  apply clarify
apply(simp add:cp-def)
apply(rule tl-zero)
  apply(erule mp)
  apply(case-tac lift Q (xs!m),simp add:snd-lift)
  apply(case-tac xs!m,case-tac fst(xs!m),simp add:lift-def snd-lift)
  apply(case-tac ys,simp+)
  apply(simp add:lift-def)
  apply simp
  apply force
apply clarify
apply(rule tl-zero)
  apply(rule tl-zero)
    apply (subgoal-tac i-m=Suc(i-Suc m))
    apply simp
    apply(erule mp)
    apply(case-tac ys,simp+)
  apply force
  apply arith
apply force
apply clarify
apply(case-tac (map (lift Q) xs @ tl ys)≠[])
apply(drule last-conv-nth)
apply(simp add: snd-lift nth-append)
apply(rule conjI,clarify)
  apply(case-tac ys,simp+)
apply clarify
apply(case-tac ys,simp+)
done

```

Soundness of the While rule

lemma *last-append*[*rule-format*]:
 $\forall xs. ys \neq [] \longrightarrow ((xs @ ys)!(length (xs @ ys) - (Suc 0))) = (ys!(length ys - (Suc 0)))$
apply (*induct* *ys*)
apply *simp*
apply *clarify*
apply (*simp* *add:nth-append*)
done

lemma *assum-after-body*:
 $\llbracket \models P \text{ sat } [pre \cap b, \text{rely}, \text{guar}, \text{pre}];$
 $(Some\ P, s) \# xs \in \text{cptn-mod}; \text{fst}(\text{last}((Some\ P, s) \# xs)) = \text{None}; s \in b;$
 $(Some\ (\text{While}\ b\ P), s) \# (Some\ (\text{Seq}\ P\ (\text{While}\ b\ P)), s) \#$
 $\text{map}(\text{lift}\ (\text{While}\ b\ P))\ xs\ @\ ys \in \text{assum}\ (\text{pre}, \text{rely}) \rrbracket$
 $\implies (Some\ (\text{While}\ b\ P), \text{snd}(\text{last}((Some\ P, s) \# xs))) \# ys \in \text{assum}\ (\text{pre}, \text{rely})$
apply (*simp* *add:assum-def* *com-validity-def* *cp-def* *cptn-iff-cptn-mod*)
apply *clarify*
apply (*erule-tac* *x=s* **in** *allE*)
apply (*drule-tac* *c=(Some P, s) # xs* **in** *subsetD, simp*)
apply *clarify*
apply (*erule-tac* *x=Suc i* **in** *allE*)
apply *simp*
apply (*simp* *add:Cons-lift-append* *nth-append* *snd-lift* *del:list.map*)
apply (*erule* *mp*)
apply (*erule* *etranE, simp*)
apply (*case-tac* *fst(((Some P, s) # xs) ! i)*)
apply (*force* *intro:Env* *simp* *add:lift-def*)
apply (*force* *intro:Env* *simp* *add:lift-def*)
apply (*rule* *conjI*)
apply *clarify*
apply (*simp* *add:comm-def* *last-length*)
apply *clarify*
apply (*rule* *conjI*)
apply (*simp* *add:comm-def*)
apply *clarify*
apply (*erule-tac* *x=Suc(length xs + i)* **in** *allE, simp*)
apply (*case-tac* *i, simp* *add:nth-append* *Cons-lift-append* *snd-lift* *last-conv-nth* *lift-def* *split-def*)
apply (*simp* *add:Cons-lift-append* *nth-append* *snd-lift*)
done

lemma *While-sound-aux* [*rule-format*]:
 $\llbracket pre \cap - b \subseteq post; \models P \text{ sat } [pre \cap b, \text{rely}, \text{guar}, \text{pre}]; \forall s. (s, s) \in \text{guar};$
 $\text{stable}\ pre\ \text{rely}; \text{stable}\ post\ \text{rely}; x \in \text{cptn-mod} \rrbracket$
 $\implies \forall s\ xs. x = (Some(\text{While}\ b\ P), s) \# xs \longrightarrow x \in \text{assum}(pre, \text{rely}) \longrightarrow x \in \text{comm}$
 $(\text{guar}, \text{post})$
supply [*simproc* *del: defined-all*]
apply (*erule* *cptn-mod.induct*)
apply *safe*

```

apply (simp-all del:last.simps)
— 5 subgoals left
apply(simp add:comm-def)
— 4 subgoals left
apply(rule etran-in-comm)
apply(erule mp)
apply(erule tl-of-assum-in-assum, simp)
— While-None
apply(ind-cases ((Some (While b P)), s), None, t) in ctran for s t)
apply(simp add:comm-def)
apply(simp add:cptn-iff-cptn-mod [THEN sym])
apply(rule conjI, clarify)
  apply(force simp add:assum-def)
apply clarify
apply(rule conjI, clarify)
  apply(case-tac i, simp, simp)
  apply(force simp add:not-ctran-None2)
apply(subgoal-tac  $\forall i. \text{Suc } i < \text{length } ((\text{None}, t) \# xs) \longrightarrow (((\text{None}, t) \# xs) ! i, (\text{None}, t) \# xs) ! \text{Suc } i) \in \text{etran}$ )
  prefer 2
  apply clarify
  apply(rule-tac m=length ((None, s)  $\#$  xs) in etran-or-ctran, simp+)
  apply(erule not-ctran-None2, simp)
  apply simp+
apply(frule-tac j=0 and k=length ((None, s)  $\#$  xs)  $- 1$  and p=post in stability, simp+)
  apply(force simp add:assum-def subsetD)
  apply(simp add:assum-def)
  apply clarify
  apply(erule-tac x=i in allE, simp)
  apply(erule-tac x=Suc i in allE, simp)
  apply simp
apply clarify
apply (simp add:last-length)
— WhileOne
apply(rule ctran-in-comm, simp)
apply(simp add:Cons-lift del:list.map)
apply(simp add:comm-def del:list.map)
apply(rule conjI)
apply clarify
apply(case-tac fst((Some P, sa)  $\#$  xs) ! i)
apply(case-tac ((Some P, sa)  $\#$  xs) ! i)
apply (simp add:lift-def)
apply(ind-cases (Some (While b P)), ba)  $-c \rightarrow t$  for ba t)
  apply simp
  apply simp
apply(simp add:snd-lift del:list.map)
apply(simp only:com-validity-def cp-def cptn-iff-cptn-mod)
apply(erule-tac x=sa in allE)

```

```

apply(drule-tac  $c=(\text{Some } P, sa) \# xs$  in subsetD)
apply (simp add:assum-def del:list.map)
apply clarify
apply(erule-tac  $x=\text{Suc } ia$  in allE, simp add:snd-lift del:list.map)
apply(erule mp)
apply(case-tac  $\text{fst}(((\text{Some } P, sa) \# xs) ! ia)$ )
  apply(erule etranE, simp add:lift-def)
  apply(rule Env)
apply(erule etranE, simp add:lift-def)
apply(rule Env)
apply (simp add:comm-def del:list.map)
apply clarify
apply(erule allE, erule impE, assumption)
apply(erule mp)
apply(case-tac  $((\text{Some } P, sa) \# xs) ! i$ )
apply(case-tac  $xs!i$ )
apply(simp add:lift-def)
apply(case-tac  $\text{fst}(xs!i)$ )
  apply force
apply force
— last=None
apply clarify
apply(subgoal-tac ( $\text{map } (\text{lift } (\text{While } b P)) ((\text{Some } P, sa) \# xs) \neq []$ ))
  apply(drule last-conv-nth)
  apply (simp del:list.map)
  apply(simp only:last-lift-not-None)
apply simp
— WhileMore
apply(rule ctran-in-comm, simp del:last.simps)
— metiendo la hipotesis antes de dividir la conclusion.
apply(subgoal-tac ( $\text{Some } (\text{While } b P), \text{snd } (\text{last } ((\text{Some } P, sa) \# xs)) \# ys \in$ 
assum (pre, rely))
  apply (simp del:last.simps)
prefer 2
apply(erule assum-after-body)
  apply (simp del:last.simps)+
— lo de antes.
apply(simp add:comm-def del:list.map last.simps)
apply(rule conjI)
apply clarify
apply(simp only:Cons-lift-append)
apply(case-tac  $i < \text{length } xs$ )
apply(simp add:nth-append del:list.map last.simps)
apply(case-tac  $\text{fst}(((\text{Some } P, sa) \# xs) ! i)$ )
  apply(case-tac  $((\text{Some } P, sa) \# xs) ! i$ )
  apply (simp add:lift-def del:last.simps)
apply(ind-cases ( $\text{Some } (\text{While } b P), ba$ )  $-c \rightarrow t$  for  $ba t$ )
  apply simp
apply simp

```

```

apply(simp add:snd-lift del:list.map last.simps)
apply(thin-tac  $\forall i. i < \text{length } ys \rightarrow P i$  for P)
apply(simp only:com-validity-def cp-def cptn-iff-cptn-mod)
apply(erule-tac x=sa in allE)
apply(drule-tac c=(Some P, sa) # xs in subsetD)
  apply (simp add:assum-def del:list.map last.simps)
  apply clarify
    apply(erule-tac x=Suc ia in allE,simp add:nth-append snd-lift del:list.map
last.simps, erule mp)
    apply(case-tac fst(((Some P, sa) # xs) ! ia))
      apply(erule etranE,simp add:lift-def)
      apply(rule Env)
    apply(erule etranE,simp add:lift-def)
    apply(rule Env)
apply (simp add:comm-def del:list.map)
apply clarify
apply(erule allE,erule impE,assumption)
apply(erule mp)
apply(case-tac ((Some P, sa) # xs) ! i)
apply(case-tac xs!i)
apply(simp add:lift-def)
apply(case-tac fst(xs!i))
  apply force
apply force
—  $i \geq \text{length } xs$ 
apply(subgoal-tac  $i - \text{length } xs < \text{length } ys$ )
  prefer 2
  apply arith
apply(erule-tac x= $i - \text{length } xs$  in allE,clarify)
apply(case-tac  $i = \text{length } xs$ )
  apply (simp add:nth-append snd-lift del:list.map last.simps)
  apply(simp add:last-length del:last.simps)
  apply(erule mp)
  apply(case-tac last((Some P, sa) # xs))
  apply(simp add:lift-def del:last.simps)
—  $i > \text{length } xs$ 
apply(case-tac  $i - \text{length } xs$ )
  apply arith
apply(simp add:nth-append del:list.map last.simps)
apply(rotate-tac -3)
apply(subgoal-tac  $i - \text{Suc } (\text{length } xs) = \text{nat}$ )
  prefer 2
  apply arith
apply simp
— last=None
apply clarify
apply(case-tac ys)
  apply(simp add:Cons-lift del:list.map last.simps)
  apply(subgoal-tac (map (lift (While b P)) ((Some P, sa) # xs))≠[])

```



```

apply(drule last-conv-nth)
apply (simp del:list.map)
apply(simp only:last-lift-not-None)
apply simp
apply(subgoal-tac ((Some (Seq P (While b P))), sa) # map (lift (While b P)) xs
@ ys)≠[])
apply(drule last-conv-nth)
apply (simp del:list.map last.simps)
apply(simp add:nth-append del:last.simps)
apply(rename-tac a list)
apply(subgoal-tac ((Some (While b P)), snd (last ((Some P, sa) # xs))) # a #
list)≠[])
apply(drule last-conv-nth)
apply (simp del:list.map last.simps)
apply simp
apply simp
done

```

lemma *While-sound*:

```

[[stable pre rely; pre ∩ b ⊆ post; stable post rely;
  ⊨ P sat [pre ∩ b, rely, guar, pre]; ∀ s. (s,s)∈guar]
⇒ ⊨ While b P sat [pre, rely, guar, post]
apply(unfold com-validity-def)
apply clarify
apply(erule-tac xs=tl x in While-sound-aux)
apply(simp add:com-validity-def)
apply force
apply simp-all
apply(simp add:cptn-iff-cptn-mod cp-def)
apply(simp add:cp-def)
apply clarify
apply(rule nth-equalityI)
apply simp-all
apply(case-tac x, simp+)
apply(case-tac i, simp+)
apply(case-tac x, simp+)
done

```

Soundness of the Rule of Consequence

lemma *Conseq-sound*:

```

[[pre ⊆ pre'; rely ⊆ rely'; guar' ⊆ guar; post' ⊆ post;
  ⊨ P sat [pre', rely', guar', post']]
⇒ ⊨ P sat [pre, rely, guar, post]
apply(simp add:com-validity-def assum-def comm-def)
apply clarify
apply(erule-tac x=s in allE)
apply(drule-tac c=x in subsetD)
apply force

```

apply *force*
done

Soundness of the system for sequential component programs

theorem *rgsound*:

$\vdash P \text{ sat } [pre, rely, guar, post] \implies \models P \text{ sat } [pre, rely, guar, post]$
apply(*erule* *rghoare.induct*)
apply(*force* *elim:Basic-sound*)
apply(*force* *elim:Seq-sound*)
apply(*force* *elim:Cond-sound*)
apply(*force* *elim:While-sound*)
apply(*force* *elim:Await-sound*)
apply(*erule* *Conseq-sound,simp+*)
done

3.5.2 Soundness of the System for Parallel Programs

definition *ParallelCom* :: ('a *rgformula*) *list* \Rightarrow 'a *par-com* **where**
ParallelCom *Ps* \equiv *map* (*Some* \circ *fst*) *Ps*

lemma *two*:

$\llbracket \forall i < \text{length } xs. \text{ rely} \cup (\bigcup j \in \{j. j < \text{length } xs \wedge j \neq i\}. \text{ Guar } (xs ! j))$
 $\subseteq \text{ Rely } (xs ! i);$
 $pre \subseteq (\bigcap i \in \{i. i < \text{length } xs\}. \text{ Pre } (xs ! i));$
 $\forall i < \text{length } xs.$
 $\models \text{ Com } (xs ! i) \text{ sat } [\text{Pre } (xs ! i), \text{ Rely } (xs ! i), \text{ Guar } (xs ! i), \text{ Post } (xs ! i)];$
 $\text{length } xs = \text{length } \text{clist}; x \in \text{par-cp } (\text{ParallelCom } xs) \text{ } s; x \in \text{par-assum}(pre, rely);$
 $\forall i < \text{length } \text{clist}. \text{clist} ! i \in \text{cp } (\text{Some}(\text{Com}(xs ! i))) \text{ } s; x \propto \text{clist} \rrbracket$
 $\implies \forall j i. i < \text{length } \text{clist} \wedge \text{Suc } j < \text{length } x \longrightarrow (\text{clist} ! i ! j) -c \rightarrow (\text{clist} ! i ! \text{Suc } j)$
 $\longrightarrow (\text{snd}(\text{clist} ! i ! j), \text{snd}(\text{clist} ! i ! \text{Suc } j)) \in \text{Guar}(xs ! i)$

apply(*unfold* *par-cp-def*)

apply (*rule* *ccontr*)

— By contradiction:

apply *simp*

apply(*erule* *exE*)

— the first c-tran that does not satisfy the guarantee-condition is from σ -*i* at step *m*.

apply(*drule-tac* *n=j and P= λ j. $\exists i. H i j$ for H in Ex-first-occurrence*)

apply(*erule* *exE*)

apply *clarify*

— σ -*i* $\in A(pre, rely-1)$

apply(*subgoal-tac* *take* (*Suc* (*Suc* *m*)) (*clist* ! *i*) \in *assum*(*Pre*(*xs* ! *i*), *Rely*(*xs* ! *i*)))

— but this contradicts $\models \sigma$ -*i* *sat* [*pre-i*, *rely-i*, *guar-i*, *post-i*]

apply(*erule-tac* *x=i and P= λ i. H i $\longrightarrow \models (J i) \text{ sat } [I i, K i, M i, N i]$ for H J I*
K M N in allE,erule impE,assumption)

apply(*simp* *add:com-validity-def*)

apply(*erule-tac* *x=s in allE*)

apply(*simp* *add:cp-def comm-def*)

apply(*drule-tac* *c=take* (*Suc* (*Suc* *m*)) (*clist* ! *i*) **in** *subsetD*)

apply *simp*
apply (*blast intro: takecptn-is-cptn*)
apply *simp*
apply *clarify*
apply(*erule-tac x=m and P= $\lambda j. I j \wedge J j \longrightarrow H j$ for $I J H$ in $allE$)
apply (*simp add:conjoin-def same-length-def*)
apply(*simp add:assum-def*)
apply(*rule conjI*)
apply(*erule-tac x=i and P= $\lambda j. H j \longrightarrow I j \in cp (K j) (J j)$ for $H I K J$ in $allE$)
apply(*simp add:cp-def par-assum-def*)
apply(*erule-tac c=s in subsetD, simp*)
apply *simp*
apply *clarify*
apply(*erule-tac x=i and P= $\lambda j. H j \longrightarrow M \cup \bigcup((T j) ' (S j)) \subseteq (L j)$ for $H M S T L$ in $allE$)
apply *simp*
apply(*erule subsetD*)
apply *simp*
apply(*simp add:conjoin-def compat-label-def*)
apply *clarify*
apply(*erule-tac x=ia and P= $\lambda j. H j \longrightarrow (P j) \vee Q j$ for $H P Q$ in $allE, simp$)
— each etran in $\sigma^{-1}[0..m]$ corresponds to
apply(*erule disjE*)
— a c-tran in some $\sigma\text{-}\{ib\}$
apply *clarify*
apply(*case-tac i=ib, simp*)
apply(*erule etranE, simp*)
apply(*erule-tac x=ib and P= $\lambda i. H i \longrightarrow (I i) \vee (J i)$ for $H I J$ in $allE$)
apply (*erule etranE*)
apply(*case-tac ia=m, simp*)
apply *simp*
apply(*erule-tac x=ia and P= $\lambda j. H j \longrightarrow (\forall i. P i j)$ for $H P$ in $allE$)
apply(*subgoal-tac ia<m, simp*)
prefer 2
apply *arith*
apply(*erule-tac x=ib and P= $\lambda j. (I j, H j) \in ctran \longrightarrow P i j$ for $I H P$ in $allE, simp$)
apply(*simp add:same-state-def*)
apply(*erule-tac x=i and P= $\lambda j. (T j) \longrightarrow (\forall i. (H j i) \longrightarrow (snd (d j i))=(snd (e j i)))$ for $T H d e$ in $all-dupE$)
apply(*erule-tac x=ib and P= $\lambda j. (T j) \longrightarrow (\forall i. (H j i) \longrightarrow (snd (d j i))=(snd (e j i)))$ for $T H d e$ in $allE, simp$)
— or an e-tran in σ , therefore it satisfies *rely* \vee *guar*- $\{ib\}$
apply (*force simp add:par-assum-def same-state-def*)
done*********

lemma *three* [*rule-format*]:

$\llbracket xs \neq []; \forall i < \text{length } xs. \text{rely} \cup (\bigcup j \in \{j. j < \text{length } xs \wedge j \neq i\}. \text{Guar } (xs ! j))$

$\subseteq \text{Rely } (xs ! i)$;
 $pre \subseteq (\bigcap i \in \{i. i < \text{length } xs\}. Pre (xs ! i))$;
 $\forall i < \text{length } xs.$
 $\models Com (xs ! i) \text{ sat } [Pre (xs ! i), Rely (xs ! i), Guar (xs ! i), Post (xs ! i)]$;
 $\text{length } xs = \text{length } clist$; $x \in \text{par-cp } (ParallelCom \text{ } xs) \text{ } s$; $x \in \text{par-assum}(pre, rely)$;
 $\forall i < \text{length } clist. clist!i \in cp (Some(Com(xs!i))) \text{ } s$; $x \propto clist \]$
 $\implies \forall j \text{ } i. i < \text{length } clist \wedge Suc \ j < \text{length } x \longrightarrow (clist!i!j) -e \rightarrow (clist!i!Suc \ j)$
 $\longrightarrow (snd(clist!i!j), snd(clist!i!Suc \ j)) \in rely \cup (\bigcup j \in \{j. j < \text{length } xs \wedge j \neq i\}.$
 $Guar (xs ! j))$
apply(drule two)
apply simp-all
apply clarify
apply(simp add: conjoin-def compat-label-def)
apply clarify
apply(erule-tac $x=j$ and $P=\lambda j. H \ j \longrightarrow (J \ j \wedge (\exists i. P \ i \ j)) \vee I \ j$ for $H \ J \ P \ I$ in
 $allE, simp$)
apply(erule disjE)
prefer 2
apply(force simp add: same-state-def par-assum-def)
apply clarify
apply(case-tac $i=ia, simp$)
apply(erule etranE, simp)
apply(erule-tac $x=ia$ and $P=\lambda i. H \ i \longrightarrow (I \ i) \vee (J \ i)$ for $H \ I \ J$ in $allE, simp$)
apply(erule-tac $x=j$ and $P=\lambda j. \forall i. S \ j \ i \longrightarrow (I \ j \ i, H \ j \ i) \in ctran \longrightarrow P \ i \ j$ for
 $S \ I \ H \ P$ in $allE$)
apply(erule-tac $x=ia$ and $P=\lambda j. S \ j \longrightarrow (I \ j, H \ j) \in ctran \longrightarrow P \ j$ for $S \ I \ H \ P$
in $allE$)
apply(simp add: same-state-def)
apply(erule-tac $x=i$ and $P=\lambda j. T \ j \longrightarrow (\forall i. H \ j \ i \longrightarrow (snd \ (d \ j \ i)) = (snd \ (e \ j \ i)))$ for $T \ H \ d \ e$ in $all-dupE$)
apply(erule-tac $x=ia$ and $P=\lambda j. T \ j \longrightarrow (\forall i. H \ j \ i \longrightarrow (snd \ (d \ j \ i)) = (snd \ (e \ j \ i)))$ for $T \ H \ d \ e$ in $allE, simp$)
done

lemma four:

$\llbracket xs \neq [] \rrbracket$; $\forall i < \text{length } xs. rely \cup (\bigcup j \in \{j. j < \text{length } xs \wedge j \neq i\}. Guar (xs ! j))$
 $\subseteq \text{Rely } (xs ! i)$;
 $(\bigcup j \in \{j. j < \text{length } xs\}. Guar (xs ! j)) \subseteq guar$;
 $pre \subseteq (\bigcap i \in \{i. i < \text{length } xs\}. Pre (xs ! i))$;
 $\forall i < \text{length } xs.$
 $\models Com (xs ! i) \text{ sat } [Pre (xs ! i), Rely (xs ! i), Guar (xs ! i), Post (xs ! i)]$;
 $x \in \text{par-cp } (ParallelCom \text{ } xs) \text{ } s$; $x \in \text{par-assum } (pre, rely)$; $Suc \ i < \text{length } x$;
 $x ! i -pc \rightarrow x ! Suc \ i \]$
 $\implies (snd \ (x ! i), snd \ (x ! Suc \ i)) \in guar$
apply(simp add: ParallelCom-def)
apply(subgoal-tac (map (Some o fst) xs) $\neq []$)
prefer 2
apply simp
apply(frule rev-subsetD)

```

apply(erule one [THEN equalityD1])
apply(erule subsetD)
apply simp
apply clarify
apply(drule-tac pre=pre and rely=rely and x=x and s=s and xs=xs and clist=clist
in two)
apply(assumption+)
  apply(erule sym)
  apply(simp add:ParallelCom-def)
  apply assumption
  apply(simp add:Com-def)
apply assumption
apply(simp add:conjoin-def same-program-def)
apply clarify
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow \text{fst}(I j)=(J j)$  for H I J in all-dupE)
apply(erule-tac x=Suc i and P= $\lambda j. H j \longrightarrow \text{fst}(I j)=(J j)$  for H I J in allE)
apply(erule par-ctranE,simp)
apply(erule-tac x=i and P= $\lambda j. \forall i. S j i \longrightarrow (I j i, H j i) \in \text{ctran} \longrightarrow P i j$  for
S I H P in allE)
apply(erule-tac x=ia and P= $\lambda j. S j \longrightarrow (I j, H j) \in \text{ctran} \longrightarrow P j$  for S I H P
in allE)
apply(rule-tac x=ia in exI)
apply(simp add:same-state-def)
apply(erule-tac x=ia and P= $\lambda j. T j \longrightarrow (\forall i. H j i \longrightarrow (\text{snd}(d j i))=(\text{snd}(e j
i)))$  for T H d e in all-dupE,simp)
apply(erule-tac x=ia and P= $\lambda j. T j \longrightarrow (\forall i. H j i \longrightarrow (\text{snd}(d j i))=(\text{snd}(e j
i)))$  for T H d e in allE,simp)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{snd}(d j))=(\text{snd}(e j))$  for H d e in
all-dupE)
apply(erule-tac x=i and P= $\lambda j. H j \longrightarrow (\text{snd}(d j))=(\text{snd}(e j))$  for H d e in
all-dupE,simp)
apply(erule-tac x=Suc i and P= $\lambda j. H j \longrightarrow (\text{snd}(d j))=(\text{snd}(e j))$  for H d e in
allE,simp)
apply(erule mp)
apply(subgoal-tac r=fst(clist ! ia ! Suc i),simp)
apply(drule-tac i=ia in list-eq-if)
back
apply simp-all
done

```

```

lemma parcptn-not-empty [simp]: []  $\notin$  par-cptn
apply(force elim:par-cptn.cases)
done

```

lemma five:

```

[[xs $\neq$ []];  $\forall i < \text{length } xs. \text{rely} \cup (\bigcup j \in \{j. j < \text{length } xs \wedge j \neq i\}. \text{Guar}(xs ! j))$ 
 $\subseteq \text{Rely}(xs ! i)$ ;
pre  $\subseteq (\bigcap i \in \{i. i < \text{length } xs\}. \text{Pre}(xs ! i))$ ;
 $(\bigcap i \in \{i. i < \text{length } xs\}. \text{Post}(xs ! i)) \subseteq \text{post}$ ;

```

```

 $\forall i < \text{length } xs.$ 
 $\models \text{Com } (xs ! i) \text{ sat } [\text{Pre } (xs ! i), \text{Rely } (xs ! i), \text{Guar } (xs ! i), \text{Post } (xs ! i)];$ 
 $x \in \text{par-cp } (\text{ParallelCom } xs) \text{ s}; x \in \text{par-assum } (\text{pre}, \text{rely});$ 
 $\text{All-None } (\text{fst } (\text{last } x)) \text{ ] } \implies \text{snd } (\text{last } x) \in \text{post}$ 
apply(simp add: ParallelCom-def)
apply(subgoal-tac (map (Some  $\circ$  fst) xs) $\neq$ [])
prefer 2
apply simp
apply(frule rev-subsetD)
apply(erule one [THEN equalityD1])
apply(erule subsetD)
apply simp
apply clarify
apply(subgoal-tac  $\forall i < \text{length } \text{clist}. \text{clist}!i \in \text{assum}(\text{Pre}(xs!i), \text{Rely}(xs!i))$ )
apply(erule-tac  $x=xa$  and  $P=\lambda i. H i \longrightarrow \models (J i) \text{ sat } [I i, K i, M i, N i]$  for  $H J$ 
 $I K M N$  in  $\text{allE}, \text{erule impE}, \text{assumption}$ )
apply(simp add: com-validity-def)
apply(erule-tac  $x=s$  in  $\text{allE}$ )
apply(erule-tac  $x=xa$  and  $P=\lambda j. H j \longrightarrow (I j) \in \text{cp } (J j) \text{ s}$  for  $H I J$  in  $\text{allE}, \text{simp}$ )
apply(drule-tac  $c=\text{clist}!xa$  in  $\text{subsetD}$ )
apply (force simp add: Com-def)
apply(simp add: comm-def conjoin-def same-program-def del: last.simps)
apply clarify
apply(erule-tac  $x=\text{length } x - 1$  and  $P=\lambda j. H j \longrightarrow \text{fst}(I j)=(J j)$  for  $H I J$  in
 $\text{allE}$ )
apply (simp add: All-None-def same-length-def)
apply(erule-tac  $x=xa$  and  $P=\lambda j. H j \longrightarrow \text{length}(J j)=(K j)$  for  $H J K$  in  $\text{allE}$ )
apply(subgoal-tac  $\text{length } x - 1 < \text{length } x, \text{simp}$ )
apply(case-tac  $x \neq []$ )
apply(simp add: last-conv-nth)
apply(erule-tac  $x=\text{clist}!xa$  in  $\text{ballE}$ )
apply(simp add: same-state-def)
apply(subgoal-tac  $\text{clist}!xa \neq []$ )
apply(simp add: last-conv-nth)
apply(case-tac  $x$ )
apply (force simp add: par-cp-def)
apply (force simp add: par-cp-def)
apply force
apply (force simp add: par-cp-def)
apply(case-tac  $x$ )
apply (force simp add: par-cp-def)
apply (force simp add: par-cp-def)
apply clarify
apply(simp add: assum-def)
apply(rule conjI)
apply(simp add: conjoin-def same-state-def par-cp-def)
apply clarify
apply(erule-tac  $x=i$  and  $P=\lambda j. T j \longrightarrow (\forall i. H j i \longrightarrow (\text{snd } (d j i))=(\text{snd } (e j$ 
 $i)))$  for  $T H d e$  in  $\text{allE}, \text{simp}$ )

```

```

apply(erule-tac  $x=0$  and  $P=\lambda j. H j \longrightarrow (snd (d j))=(snd (e j))$  for  $H d e$  in
allE)
apply(case-tac  $x, simp+$ )
apply (simp add:par-assum-def)
apply clarify
apply(erule-tac  $c=snd (clist ! i ! 0)$  in subsetD)
apply assumption
apply simp
apply clarify
apply(erule-tac  $x=i$  in all-dupE)
apply(rule subsetD, erule mp, assumption)
apply(erule-tac  $pre=pre$  and  $rely=rely$  and  $x=x$  and  $s=s$  in three)
apply(erule-tac  $x=ib$  in allE,erule mp)
apply simp-all
apply(simp add:ParallelCom-def)
apply(force simp add:Com-def)
apply(simp add:conjoin-def same-length-def)
done

```

```

lemma ParallelEmpty [rule-format]:
   $\forall i s. x \in par-cp (ParallelCom []) s \longrightarrow$ 
   $Suc i < length x \longrightarrow (x ! i, x ! Suc i) \notin par-ctran$ 
apply(induct-tac  $x$ )
apply(simp add:par-cp-def ParallelCom-def)
apply clarify
apply(case-tac list, simp, simp)
apply(case-tac  $i$ )
apply(simp add:par-cp-def ParallelCom-def)
apply(erule par-ctranE, simp)
apply(simp add:par-cp-def ParallelCom-def)
apply clarify
apply(erule par-cptn.cases, simp)
apply simp
apply(erule par-ctranE)
back
apply simp
done

```

```

theorem par-rgsound:
   $\vdash c SAT [pre, rely, guar, post] \implies$ 
   $\models (ParallelCom c) SAT [pre, rely, guar, post]$ 
apply(erule par-rghoare.induct)
apply(case-tac  $xs, simp$ )
apply(simp add:par-com-validity-def par-comm-def)
apply clarify
apply(case-tac post=UNIV, simp)
apply clarify
apply(erule ParallelEmpty)
apply assumption

```

```

  apply simp
  apply clarify
  apply simp
  apply(subgoal-tac xs≠[])
  prefer 2
  apply simp
  apply(rename-tac a list)
  apply(thin-tac xs = a # list)
  apply(simp add:par-com-validity-def par-comm-def)
  apply clarify
  apply(rule conjI)
  apply clarify
  apply(erule-tac pre=pre and rely=rely and guar=guar and x=x and s=s and
xs=xs in four)
    apply(assumption+)
    apply clarify
    apply (erule allE, erule impE, assumption,erule rgsound)
    apply(assumption+)
  apply clarify
  apply(erule-tac pre=pre and rely=rely and post=post and x=x and s=s and
xs=xs in five)
    apply(assumption+)
    apply clarify
    apply (erule allE, erule impE, assumption,erule rgsound)
    apply(assumption+)
done

end

```

3.6 Concrete Syntax

```

theory RG-Syntax
imports RG-Hoare Quote-Antiquote
begin

```

```

abbreviation Skip :: 'a com (⟨SKIP⟩)
  where SKIP ≡ Basic id

```

```

notation Seq (⟨(-;/-)⟩ [60,61] 60)

```

```

syntax

```

```

-Assign    :: idt ⇒ 'b ⇒ 'a com                (⟨' - :=/ -⟩ [70, 65] 61)
-Cond      :: 'a bexp ⇒ 'a com ⇒ 'a com ⇒ 'a com (⟨(OIF -/ THEN -/ ELSE
-/FI)⟩ [0, 0, 0] 61)
-Cond2     :: 'a bexp ⇒ 'a com ⇒ 'a com          (⟨(OIF - THEN - FI)⟩ [0,0]
56)
-While     :: 'a bexp ⇒ 'a com ⇒ 'a com          (⟨(OWHILE - /DO - /OD)⟩
[0, 0] 61)
-Await     :: 'a bexp ⇒ 'a com ⇒ 'a com          (⟨(OAWAIT - /THEN -/

```


$\langle \text{/END} \rangle \ [0,0] \ 61)$
 $\text{-Atom} \quad \quad \quad \text{:: } 'a \text{ com} \Rightarrow 'a \text{ com} \quad \quad \quad (\langle \langle (-) \rangle \rangle \ 61)$
 $\text{-Wait} \quad \quad \quad \text{:: } 'a \text{ bexp} \Rightarrow 'a \text{ com} \quad \quad \quad (\langle \langle 0\text{WAIT} - \text{END} \rangle \rangle \ 61)$

translations

$'x := a \rightarrow \text{CONST Basic} \langle '(-\text{update-name } x (\lambda-. a)) \rangle$
 $\text{IF } b \text{ THEN } c1 \text{ ELSE } c2 \text{ FI} \rightarrow \text{CONST Cond} \{b\} \ c1 \ c2$
 $\text{IF } b \text{ THEN } c \text{ FI} \Leftrightarrow \text{IF } b \text{ THEN } c \text{ ELSE SKIP FI}$
 $\text{WHILE } b \text{ DO } c \text{ OD} \rightarrow \text{CONST While} \{b\} \ c$
 $\text{AWAIT } b \text{ THEN } c \text{ END} \Leftrightarrow \text{CONST Await} \{b\} \ c$
 $\langle c \rangle \Leftrightarrow \text{AWAIT CONST True THEN } c \text{ END}$
 $\text{WAIT } b \text{ END} \Leftrightarrow \text{AWAIT } b \text{ THEN SKIP END}$

nonterminal prgs

syntax

$\text{-PAR} \quad \quad \quad \text{:: } prgs \Rightarrow 'a \quad \quad \quad (\langle \text{COBEGIN} // - // \text{COEND} \rangle \ 60)$
 $\text{-prg} \quad \quad \quad \text{:: } 'a \Rightarrow prgs \quad \quad \quad (\langle \rightarrow \ 57)$
 $\text{-prgs} \quad \quad \quad \text{:: } ['a, prgs] \Rightarrow prgs \quad \quad \quad (\langle - // || // - \rangle \ [60,57] \ 57)$

translations

$\text{-prg } a \rightarrow [a]$
 $\text{-prgs } a \ ps \rightarrow a \ \# \ ps$
 $\text{-PAR } ps \rightarrow ps$

syntax

$\text{-prg-scheme} \text{:: } ['a, 'a, 'a, 'a] \Rightarrow prgs \ (\langle \text{SCHEME} [- \leq - < -] \rightarrow [0,0,0,60] \ 57)$

translations

$\text{-prg-scheme } j \ i \ k \ c \Leftrightarrow (\text{CONST map } (\lambda i. \ c) \ [j..<k])$

Translations for variables before and after a transition:

syntax

$\text{-before} \text{:: } id \Rightarrow 'a \ (\langle {}^{\circ}\rightarrow)$
 $\text{-after} \text{:: } id \Rightarrow 'a \ (\langle {}^{\text{a}}\rightarrow)$

translations

${}^{\circ}x \Leftrightarrow x \ ' \ \text{CONST fst}$
 ${}^{\text{a}}x \Leftrightarrow x \ ' \ \text{CONST snd}$

print-translation \langle

let
 $\text{fun quote-tr}' f (t :: ts) =$
 $\quad \text{Term.list-comb } (f \ \$ \ \text{Syntax-Trans.quote-tr}' \ \mathbf{syntax-const} \ \langle \text{-antiquote} \rangle \ t,$
 $\text{ts})$
 $\quad | \ \text{quote-tr}' \ - \ - = \text{raise Match};$
 $\text{val assert-tr}' = \text{quote-tr}' (\text{Syntax.const } \mathbf{syntax-const} \ \langle \text{-Assert} \rangle);$

```

    fun bexp-tr' name ((Const (const-syntax <Collect>, -) $ t) :: ts) =
      quote-tr' (Syntax.const name) (t :: ts)
    | bexp-tr' - - = raise Match;

    fun assign-tr' (Abs (x, -, f $ k $ Bound 0) :: ts) =
      quote-tr' (Syntax.const syntax-const <-Assign> $ Syntax-Trans.update-name-tr'
f)
      (Abs (x, dummyT, Syntax-Trans.const-abs-tr' k) :: ts)
    | assign-tr' - - = raise Match;
  in
    [(const-syntax <Collect>, K assert-tr'),
     (const-syntax <Basic>, K assign-tr'),
     (const-syntax <Cond>, K (bexp-tr' syntax-const <-Cond>)),
     (const-syntax <While>, K (bexp-tr' syntax-const <-While>))]
  end
>

end

```

3.7 Examples

```

theory RG-Examples
imports RG-Syntax
begin

```

```

lemmas definitions [simp] = stable-def Pre-def Rely-def Guar-def Post-def Com-def

```

3.7.1 Set Elements of an Array to Zero

```

lemma le-less-trans2: [(j::nat) < k; i ≤ j] ⇒ i < k
by simp

```

```

lemma add-le-less-mono: [(a::nat) < c; b ≤ d] ⇒ a + b < c + d
by simp

```

```

record Example1 =
  A :: nat list

```

```

lemma Example1:
  ⊢ COBEGIN
    SCHEME [0 ≤ i < n]
    ( 'A := 'A [i := 0],
      { n < length 'A },
      { length °A = length °A ∧ °A ! i = °A ! i },
      { length °A = length °A ∧ (∀ j < n. i ≠ j → °A ! j = °A ! j) },
      { 'A ! i = 0 } )
  COEND
  SAT [{ n < length 'A }, { °A = °A }, { True }, { ∀ i < n. 'A ! i = 0 }]
apply(rule Parallel)

```

apply (*auto intro!*: *Basic*)
done

lemma *Example1-parameterized*:

$k < t \implies$

⊢ *COBEGIN*

SCHEME [$k*n \leq i < (Suc\ k)*n$] ($\acute{A} := \acute{A}[i:=0]$),

$\{\{t*n < length\ \acute{A}\}$,

$\{\{t*n < length\ \circ A \wedge length\ \circ A = length\ \mathring{A} \wedge \circ A!i = \mathring{A}!i\}$,

$\{\{t*n < length\ \circ A \wedge length\ \circ A = length\ \mathring{A} \wedge (\forall j < length\ \circ A . i \neq j \longrightarrow \circ A!j = \mathring{A}!j)\}$,

$\{\{ \acute{A}!i=0\}\}$)

COEND

SAT [$\{\{t*n < length\ \acute{A}\}$,

$\{\{t*n < length\ \circ A \wedge length\ \circ A = length\ \mathring{A} \wedge (\forall i < n . \circ A!(k*n+i) = \mathring{A}!(k*n+i))\}$,

$\{\{t*n < length\ \circ A \wedge length\ \circ A = length\ \mathring{A} \wedge$

$(\forall i < length\ \circ A . (i < k*n \longrightarrow \circ A!i = \mathring{A}!i) \wedge ((Suc\ k)*n \leq i \longrightarrow \circ A!i = \mathring{A}!i))\}$,

$\{\{\forall i < n . \acute{A}!(k*n+i) = 0\}\}$]

apply(*rule Parallel*)

apply *auto*

apply(*erule-tac* $x=k*n+i$ **in** *allE*)

apply(*subgoal-tac* $k*n+i < length\ (A\ b)$)

apply *force*

apply(*erule le-less-trans2*)

apply(*case-tac* $t, simp+$)

apply (*simp add: add.commute*)

apply(*simp add: add-le-mono*)

apply(*rule Basic*)

apply *simp*

apply *clarify*

apply (*subgoal-tac* $k*n+i < length\ (A\ x)$)

apply *simp*

apply(*erule le-less-trans2*)

apply(*case-tac* $t, simp+$)

apply (*simp add: add.commute*)

apply(*rule add-le-mono, auto*)

done

3.7.2 Increment a Variable in Parallel

Two components

record *Example2* =

$x :: nat$

$c-0 :: nat$

$c-1 :: nat$

lemma *Example2*:

⊢ *COBEGIN*

($\langle \acute{x} := \acute{x}+1;; \acute{c-0} := \acute{c-0} + 1 \rangle$),

$$\begin{aligned}
& \{ \{ x = 'c-0 + 'c-1 \wedge 'c-0=0 \}, \\
& \{ \{ {}^o c-0 = {}^a c-0 \wedge \\
& \quad ({}^o x = {}^o c-0 + {}^o c-1 \\
& \quad \rightarrow {}^a x = {}^a c-0 + {}^a c-1) \}, \\
& \{ \{ {}^o c-1 = {}^a c-1 \wedge \\
& \quad ({}^o x = {}^o c-0 + {}^o c-1 \\
& \quad \rightarrow {}^a x = {}^a c-0 + {}^a c-1) \}, \\
& \{ \{ x = 'c-0 + 'c-1 \wedge 'c-0=1 \} \} \\
\parallel \\
& (\langle 'x := 'x+1;; 'c-1 := 'c-1+1 \rangle, \\
& \{ \{ x = 'c-0 + 'c-1 \wedge 'c-1=0 \}, \\
& \{ \{ {}^o c-1 = {}^a c-1 \wedge \\
& \quad ({}^o x = {}^o c-0 + {}^o c-1 \\
& \quad \rightarrow {}^a x = {}^a c-0 + {}^a c-1) \}, \\
& \{ \{ {}^o c-0 = {}^a c-0 \wedge \\
& \quad ({}^o x = {}^o c-0 + {}^o c-1 \\
& \quad \rightarrow {}^a x = {}^a c-0 + {}^a c-1) \}, \\
& \{ \{ x = 'c-0 + 'c-1 \wedge 'c-1=1 \} \} \\
& \text{COEND} \\
& \text{SAT} [\{ \{ x=0 \wedge 'c-0=0 \wedge 'c-1=0 \}, \\
& \quad \{ \{ {}^o x = {}^a x \wedge {}^o c-0 = {}^a c-0 \wedge {}^o c-1 = {}^a c-1 \}, \\
& \quad \{ \text{True} \}, \\
& \quad \{ \{ x=2 \} \}] \\
& \text{apply}(\text{rule Parallel}) \\
& \quad \text{apply simp-all} \\
& \quad \text{apply clarify} \\
& \quad \text{apply}(\text{case-tac } i) \\
& \quad \text{apply simp} \\
& \quad \text{apply}(\text{rule conjI}) \\
& \quad \text{apply clarify} \\
& \quad \text{apply simp} \\
& \quad \text{apply clarify} \\
& \quad \text{apply simp} \\
& \quad \text{apply simp} \\
& \quad \text{apply}(\text{rule conjI}) \\
& \quad \text{apply clarify} \\
& \quad \text{apply simp} \\
& \quad \text{apply clarify} \\
& \quad \text{apply simp} \\
& \quad \text{apply}(\text{subgoal-tac } x=0) \\
& \quad \text{apply simp} \\
& \quad \text{apply arith} \\
& \quad \text{apply clarify} \\
& \quad \text{apply}(\text{case-tac } xa, \text{simp}, \text{simp}) \\
& \quad \text{apply clarify} \\
& \quad \text{apply simp} \\
& \quad \text{apply}(\text{erule-tac } x=0 \text{ in all-dupE}) \\
& \quad \text{apply}(\text{erule-tac } x=1 \text{ in allE,simp}) \\
& \quad \text{apply clarify}
\end{aligned}$$

```

apply(case-tac i, simp)
apply(rule Await)
  apply simp-all
apply(clarify)
apply(rule Seq)
  prefer 2
  apply(rule Basic)
    apply simp-all
  apply(rule subset-refl)
apply(rule Basic)
apply simp-all
apply clarify
apply simp
apply(rule Await)
  apply simp-all
apply(clarify)
apply(rule Seq)
  prefer 2
  apply(rule Basic)
    apply simp-all
  apply(rule subset-refl)
apply(auto intro! : Basic)
done

```

Parameterized

```

lemma Example2-lemma2-aux: j < n  $\implies$ 
  ( $\sum i=0..<n. (b \ i::nat)$ ) =
  ( $\sum i=0..<j. b \ i$ ) +  $b \ j$  + ( $\sum i=0..<n-(Suc \ j) . b \ (Suc \ j + i)$ )
apply(induct n)
apply simp-all
apply(simp add:less-Suc-eq)
apply(auto)
apply(subgoal-tac n - j = Suc(n- Suc j))
  apply simp
apply arith
done

```

```

lemma Example2-lemma2-aux2:
   $j \leq s \implies (\sum i::nat=0..<j. (b \ (s:=t)) \ i) = (\sum i=0..<j. b \ i)$ 
  by (induct j) simp-all

```

```

lemma Example2-lemma2:
   $\llbracket j < n; b \ j = 0 \rrbracket \implies Suc \ (\sum i::nat=0..<n. b \ i) = (\sum i=0..<n. (b \ (j := Suc \ 0)) \ i)$ 
apply(frule-tac b=(b \ (j:=(Suc \ 0))) in Example2-lemma2-aux)
apply(erule-tac t=sum \ b \ (j := (Suc \ 0)) \ {0..<n} in ssubst)
apply(frule-tac b=b in Example2-lemma2-aux)
apply(erule-tac t=sum \ b \ {0..<n} in ssubst)
apply(subgoal-tac Suc \ (sum \ b \ {0..<j} + b \ j + (\sum i=0..<n - Suc \ j. b \ (Suc \ j +

```

```

i)))=(sum b {0..<j} + Suc (b j) + (sum i=0..<n - Suc j. b (Suc j + i))))
apply(rotate-tac -1)
apply(erule ssubst)
apply(subgoal-tac j≤j)
  apply(drule-tac b=b and t=(Suc 0) in Example2-lemma2-aux2)
apply(rotate-tac -1)
apply(erule ssubst)
apply simp-all
done

```

```

lemma Example2-lemma2-Suc0:  $\llbracket j < n; b\ j = 0 \rrbracket \implies$ 
  Suc (sum i::nat=0..<n. b i)=(sum i=0..<n. (b (j:=Suc 0)) i)
by(simp add:Example2-lemma2)

```

```

record Example2-parameterized =
  C :: nat ⇒ nat
  y :: nat

```

```

lemma Example2-parameterized:  $0 < n \implies$ 
  ⊢ COBEGIN SCHEME  $[0 \leq i < n]$ 
    ( $\langle \text{' } y := \text{' } y + 1;; \text{' } C := \text{' } C\ (i := 1) \rangle$ ),
     $\{\text{' } y = (\sum i = 0..<n. \text{' } C\ i) \wedge \text{' } C\ i = 0\}$ ,
     $\{\text{' } C\ i = \text{' } C\ i \wedge$ 
       $(\text{' } y = (\sum i = 0..<n. \text{' } C\ i) \longrightarrow \text{' } y = (\sum i = 0..<n. \text{' } C\ i))\}$ ,
     $\{\forall j < n. i \neq j \longrightarrow \text{' } C\ j = \text{' } C\ j\} \wedge$ 
       $(\text{' } y = (\sum i = 0..<n. \text{' } C\ i) \longrightarrow \text{' } y = (\sum i = 0..<n. \text{' } C\ i))\}$ ,
     $\{\text{' } y = (\sum i = 0..<n. \text{' } C\ i) \wedge \text{' } C\ i = 1\}$ 
  COEND
  SAT  $[\{\text{' } y = 0 \wedge (\sum i = 0..<n. \text{' } C\ i) = 0\}, \{\text{' } C = \text{' } C \wedge \text{' } y = \text{' } y\}, \{\text{True}\}, \{\text{' } y = n\}]$ 
apply(rule Parallel)
apply force
apply force
apply(force)
apply clarify
apply simp
apply simp
apply clarify
apply simp
apply(rule Await)
apply simp-all
apply clarify
apply(rule Seq)
prefer 2
apply(rule Basic)
apply(rule subset-refl)
apply simp+
apply(rule Basic)
apply simp
apply clarify

```

```

apply simp
apply(simp add:Example2-lemma2-Suc0 cong:if-cong)
apply simp-all
done

```

3.7.3 Find Least Element

A previous lemma:

lemma *mod-aux* : $\llbracket i < (n::nat); a \bmod n = i; j < a + n; j \bmod n = i; a < j \rrbracket \implies \text{False}$

```

apply(subgoal-tac a=a div n*n + a mod n )
prefer 2 apply (simp (no-asm-use))
apply(subgoal-tac j=j div n*n + j mod n)
prefer 2 apply (simp (no-asm-use))
apply simp
apply(subgoal-tac a div n*n < j div n*n)
prefer 2 apply arith
apply(subgoal-tac j div n*n < (a div n + 1)*n)
prefer 2 apply simp
apply (simp only:mult-less-cancel2)
apply arith
done

```

```

record Example3 =
  X :: nat => nat
  Y :: nat => nat

```

lemma *Example3*: $m \bmod n = 0 \implies$

```

COBEGIN
SCHEME [0 ≤ i < n]
( WHILE (∀ j < n. 'X i < 'Y j) DO
  IF P(B!( 'X i)) THEN 'Y := 'Y (i := 'X i)
  ELSE 'X := 'X (i := ('X i) + n) FI
OD,
{ ('X i) mod n = i ∧ (∀ j < 'X i. j mod n = i → ¬P(B!j)) ∧ ('Y i < m → P(B!( 'Y i)) ∧ 'Y i ≤ m+i) },
{ (∀ j < n. i ≠ j → aY j ≤ oY j) ∧ oX i = aX i ∧ oY i = aY i },
{ (∀ j < n. i ≠ j → oX j = aX j ∧ oY j = aY j) ∧ aY i ≤ oY i },
{ ('X i) mod n = i ∧ (∀ j < 'X i. j mod n = i → ¬P(B!j)) ∧ ('Y i < m → P(B!( 'Y i)) ∧ 'Y i ≤ m+i) ∧ (∃ j < n. 'Y j ≤ 'X i) }
COEND
SAT [ { ∀ i < n. 'X i = i ∧ 'Y i = m+i }, { oX = aX ∧ oY = aY }, { True },
{ ∀ i < n. ('X i) mod n = i ∧ (∀ j < 'X i. j mod n = i → ¬P(B!j)) ∧ ('Y i < m → P(B!( 'Y i)) ∧ 'Y i ≤ m+i) ∧ (∃ j < n. 'Y j ≤ 'X i) } ]
apply(rule Parallel)
— 5 subgoals left
apply force+

```

```

apply clarify
apply simp
apply(rule While)
  apply force
  apply force
  apply force
  apply (erule dvdE)
  apply(rule-tac pre'={ 'X i mod n = i ∧ (∀j. j < 'X i → j mod n = i →
¬P(B!j)) ∧ ('Y i < n * k → P (B!( 'Y i))) ∧ 'X i < 'Y i } in Conseq)
  apply force
  apply(rule subset-refl)+
apply(rule Cond)
  apply force
  apply(rule Basic)
  apply force
  apply fastforce
  apply force
  apply force
apply(rule Basic)
  apply simp
  apply clarify
  apply simp
  apply (case-tac X x (j mod n) ≤ j)
  apply (erule le-imp-less-or-eq)
  apply (erule disjE)
  apply (erule-tac j=j and n=n and i=j mod n and a=X x (j mod n) in
mod-aux)
  apply auto
done

```

Same but with a list as auxiliary variable:

```

record Example3-list =
  X :: nat list
  Y :: nat list

```

```

lemma Example3-list: m mod n = 0 ⇒ ⊢ (COBEGIN SCHEME [0 ≤ i < n]
(WHILE (∀j < n. 'X!i < 'Y!j) DO
  IF P(B!( 'X!i)) THEN 'Y := 'Y[i := 'X!i] ELSE 'X := 'X[i := ('X!i) + n] FI
OD,
{ n < length 'X ∧ n < length 'Y ∧ ('X!i) mod n = i ∧ (∀j < 'X!i. j mod n = i →
¬P(B!j)) ∧ ('Y!i < m → P(B!( 'Y!i)) ∧ 'Y!i ≤ m + i }},
{ (∀j < n. i ≠ j → aY!j ≤ oY!j) ∧ oX!i = aX!i ∧
oY!i = aY!i ∧ length oX = length aX ∧ length oY = length aY },
{ (∀j < n. i ≠ j → oX!j = aX!j ∧ oY!j = aY!j) ∧
aY!i ≤ oY!i ∧ length oX = length aX ∧ length oY = length aY },
{ ('X!i) mod n = i ∧ (∀j < 'X!i. j mod n = i → ¬P(B!j)) ∧ ('Y!i < m → P(B!( 'Y!i))
∧ 'Y!i ≤ m + i) ∧ (∃j < n. 'Y!j ≤ 'X!i) } ) COEND)
SAT [ { n < length 'X ∧ n < length 'Y ∧ (∀i < n. 'X!i = i ∧ 'Y!i = m + i) },
{ oX = aX ∧ oY = aY },

```



```

    {True},
    { $\forall i < n. ('X!i) \text{ mod } n = i \wedge (\forall j < 'X!i. j \text{ mod } n = i \longrightarrow \neg P(B!j)) \wedge$ 
      ( $'Y!i < m \longrightarrow P(B!('Y!i)) \wedge 'Y!i \leq m+i \wedge (\exists j < n. 'Y!j \leq 'X!i)$ )}]
  apply (rule Parallel)
apply (auto cong del: image-cong-simp)
apply force
apply (rule While)
  apply force
  apply force
  apply force
  apply (erule dvdE)
  apply (rule-tac pre'= $\{n < \text{length } 'X \wedge n < \text{length } 'Y \wedge 'X!i \text{ mod } n = i \wedge (\forall j. j$ 
     $< 'X!i \longrightarrow j \text{ mod } n = i \longrightarrow \neg P(B!j)) \wedge ('Y!i < n * k \longrightarrow P(B!('Y!$ 
     $i)) \wedge 'X!i < 'Y!i\}$  in Conseq)
    apply force
    apply (rule subset-refl)+
  apply (rule Cond)
    apply force
  apply (rule Basic)
    apply force
    apply force
    apply force
    apply force
  apply (rule Basic)
    apply simp
    apply clarify
    apply simp
    apply (rule allI)
    apply (rule impI)+
    apply (case-tac  $X\ x!\ i \leq j$ )
      apply (drule le-imp-less-or-eq)
      apply (erule disjE)
      apply (drule-tac  $j=j$  and  $n=n$  and  $i=i$  and  $a=X\ x!\ i$  in mod-aux)
      apply auto
done

end
theory Hoare-Parallel
imports OG-Examples Gar-Coll Mul-Gar-Coll RG-Examples
begin

end

```

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